

24) Aspect-Oriented Programming with Aspect/J

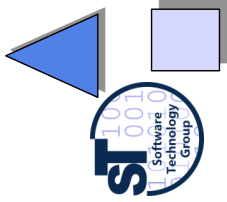
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1. The Problem of Crosscutting
2. Aspect-Oriented Programming
3. Composition Operators and Point-Cuts
4. AOSD
5. Evaluation as Composition System

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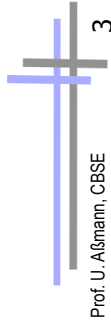
Literature

- ▶ <http://www.eclipse.org/aspectj/>
- ▶ <http://aosd.net/>
- ▶ [KLM+97] G. Kiczales, J. Lamping, A. Mendhekar, C. Maeda, C. Videira Lopes, J.-M. Loingtier, J. Irwin. *Aspect-Oriented Programming*. 1997
- ▶ R. Laddad. *Aspect/J in Action*. Manning Publishers. 2003. Book with many details and applications of Aspect/J.



Other literature

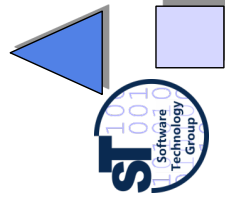
- C. V. Lopes. *Aspect-Oriented Programming: An Historical Perspective (What's in a Name?)*. 2002
http://www.isr.uci.edu/tech_reports/UCI-ISR-02-5.pdf
- G. Kiczales. *Aspect Oriented Programming - Radical Research in Modularity*. Google Tech Talk, 57 min
<http://video.google.com/videosearch?q=Kiczales>
- Jendrik Johannes. *Component-Based Model-Driven Software Development*. PhDthesis, Dresden University of Technology, December 2010.
- Jendrik Johannes and Uwe Aßmann. *Concern-based (de)composition of model-driven software development processes*. In D. C. Petriu, N. Rouquette, and O.Haugen, editors, *MoDELS (2)*, volume 6395 of *Lecture Notes in Computer Science*, pages 47-62. Springer, 2010.



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24.1 The Problem of Crosscutting



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XML parsing in org.apache.tomcat



[Picture taken from the aspectj.org website]

Good modularity:

The „red“ concern is handled by code in one class



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URL pattern matching in org.apache.tomcat



[Picture taken from the aspectj.org website]

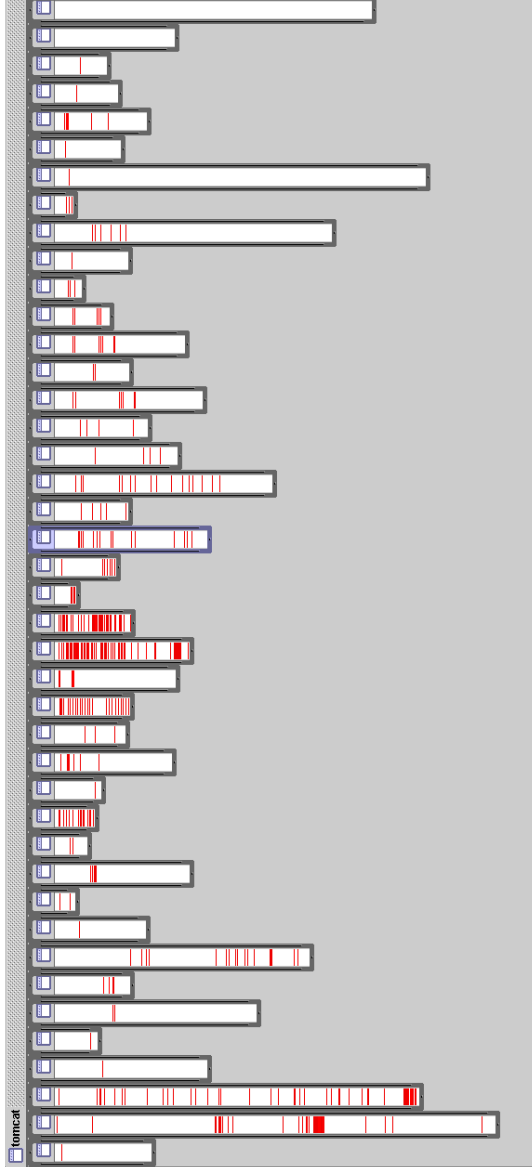
Good modularity:

The “red” concern is handled by code in two classes related by inheritance



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[Picture taken from the aspectj.org website]

BAD modularity:

The concern is handled by code that is scattered over almost all classes

Comparison

Bad modularity

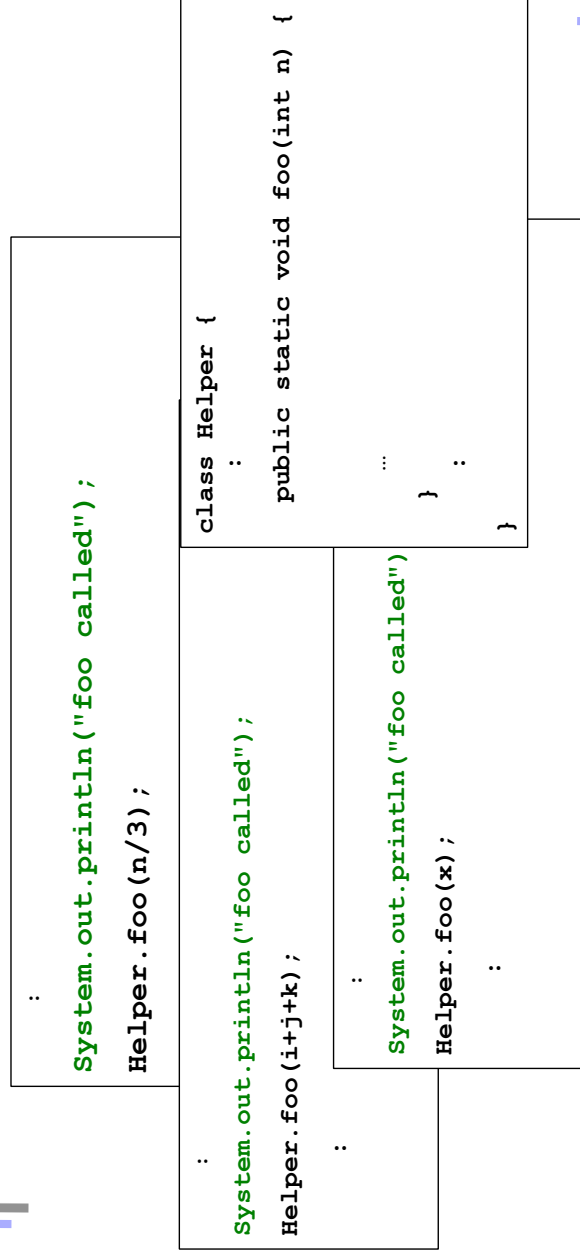
- ▶ **scattering** – code addressing one concern is spread around in the code
- ▶ **tangling** – code in one region addresses multiple concerns
- ▶ Scattering and tangling appear together; they describe different facets of the same problem
 - redundant code
 - difficult to reason about
 - difficult to change

Good Modularity

- ▶ **separated** – implementation of a concern can be treated as relatively separate entity
- ▶ **localized** – implementation of a concern appears in one part of program
- ▶ **modular** – above + has a clear, well defined interface to rest of system

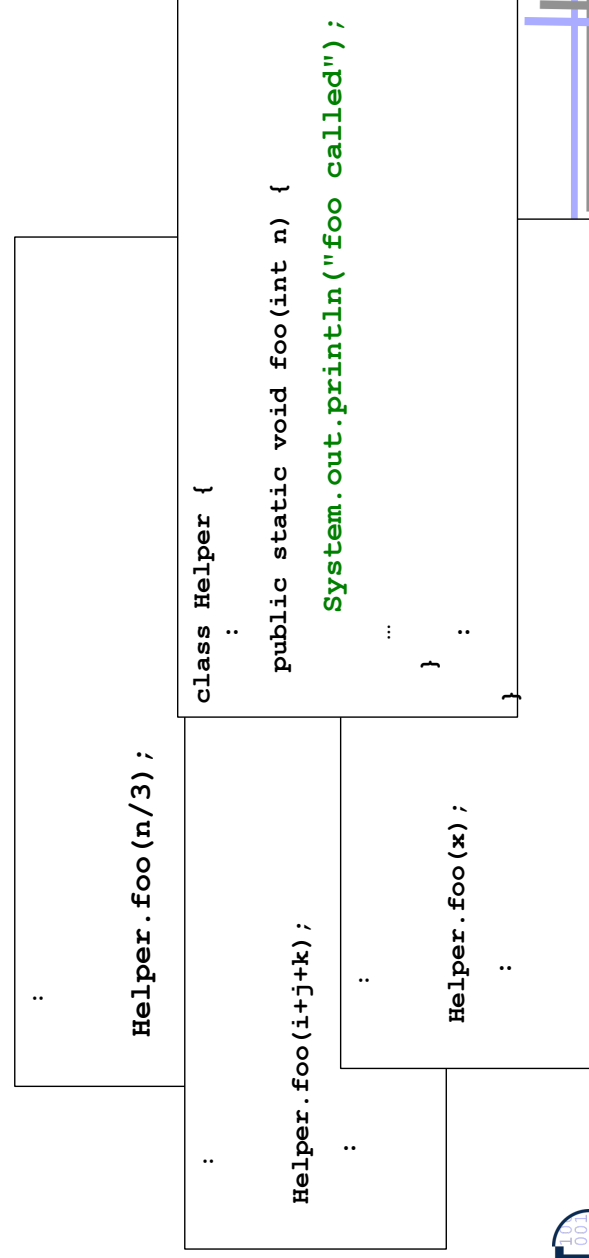
A first example for scattering

- ▶ every call to foo is preceded by a log call (scattering)



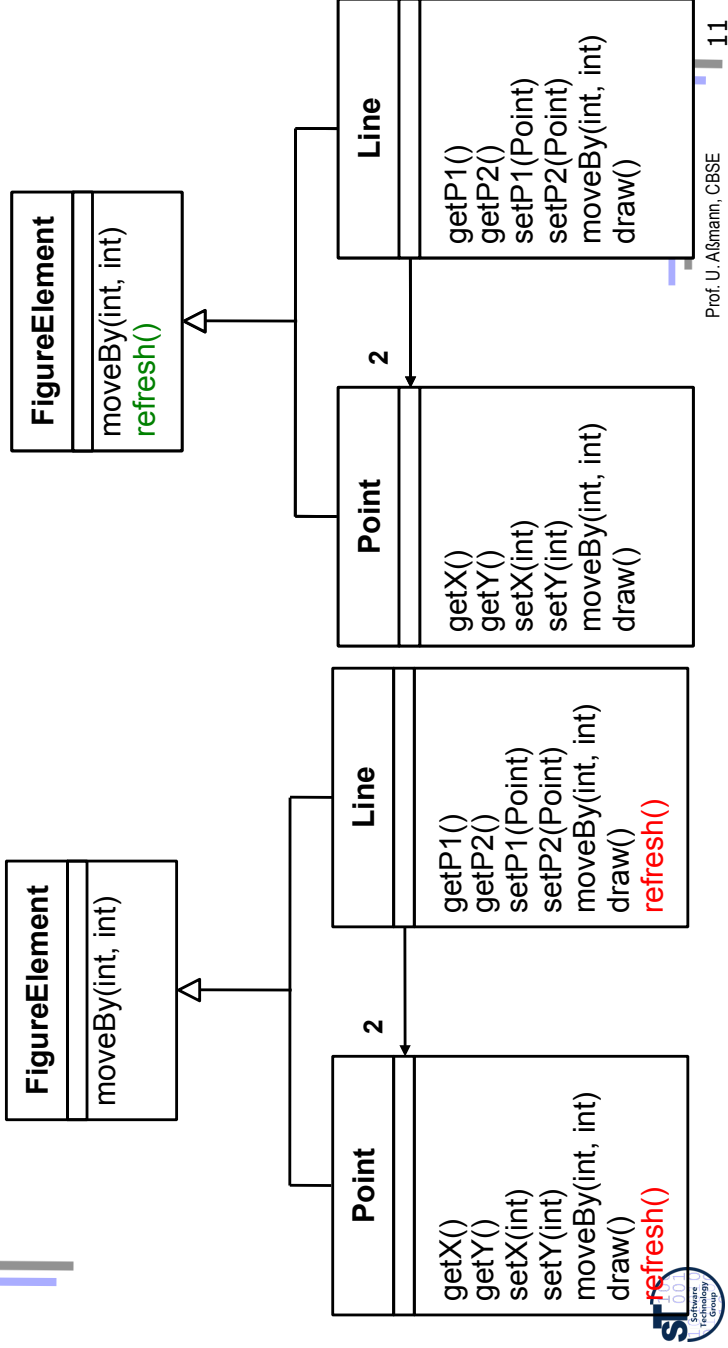
Solution: Refactoring of Scattered Calls

- ▶ Procedures can modularize this case (unless logs use calling context)
- ▶ Scattered calls can be refactored *into* called procedures



A second example of S&T

- ▶ all subclasses have an identical method
- inheritance can modularize this
- Refactoring **moveUpMethod**



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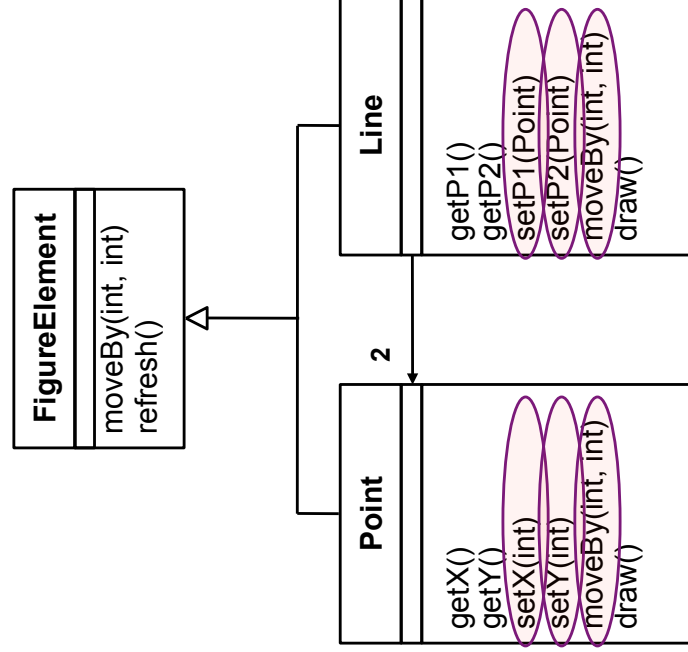
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A Final Example of S&T in the Implementation of Methods

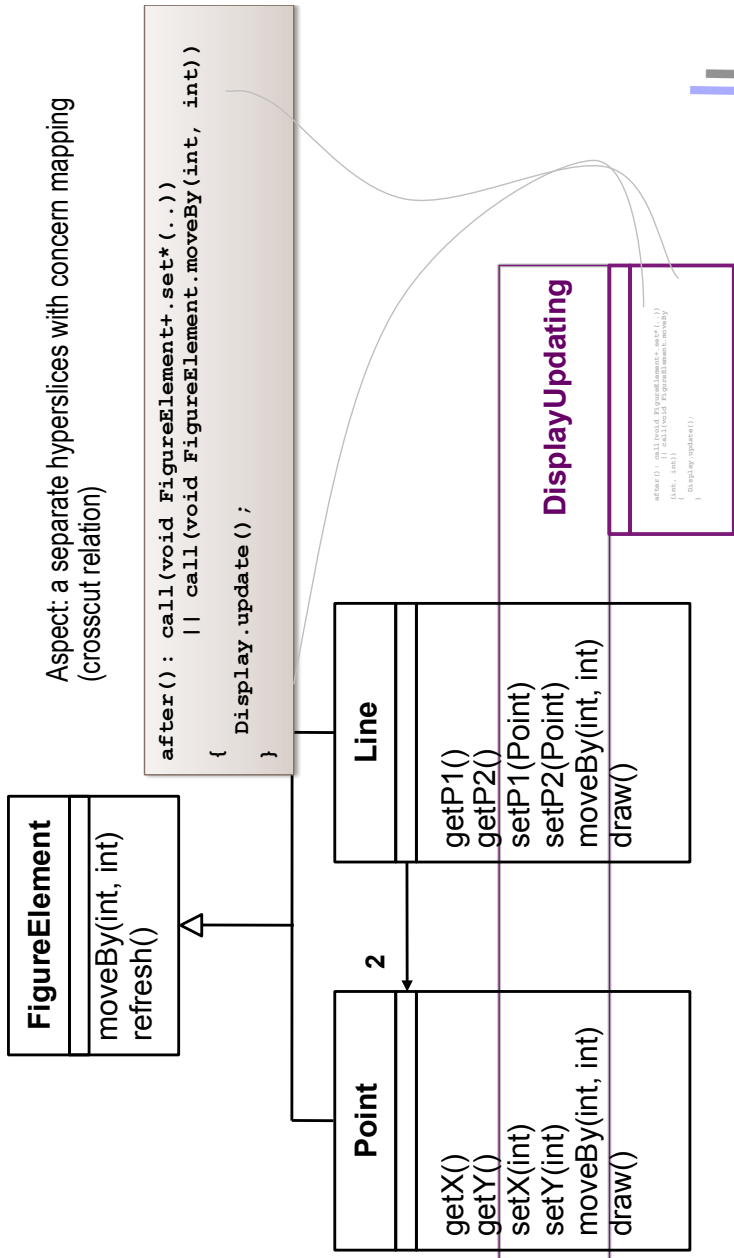
Some scatterings cannot easily be refactored.

Example:
All implementations of these methods end with call to:

Display.update () ;

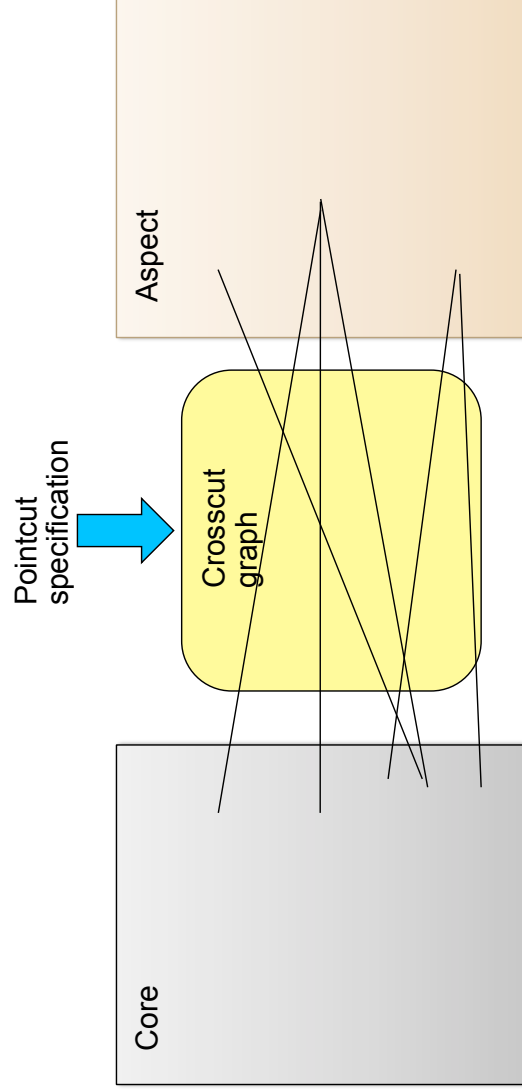


Needs AOP for a Solution



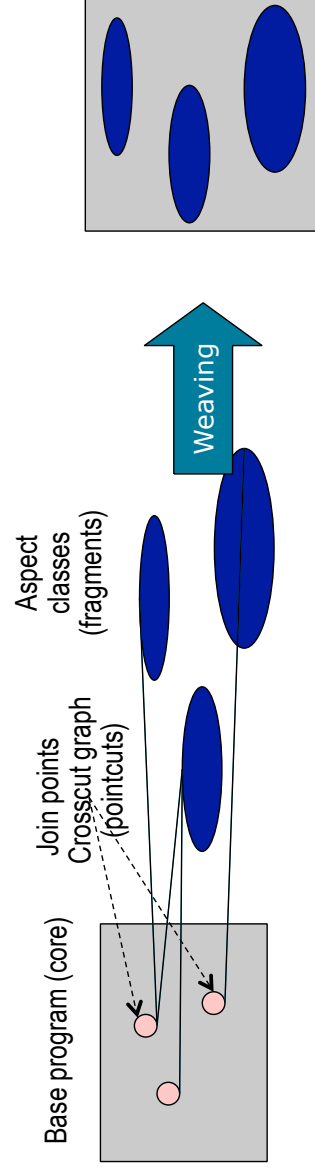
Crosscut Graphs

- Crosscuts are represented by crosscut graphs between core and aspect
- Pointcut specifications specify crosscut graphs



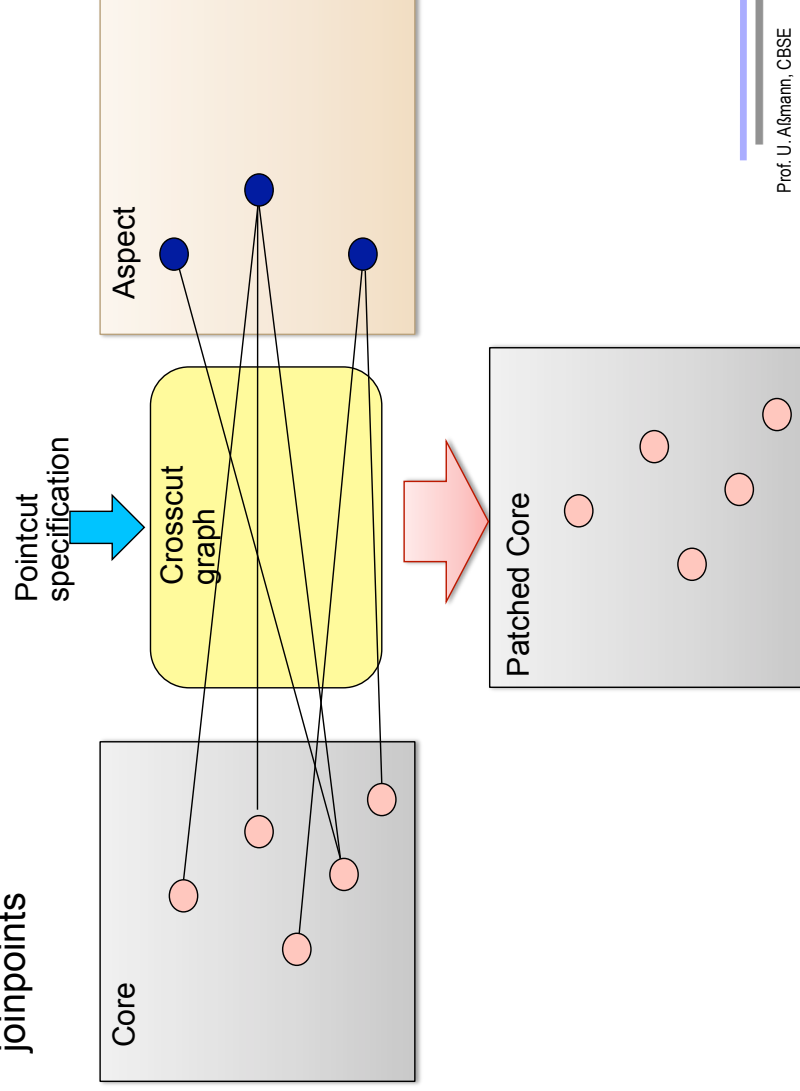
The AOP Idea (2)

- ▶ **Aspects** are separate, independent hyperslices, in which a **crosscutting concern mapping** relates fragment groups (advices) to concerns
- ▶ **Weaving** describes the composition, extending a core program at join points
 - ▶ At software development time, aspects and classes are kept as two, separate dimensions.
 - ▶ At run-time, both dimension need to be combined in some way for obtaining the final product.
- ▶ Weaving is **non-symmetric composition** (hyperslice composition is symmetric)



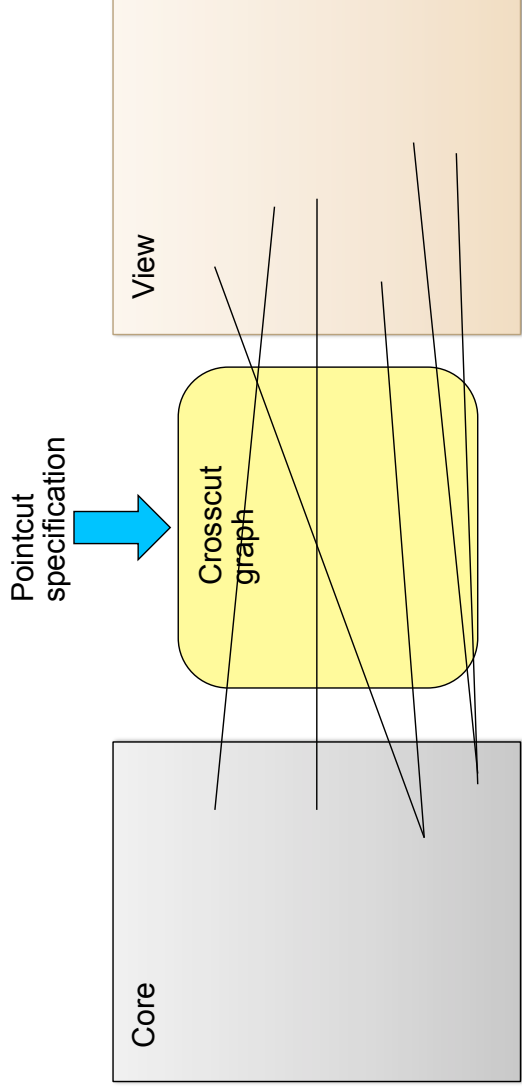
Aspects are Woven by Interpretation of the Crosscut Graphs

- Crosscut graphs are interpreted to insert advice fragments into core joinpoints



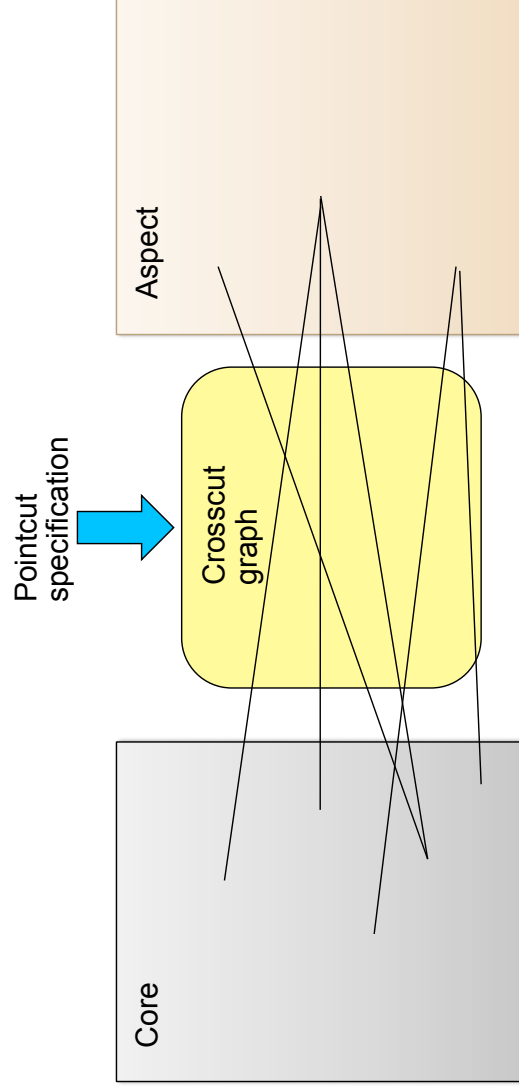
Crosscut Graphs in View-Based Programming

- Crosscut graphs are injective (View can extend only one open definition, but open definitions can be extended by many views)
- This solves *tangling*



Aspects allow for General Crosscut Graphs

- Crosscuts are non-injective, aspects can extend several open definitions
- This simulates *scattering*

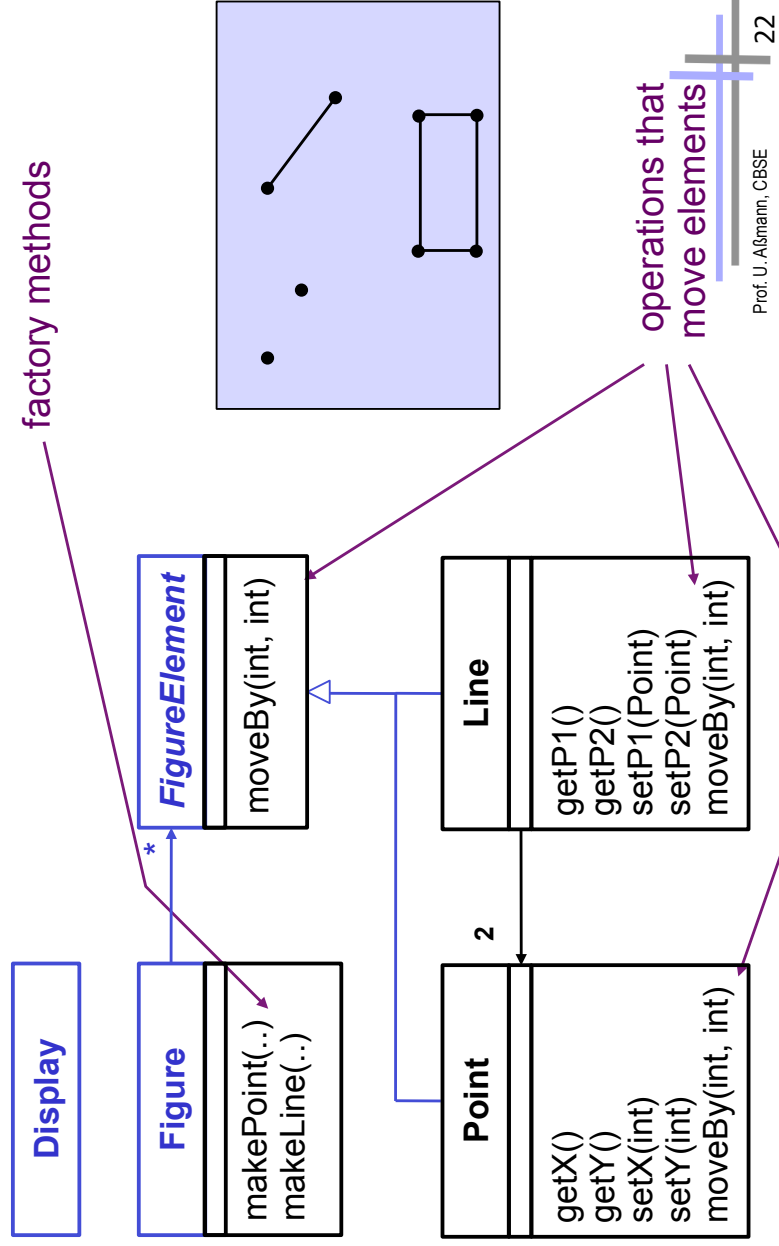


AspectJ: a Weaver for Java

- ▶ First production-quality AOP-technology
- ▶ Allows specifying hyperslices for crosscutting concerns as separate entities: Aspects
 - **Static join points** are code positions, hooks, open for extension
 - **Dynamic join points** are some points in the execution trace of an application, open for extension
 - **Pointcut**: a set of logically related join points
 - **Advice**: a fragment with behavior that should become active whenever a dynamic join point is encountered
 - **Weaving**: a technology for bringing aspects and base code together

```
// aspects are hyperslices plus integrated concern mapping
aspect <concern> {
  // introductions: fragments added to classes of the core
  // advices: fragments for extensions of methods
  // pointcuts: concern mapping from advices to
  //           joinpoints of the core
}
```

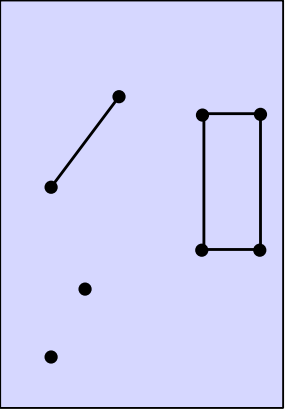
Example: A Simple Figure Editor



Example: A Simple Figure Editor (Java)

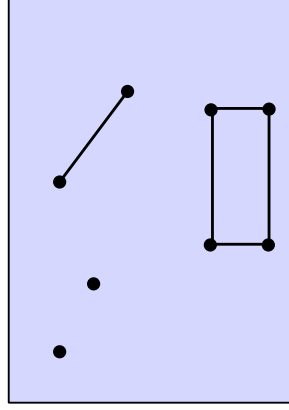
```
class Line implements FigureElement{
    private Point p1, p2;
    Point getP1() { return p1; }
    Point getP2() { return p2; }
    void setP1(Point p1) { this.p1 = p1; }
    void setP2(Point p2) { this.p2 = p2; }
    void moveBy(int dx, int dy) { ... }
}

class Point implements FigureElement {
    private int x = 0, y = 0;
    int getX() { return x; }
    int getY() { return y; }
    void setX(int x) { this.x = x; }
    void setY(int y) { this.y = y; }
    void moveBy(int dx, int dy) { ... }
}
```



Display Updating

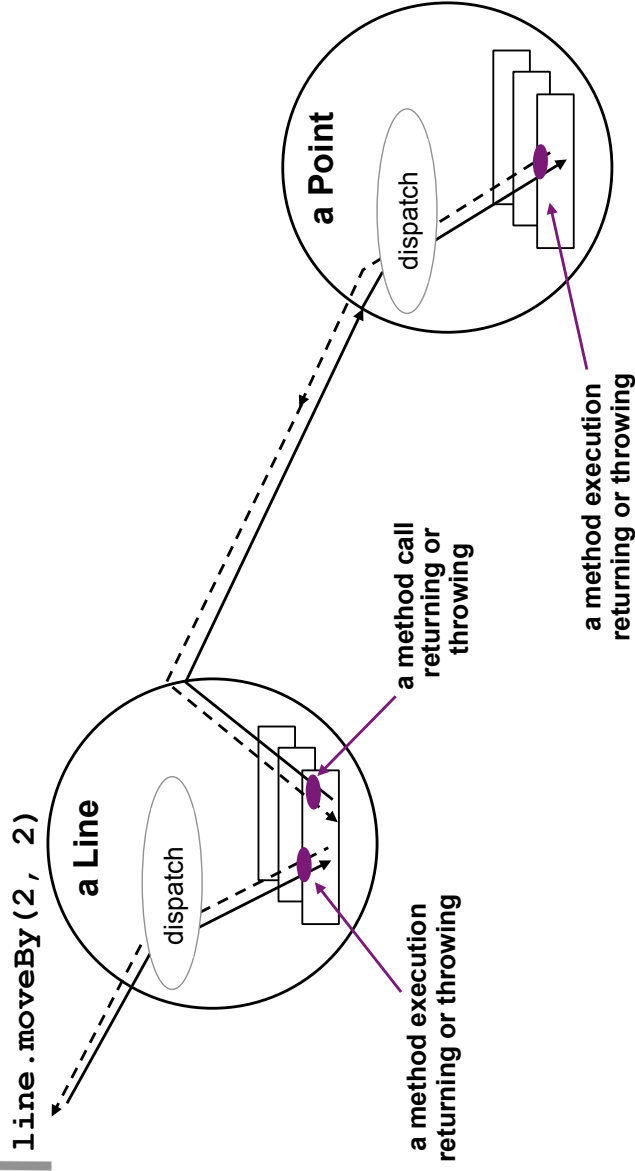
- ▶ Collection of figure elements
 - that move periodically
 - must refresh the display as needed



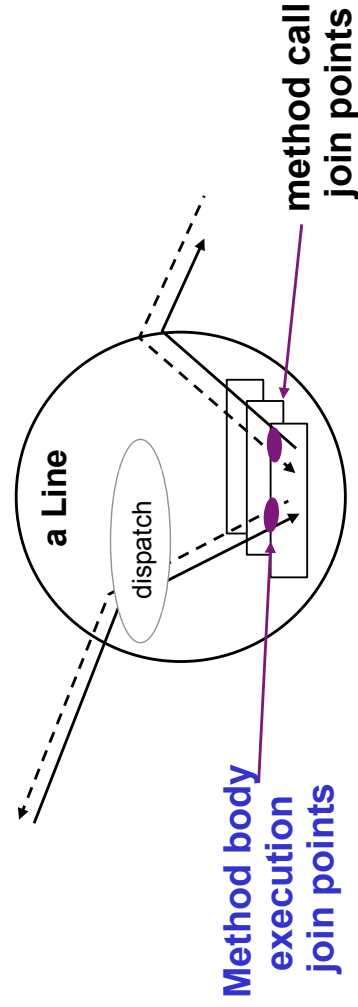
*we will assume just a
single display*

Aspect/J Dynamic Join Points (Dynamic Hooks)

- ▶ A **dynamic join point** is a point in the execution trace of a program, also in dynamic call graph



Dynamic Join Point Terminology

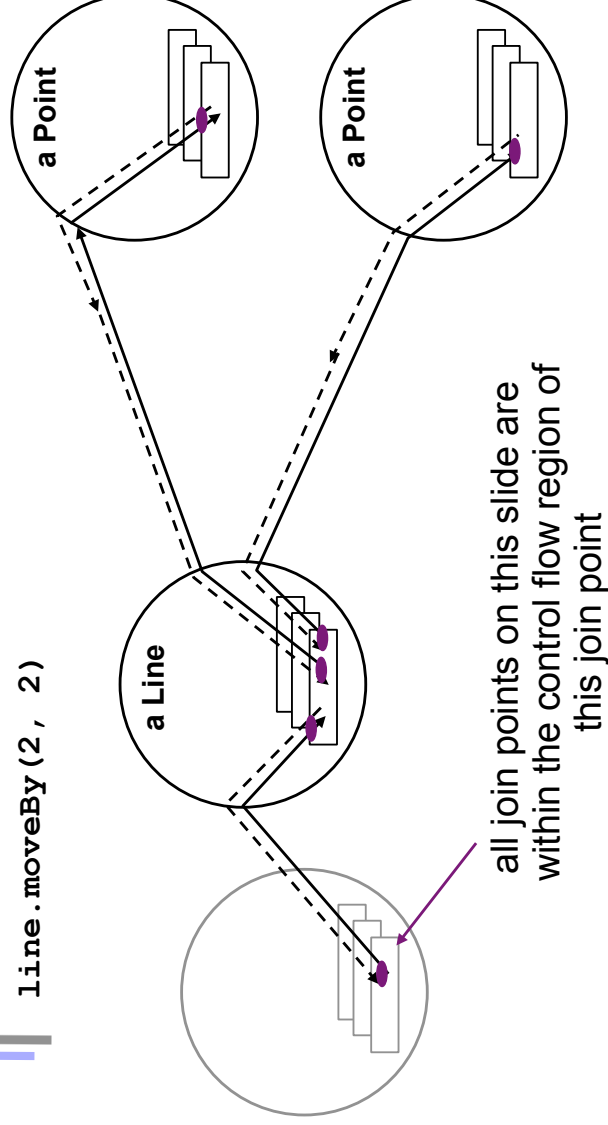


- ▶ The **join-point model** of Aspect/J defines several types of join points (join-point types)

- method & constructor call
- method & constructor execution
- field get & set
- exception handler execution
- static & dynamic initialization

Join Point Terminology

`line.moveBy(2, 2)`



Primitive Pointcuts

- ▶ A **pointcut** is an specification addressing a set of join points that:
 - can match or not match any given join point and
 - optionally, can pull out some of the values at that join point
 - “a means of identifying join points”
- ▶ Example: `call(void Line.setP1(Point))`

matches if the join point is a method call with this signature



Pointcut Composition

- ▶ Pointcuts are logical expressions in Aspect/J, they compose like predicates, using &&, || and !

a “void Line.setP1(Point)” call

```
call(void Line.setP1(Point)) ||  
call(void Line.setP2(Point));
```

Or

a “void Line.setP2(Point)” call

- whenever a Line receives a “void setP1(Point)” or “void setP2(Point)” method call



User-Defined Pointcuts

- ▶ User-defined (named) pointcuts can be used in the same way as primitive pointcuts

name parameters

```
pointcut move() :
```

```
call(void Line.setP1(Point)) ||  
call(void Line.setP2(Point));
```

*more on parameters
and how pointcut can
expose values at join
points in a few slides*



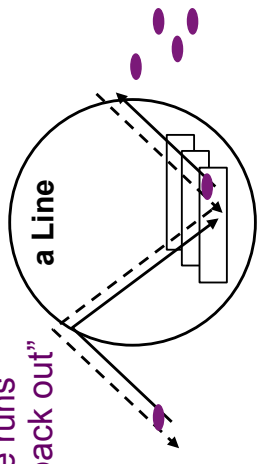


After Advice

- ▶ An **after advice** is a fragment describing the action to take after computation under join points

after advice runs

“on the way back out”



```
pointcut move() :
    call(void Line.setP1(Point)) ||
    call(void Line.setP2(Point));

after() returning: move() {
    <code here runs after each move>
}
```



A Simple Aspect

- An **aspect** defines a special class that can crosscut other classes
 - with one or several advices (fragments plus composition expression)
 - With at least one pointcut expressing the crosscut graph

```
aspect DisplayUpdating {
    pointcut move() :
        call(void Line.setP1(Point)) ||
        call(void Line.setP2(Point));
    after() returning: move() {
        Display.update();
    }
}
```



With and Without AspectJ

- Display.update calls are tangled through the code
- “what is going on” is less explicit

```
class Line {
    private Point p1, p2;
    Point getP1() { return p1; }
    Point getP2() { return p2; }

    void setP1(Point p1) {
        this.p1 = p1;
        // join point
    }
    void setP2(Point p2) {
        this.p2 = p2;
        // join point
    }
}
```

Weaver

```
aspect DisplayUpdating {
    pointcut move():
        call(void Line.setP1(Point)) ||
        call(void Line.setP2(Point));
    after() returning: move() {
        Display.update();
    }
}
```

```
class Line {
    private Point p1, p2;
    Point getP1() { return p1; }
    Point getP2() { return p2; }

    void setP1(Point p1) {
        this.p1 = p1;
        Display.update();
    }
    void setP2(Point p2) {
        this.p2 = p2;
        Display.update();
    }
}
```

A multi-class aspect

With pointcuts cutting across multiple classes

```
aspect DisplayUpdating {
    pointcut move():
        call(void FigureElement.moveBy(int, int)) ||
        call(void Line.setP1(Point))
        call(void Line.setP2(Point))
        call(void Point.setX(int))
        call(void Point.setY(int));
    after() returning: move() {
        Display.update();
    }
}
```

Using values at join points

- ▶ A pointcut can explicitly expose certain run-time values in parameters
- ▶ An advice can use the exposed value

```
pointcut move(FigureElement figElt):  
    target(figElt) &&  
    (call(void FigureElement.moveBy(int, int)) ||  
     call(void Line.setP1(Point)) ||  
     call(void Line.setP2(Point)) ||  
     call(void Point.setX(int)) ||  
     call(void Point.setY(int)));  
  
after(FigureElement fe) returning: move(fe) {  
    <fe is bound to the figure element>  
}
```

Pointcut Parameter defined and used

Pointcut parameter

Parameters of user-defined pointcut designator

- ▶ Variable is bound by user-defined pointcut declaration
 - Pointcut supplies value for variable
 - Value is available to all users of user-defined pointcut

```
pointcut move(Line l):  
    target(l) &&  
    (call(void Line.setP1(Point)) ||  
     call(void Line.setP2(Point)));  
    typed variable in place of type name
```

pointcut parameters

```
after(Line line): move(line) {  
    <line is bound to the line>  
}
```

Parameters of advice

- ▶ Variable is bound by advice declaration
 - Pointcut supplies value for variable
 - Value is available in advice body

```
pointcut move(Line l):  
    target(l) &&  
    (call(void Line.setP1(Point)) ||  
     call(void Line.setP2(Point)));  
  
advice parameters  
after(Line line): move(line) {  
    <line is bound to the line>  
}
```

Pointcut parameter



Explaining parameters...

- ▶ Value is 'pulled'
 - right to left across ':' left side : right side
 - from pointcuts to user-defined pointcuts
 - from pointcuts to advice, and then advice body

```
pointcut move(Line l):  
    target(l) &&  
    (call(void Line.setP1(Point)) ||  
     call(void Line.setP2(Point)));
```

```
after(Line line): move(line) {  
    <line is bound to the line>  
}
```



Join Point Qualifier “Target”

A **join point qualifier** does two things:

- exposes information from the context of the join point (e.g, target object of a message)
- tests a predicate on join points (e.g., a dynamic type test - any join point at which target object is an instance of type name)

```
target(<type name> | <formal reference>)
```

```
target(Point)
```

```
target(Line)
```

```
target(FigureElement)
```

“any join point” means it matches join points of all kinds:

method & constructor call join points

method & constructor execution join points

field get & set join points

exception handler execution join points

static & dynamic initialization join points

Getting target object in a polymorphic pointcut

```
target(<supertype name>) &&
```

- ▶ does not further restrict the join points
- ▶ does pick up the target object

```
pointcut move(FigureElement figElt):
```

```
target(figElt) &&
```

```
(call(void Line.setP1(Point)) ||
```

```
call(void Line.setP2(Point)) ||
```

```
call(void Point.setX(int)) ||
```

```
call(void Point.setY(int)));
```

```
after(FigureElement fe): move(fe) {
```

```
<fe is bound to the figure element>
```

```
}
```

Context & multiple classes

```
aspect DisplayUpdating {  
    pointcut move(WebElement figElt):  
        target(figElt) &&  
        (call(void WebElement.moveBy(int, int)) ||  
         call(void Line.setP1(Point)) ||  
         call(void Line.setP2(Point)) ||  
         call(void Point.setX(int)) ||  
         call(void Point.setY(int)));  
  
    after(WebElement fe): move(fe) {  
        Display.update(fe);  
    }  
}
```

Without AspectJ

```
class Line {  
    private Point p1, p2;  
  
    Point getP1() { return p1; }  
    Point getP2() { return p2; }  
  
    void setP1(Point p1) {  
        this.p1 = p1;  
        Display.update(this);  
    }  
    void setP2(Point p2) {  
        this.p2 = p2;  
        Display.update(this);  
    }  
}  
  
class Point {  
    private int x = 0, y = 0;  
  
    int getX() { return x; }  
    int getY() { return y; }  
  
    void setX(int x) {  
        this.x = x;  
        Display.update(this);  
    }  
    void setY(int y) {  
        this.y = y;  
        Display.update(this);  
    }  
}
```

► no locus of “display updating”

- evolution is cumbersome
- changes in all classes
- have to track & change all callers

With AspectJ

```
class Line {
    private Point p1, p2;
    Point getP1() { return p1; }
    Point getP2() { return p2; }
    void setP1(Point p1) {
        this.p1 = p1;
    }
    void setP2(Point p2) {
        this.p2 = p2;
    }
}

class Point {
    private int x = 0, y = 0;
    int getX() { return x; }
    int getY() { return y; }
    void setX(int x) {
        this.x = x;
    }
    void setY(int y) {
        this.y = y;
    }
}
```

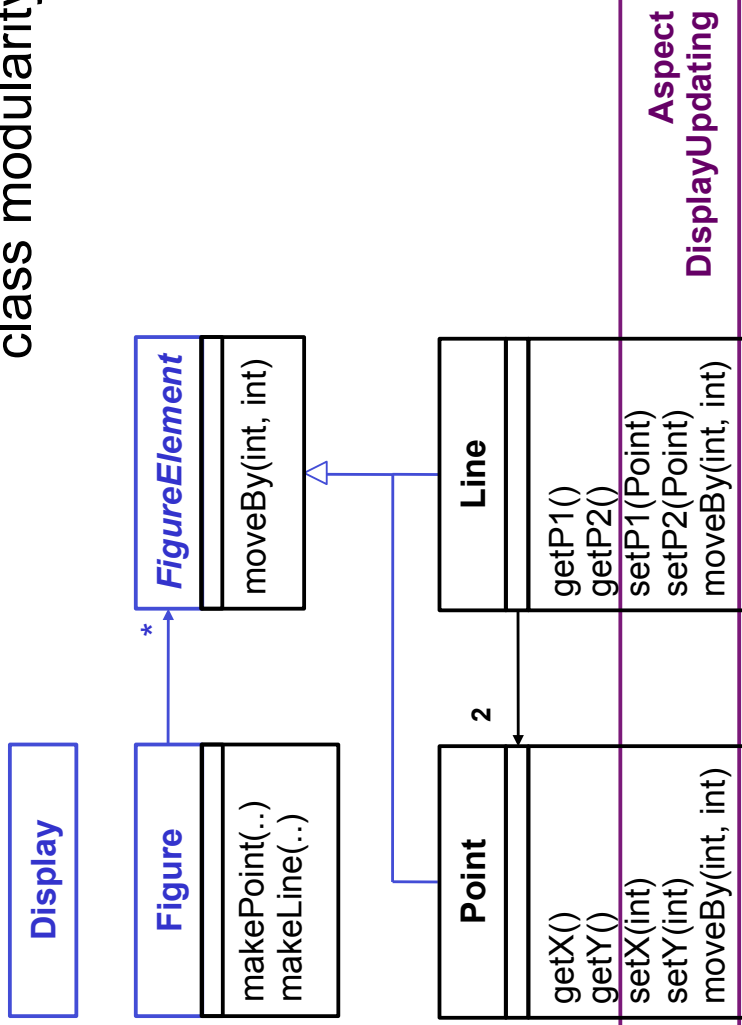
```
aspect DisplayUpdating {
    pointcut move(FigureElement figElt) :
        target(figElt) &&
        (call(void FigureElement.moveBy(int, int) ||
         call(void Line.setP1(Point))
         call(void Line.setP2(Point))
         call(void Point.setX(int))
         call(void Point.setY(int)));
    after(FigureElement fe) returning: move(fe) {
        Display.update(fe);
    }
}
```

- ▶ clear display updating module
 - all changes in single aspect
 - evolution is modular

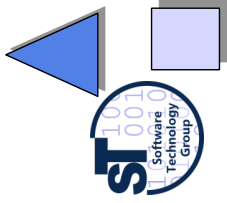


Aspects Crosscut Classes

aspect modularity cuts across
class modularity



24.3 Composition Operators and Point-Cuts



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Types of Advice Composition Operators



- ▶ before
 - ▶ after returning
 - ▶ after throwing
 - ▶ after
 - ▶ around
- ▶ before proceeding at join point
 - ▶ a value to join point
 - ▶ a throwable to join point
 - ▶ returning to join point either way
 - ▶ on arrival at join point gets explicit control over when and if program proceeds



Example: Contract checking with Aspects

- ▶ Simple application of before/after/around composition operators
- ▶ pre-conditions (assumptions)
 - check whether parameter is valid
- ▶ post-conditions (guarantees)
 - check whether values were set
- ▶ Invariants
 - ▶ Check conditions that should be true everywhere
- ▶ condition enforcement
 - force parameters to be valid and consistent



Pre-Condition (Assumption)

using before advice

```
aspect PointBoundsPreCondition {  
    before(int newX):  
        call(void Point.setX(int)) && args(newX) {  
            assert(newX >= MIN_X);  
            assert(newX <= MAX_X);  
        }  
    before(int newY):  
        call(void Point.setY(int)) && args(newY) {  
            assert(newY >= MIN_Y);  
            assert(newY <= MAX_Y);  
        }  
    private void assert(boolean v) {  
        if (!v)  
            throw new RuntimeException();  
    }  
}
```

what follows the ':' is always a pointcut – primitive or user-defined





Post-condition

using after advice

```
aspect PointBoundsPostCondition {  
    after(Point p, int newX) returning:  
        call(void Point.setX(int)) && target(p) && args(newX) {  
        assert(p.getX() == newX);  
    }  
  
    after(Point p, int newY) returning:  
        call(void Point.setY(int)) && target(p) && args(newY) {  
        assert(p.getY() == newY);  
    }  
  
    private void assert(boolean v) {  
        if ( !v )  
            throw new RuntimeException();  
    }  
}
```



Condition enforcement

using around advice

```
aspect PointBoundsEnforcement {  
    void around(int newX) {  
        call(void Point.setX(int)) && args(newX) {  
            proceed{// before the join point  
                clip(newX, MIN_X, MAX_X)  
            };  
            // after the join point  
            System.out.println("after");  
        }  
        void around(int newY):  
            call(void Point.setY(int)) && args(newY) {  
            proceed(clip(newY, MIN_Y, MAX_Y));  
        }  
        private int clip(int val, int min, int max) {  
            return Math.max(min, Math.min(max, val));  
        }  
    }  
}
```





Special Methods (Hooks in Advices)

- ▶ For each around advice with the signature
`<Tr> around(T1 arg1, T2 arg2, ...)`
- ▶ there is a special method with the signature
`<Tr> proceed(T1, T2, ...)`
- ▶ available only in around advice, meaning “*run what would have run if this around advice had not been defined*”



Property-based crosscutting (“Listener Aspects”)

```
package  
com.xerox.print;  
public class C1 {  
    ...  
    public void foo()  
        A.doSomething(...)  
    ...  
}
```

```
package  
com.xerox.scan;  
public class C2 {  
    ...  
    public int frotz()  
        A.doSomething(...)  
    ...  
    public int bar()  
        A.doSomething(...)  
    ...  
}
```

```
package  
com.xerox.copy;  
public class C3 {  
    ...  
    public String s1() {  
        A.doSomething(...);  
    }  
    ...  
}
```

- ▶ crosscuts of methods with a common property
 - public/private, return a certain value, in a particular package
- ▶ logging, debugging, profiling
 - log on entry to every public method





Property-based crosscutting

```
aspect PublicErrorLogging {  
    Log log = new Log();  
    pointcut publicInterface():  
        call(public * com.xerox.*.*(..));  
    after() throwing (Error e): publicInterface() {  
        log.write(e);  
    }  
}
```

neatly captures public interface of mypackage

- ▶ consider code maintenance
- ▶ another programmer adds a public method
 - i.e. extends public interface – this code will still work
- ▶ another programmer reads this code
 - “what’s really going on” is explicit



Wildcarding in pointcuts

```
target(Point)  
target(graphics.geom.Point)  
target(graphics.geom.*)  
target(graphics..*)
```

“*” is wild card
“..” is multi-part wild card

any type in graphics.geom
any type in any sub-package of graphics

```
call(void Point.setX(int))  
call(public * Point.*(..)) any public method on Point  
call(public * *(..)) any public method on any type
```

```
call(void Point.getX())  
call(void Point.getY())  
call(void Point.get*())  
call(void get*()) any getter
```

```
call(Point.new(int, int))  
call(new(..)) any constructor
```





Other Primitive Pointcuts

`this (<type name>)`

any join point at which currently executing object is an instance of type name

`within (<type name>)`

any join point at which currently executing code is contained within type name

`withincode (<method/constructor signature>)`

any join point at which currently executing code is specified method or constructor

`get(int Point.x)`

`set(int Point.x)`

field reference or assignment join points



Other Primitive Pointcuts

`execution (void Point.setX(int))`

method/constructor execution join points (actual running method)

`initialization (Point)`

object initialization join points

`staticinitialization (Point)`

class initialization join points (as the class is loaded)

`cflow (pointcut designator)`

all join points within the dynamic control flow of any join point in pointcut designator

`cflowbelow (pointcut designator)`

all join points within the dynamic control flow below any join point in pointcut designator, excluding thisJoinPoint



Example: Only top-level moves

```
aspect DisplayUpdating {  
  
    pointcut move (FigureElement fe):  
        target(fe) &&  
        (call(void FigureElement.moveBy(int, int)) ||  
         call(void Line.setP1(Point)) ||  
         call(void Line.setP2(Point)) ||  
         call(void Point.setX(int)) ||  
         call(void Point.setY(int)));  
  
    pointcut topLevelMove (FigureElement fe):  
        move(fe) && !cflowbelow(move (FigureElement));  
  
    after (FigureElement fe) returning: topLevelMove (fe) {  
        Display.update (fe);  
    }  
}
```

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Aspect/J Introductions

- An aspect can introduce new attributes and methods to existing classes

```
aspect PointObserving {  
    private Vector Point.observers = new Vector();  
    public static void addObserver (Point p, Screen s) {  
        p.observers.add(s); }  
  
    public static void removeObserver (Point p, Screen s) {  
        p.observers.remove(s); }  
  
    pointcut changes (Point p): target(p) && call(void Point.set*(int));  
  
    after (Point p): changes(p) {  
        Iterator iter = p.observers.iterator();  
        while ( iter.hasNext() ) {  
            updateObserver(p, (Screen)iter.next()); }  
    }  
    static void updateObserver (Point p, Screen s) {  
        s.display(p); }  
}
```

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Other approaches (1)

- ▶ <http://www.aosd.net/>
- ▶ AspectJ uses compile-time bytecode weaving,
 - but also inserts code that matches dynamic join points (dynamic weaving)
 - supports weaving aspects to existing *.class files (based on BCEL)
- ▶ AspectJ was taken over by IBM as part of the Eclipse project:
<http://www.eclipse.org/aspectj>

AspectC++ is an aspect-oriented extension to the C++ programming language.

AspectJ is a seamless aspect-oriented extension to Java that enables the modular implementation of a wide range of crosscutting concerns.

AspectWerkz is a dynamic, lightweight and high-performant AOP/AOSD framework for Java.

JAC is a Java framework for aspect-oriented distributed programming.

JBoss-AOP is the Java AOP architecture used for the JBoss application server.

Nanning is an Aspect Oriented Framework for Java based on dynamic proxies and aspects implemented as ordinary Java-classes.



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Other approaches (2)

AspectR is aspect-oriented programming for Ruby that allows you to wrap code around existing methods in your classes.

Aspects is an early prototype that enables aspect-oriented programming in the Squeak/Smalltalk environment.

CaesarJ is an aspect-oriented programming language that focuses on multi-view decomposition and aspect reusability.

DemeterJ and DJ facilitate the structure-shy encapsulation of traversal-related behavioral concerns.

JASCo is an aspect-oriented programming language tailored for component based software development.

JMangler is a framework for load-time transformation of Java programs, which supports conflict-free composition of independently developed aspects (implemented as JMangler transformer components) and their joint application to existing base classes.

MixJuice is an extension to Java, based on the difference-based module mechanism.

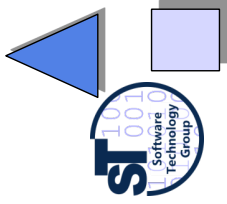
...



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24.4 AOSD



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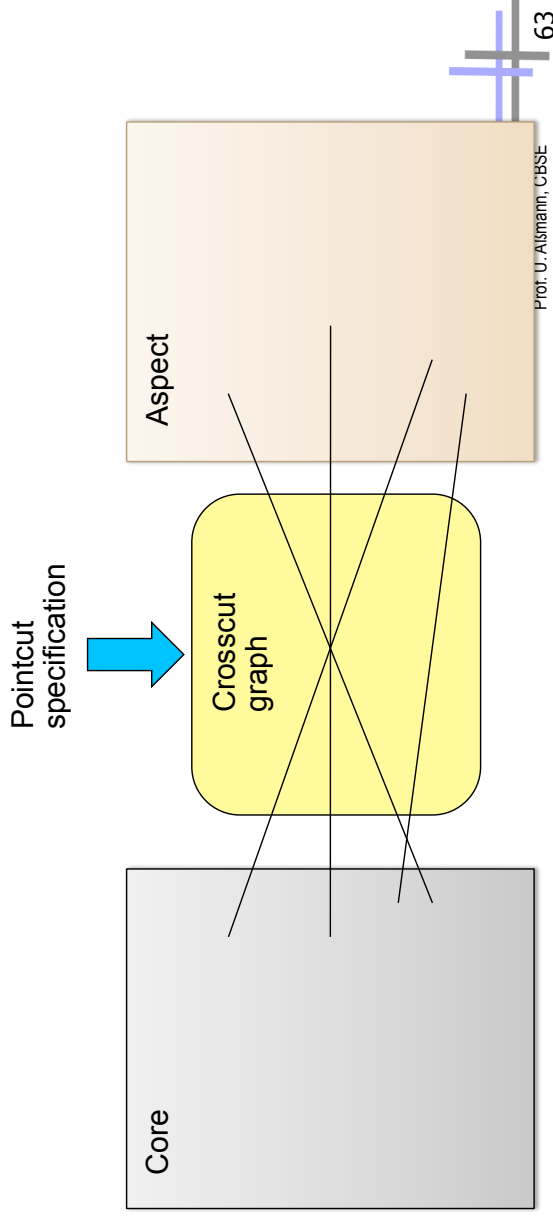


Problem of AOSD: Weaver Proliferation

- Who builds all these weavers, pointcut specification languages, extension engines, and template expanders?
- Answer:
 - Universal pointcut languages
 - Universal composable add-ons

Universal Pointcut Languages

- The specification of a pointcut is a graph-theoretic problem, and does not rely on the core nor aspect language
- Weaver proliferation can be avoided by *universal pointcut languages* for specifying crosscut graphs that *interconnect* base and aspect in any language



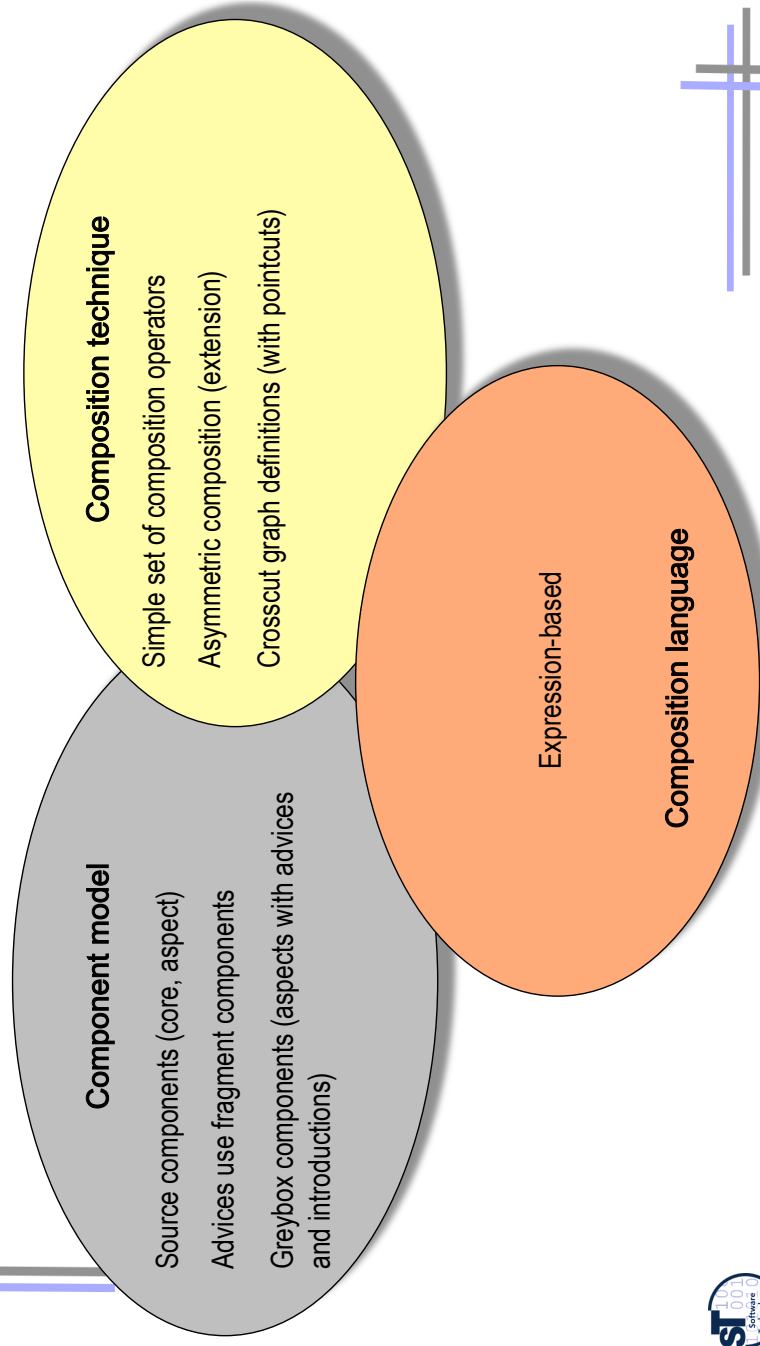
Universal Pointcut Languages

- A pointcut language connects *names* of the core and the aspect
 - does not know more concepts
- It can be used universally
- Example:
 - Xpath, can it be used as pointcut language?
 - Can you separate pointcuts from Aspect/J advices and address advice joinpoints?
 - Relational algebra, SQL, Datalog
 - Graph rewriting
 - Logic

Towards Aspect-Oriented System Development (AOSD)

- ▶ Aspects are important in the whole lifecycle
 - requirements (*early aspects*)
 - analysis
 - design (*model aspects*)
 - implementation (*code aspects*)
 - test
- ▶ **Aspect-aware development** uses crosscut graphs and their specification languages for all languages (modeling and programming)
- ▶ [Johannes] shows how to make crosscut graphs for arbitrary languages
- ▶ **Aspect-aware tools** interpret crosscut graphs
- ▶ Reuseware is a metaweaver, a generator for weavers

24.5 Evaluation: Aspects as Composition System





The End

- ▶ Many slides courtesy to Wim Vanderperren, Vrije Universitet Brussel, and the Aspect/J team

