

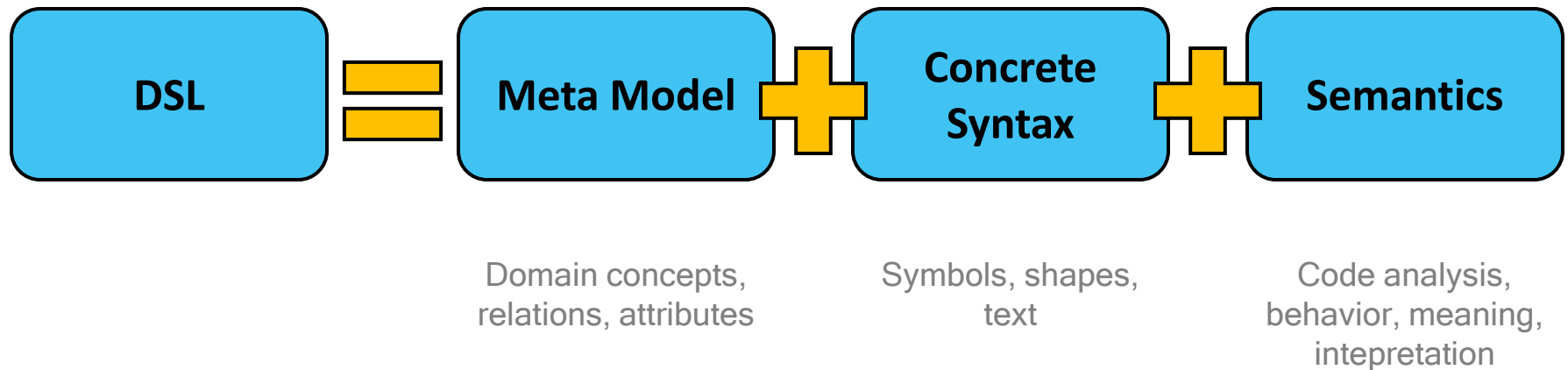


JastEMF: Reference Attribute Grammars for EMF-based DSLs

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ACSE 17.12.2012

What's in a DSL?



Benefits of Metamodelling

Metamodelling is a standardisation process with the following benefits:

- MM 1 Metamodelling Abstraction
- MM 2 Metamodelling Consistency
- MM 3 Metamodel Implementation Generators
- MM 4 Metamodel/Model Compatibility
- MM 5 Tooling Compatibility

However, metamodelling lacks convenient mechanisms for semantics specification.

Benefits of Attribute Grammars (AGs)

AGs are very convenient to specify semantics for tree structure with the following benefits:

- AG 1: Declarative Semantics Abstraction
- AG 2: Semantics Consistency
- AG 3: Semantics Generators
- AG 4: Semantics Modularity

Claim: A combination of MM and AGs enables *semantics integrated metamodelling* and leads to more successful and reliable tool implementations.

Attribute Grammars

Formalism to compute static semantics over syntax trees [Knuth68]

- Basis: context-free grammars + attributes + semantic functions
- Evaluation by tree visitors with different visiting strategies
 - Static: ordered attribute grammars (OAGs)
 - Dynamic: demand-driven evaluation
- AGs are modular and extensible

Improvements

- Higher order attribute grammars (HOAGs) [Vogt+89]
- Reference attributed grammars (RAGs) [Hedin00,Boyland05]

Formal Definition

(Short) Definition (attribute grammar): An attribute grammar (AG) is an 8-tuple $G=(\mathbf{G}_0, \mathbf{Syn}, \mathbf{Inh}, \mathbf{Syn}_x, \mathbf{Inh}_x, \mathbf{K}, \mathbf{\Omega}, \mathbf{\Phi})$ with the following components

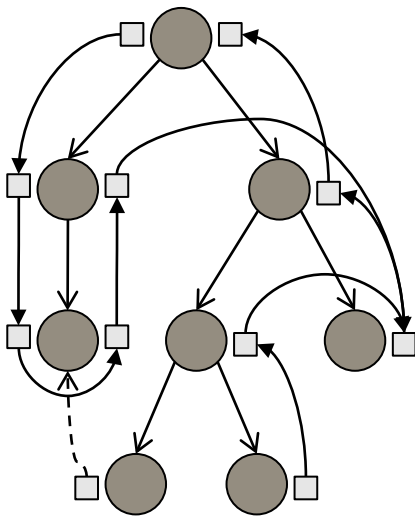
- $\mathbf{G}_0 = (N, \Sigma, P, S)$ a CFG,
- \mathbf{Syn} and \mathbf{Inh} the finite, disjoint sets of synthesized and inherited attributes,
- $\mathbf{Syn}_x : N \rightarrow P(\mathbf{Syn})$ a function that assigns a set of synthesized attributes to each nonterminal in G_0 ,
- $\mathbf{Inh}_x : N \rightarrow P(\mathbf{Inh})$ a function that assigns a set of inherited attributes to each nonterminal in G_0 ,
- \mathbf{K} a set of attribute types/sorts,
- $\mathbf{\Omega} : \mathbf{Inh} \cup \mathbf{Syn} \rightarrow \mathbf{K}$ a function assigning each attribute a $\kappa \in \mathbf{K}$,
- $\mathbf{\Phi}$ a set of semantic functions $\varphi_{(p,i,a)}$ with $p \in P$, $i \in \{0, \dots, n_p\}$, $a \in \mathbf{Syn}_x(p_i) \cup \mathbf{Inh}_x(p_i)$.

Formal Definition

Kind of AG depends on K and Φ

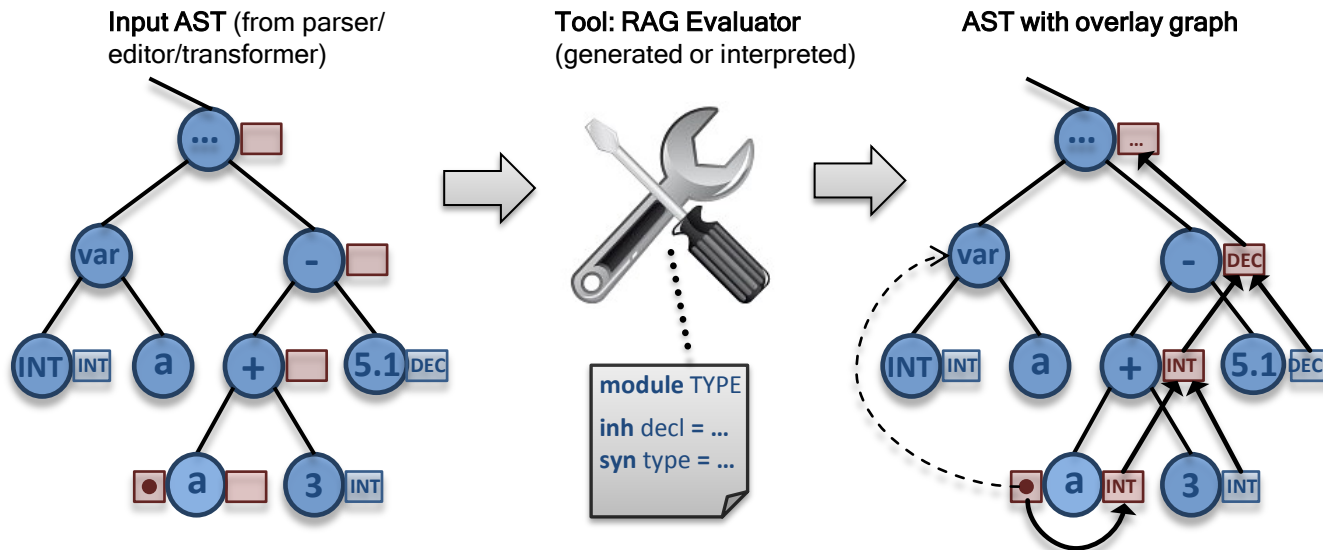
- ❑ Standard AG
 - $K \cap N = \{\}$
- ❑ Higher Order AG (HOAG)
 - $K \cap N \neq \{\}$
 - all $\varphi \in \Phi$ only yield copied or fresh values
- ❑ Reference Attribute Grammar (RAG)
 - $K \cap N \neq \{\}$
 - $\varphi \in \Phi$ yield references (pointers) to nonterminals (or value copies)

Kinds of Attributes



- **Inherited attributes** (inh): Top-down value dataflow/computation
- **Synthesized attributes** (syn): Bottom-up value dataflow/computation
- **Collection attributes** (coll): Collect values freely distributed over the AST
- **Reference attributes**: Compute references to existing nodes in the AST

Basic Working Principle of RAG Tools



How do (R)AGs Relate with Metamodeling?

Models are graphs!?

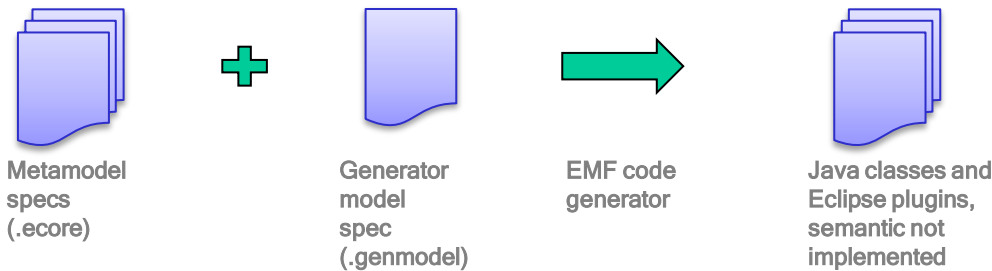
- Ecore (EMOF) metamodels are commonly build around a **tree**-based abstract syntax used by
 - **Tree** iterators, **tree** editors, transformation tools, textual editors
 - Syntax mappers like *EMFText* use the **tree** structure to compute all other information (e.g., resolving cross references)
 - Graphical editors use the **tree** structure to manage user created object hierarchies, cross references and values therein and to compute read-only information (e.g., cross references, derived values)
- ➔ Ecore (EMOF) models are ASTs with cross-references (overlay graphs) and derived information!

Mapping Between Both Worlds [Bürger+11]

Syntax in Ecore	Syntax in RAGs	} E _{Syn}
EClass	AST Node Type	
EReference[containment]	Non Terminal	
EAttribute[non-derived]	Terminal	
Semantics Interface in Ecore	Semantics in RAGs	} E _{Sem}
EAttribute[derived]	[synthesized inherited] attribute	
EAttribute[derived,multiple]	collection attribute	
EReference[non-containment]	collection attribute, reference attribute	
EOperation[side-effect free]	[synthesized inherited] attribute	

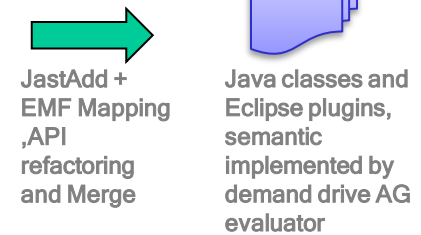
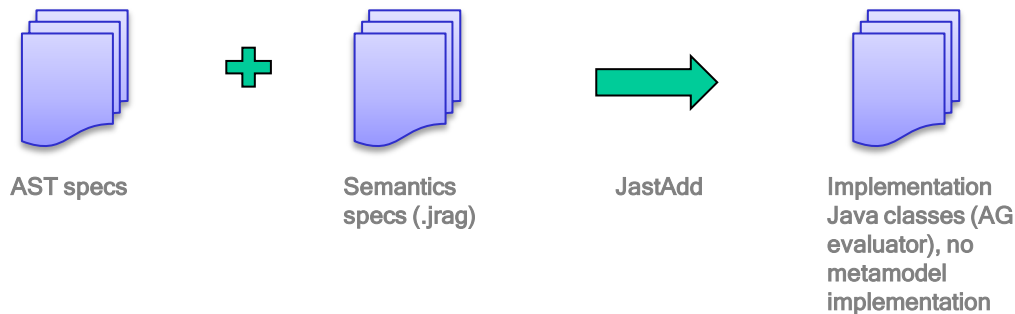
The JastEMF Tool

EMF basic generation process



JastEMF meta integration

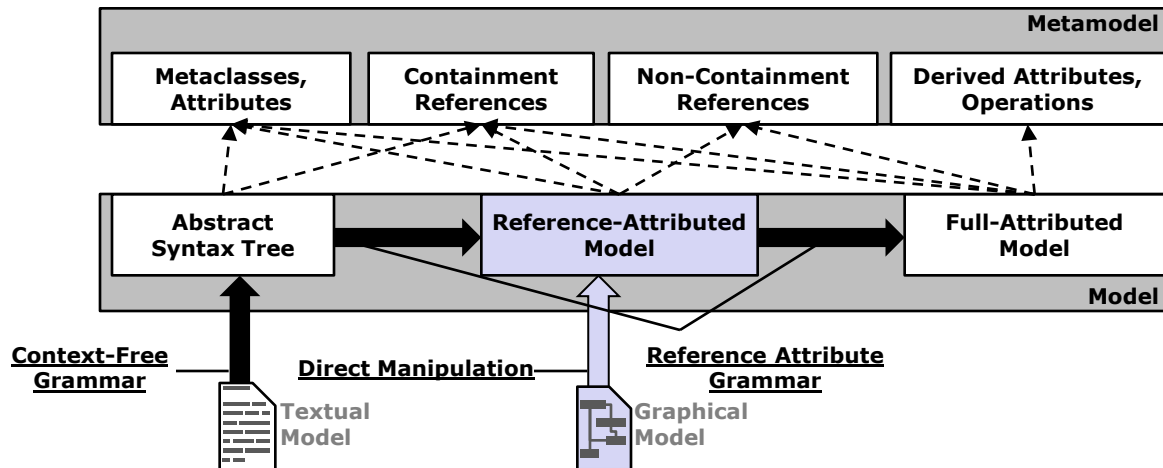
JastAdd basic generation process



The JastEMF Tool

Semantic evaluation can start from (partly) reference-attributed models

- Non-containment references can have predefined values (e.g., specified by users using a diagram editor)

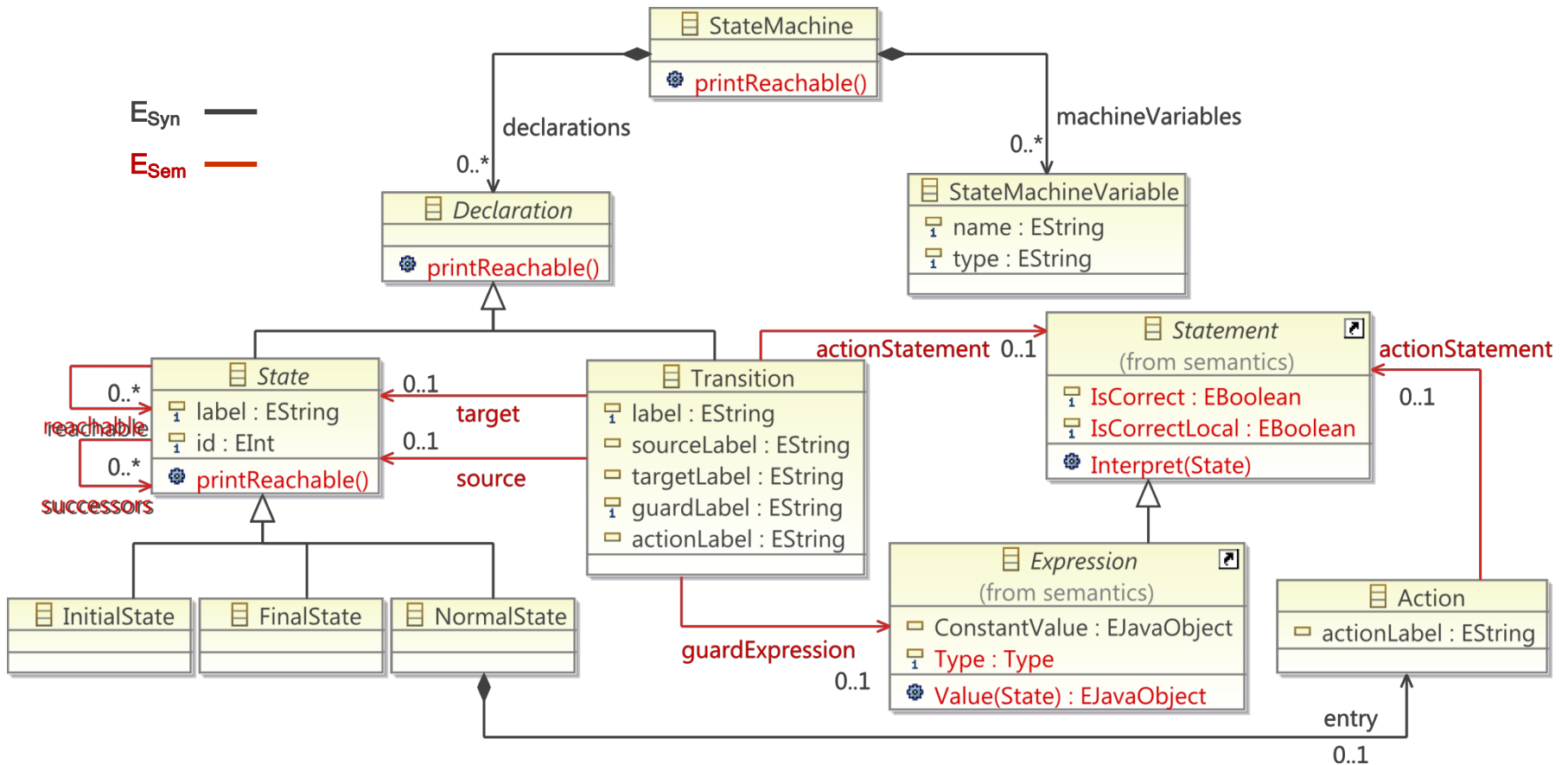


- If a value is given: Use it instead of attribute equation

Examples

1. Graphical DSL [Bürger+11]
 - Statecharts
 - Syntax in GMF
2. Textual DSL [Heidenreich+12]
 - Forms DSL (see ACSE tutorial of 04.12.12)
 - Syntax in EMFText
3. Imperative Programming Language [Bürger+11]
 - SiPLE (**S**imple **P**rogramming **L**anguage **E**xample)
 - Syntax in xText

Example 1: Statechart Metamodel



(Ecore-based, extended version of Statechart example in Hedin, G.: Generating Language Tools with JastAdd. In: GTTSE '09. LNCS, Springer (2010))

Example 1: Statechart Attributes

AST specification (partial):

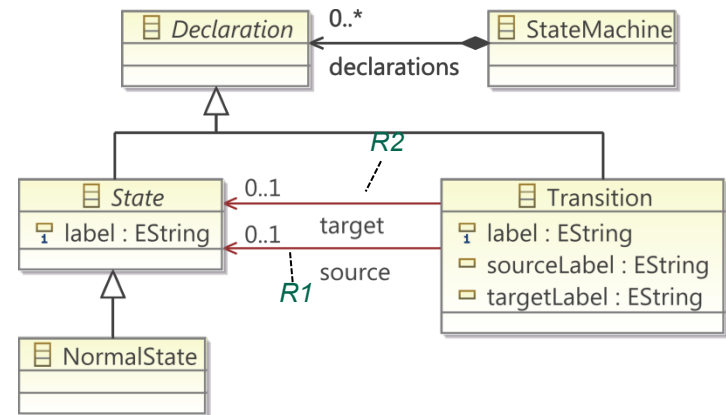
```

abstract State:Declaration ::= <label:String>;
NormalState:State;
Transition:Declaration ::=<label:String>
    <sourceLabel:String><targetLabel:String>;
    
```

Attribution example:

```

syn lazy State Transition.source() = lookup(getSourceLabel()); // R1
syn lazy State Transition.target() = lookup(getTargetLabel()); // R2
inh State Declaration.lookup(String label); // R3
eq StateMachine.getDeclarations(int i).lookup(String label) { ... } // R4
syn State Declaration.localLookup(String label) =
    (label==getLabel()) ? this : null; // R5
    
```



compute closure

reuse of
metamodels and semantics

compute transition ends from labels

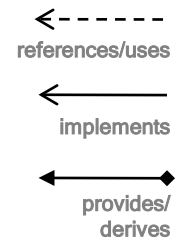
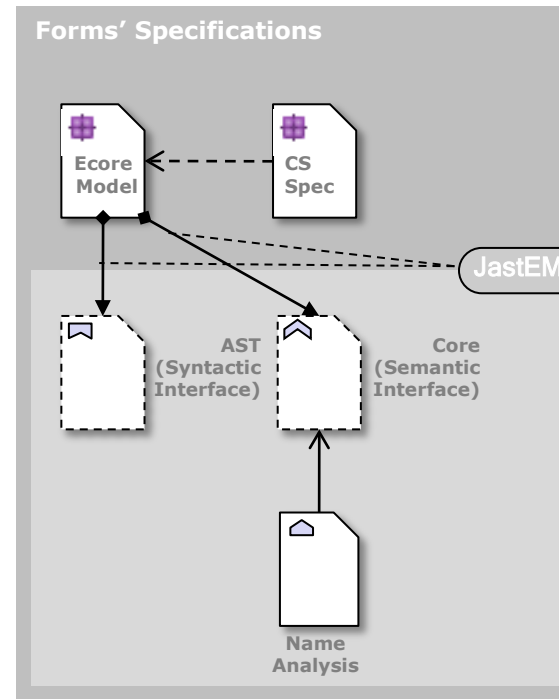
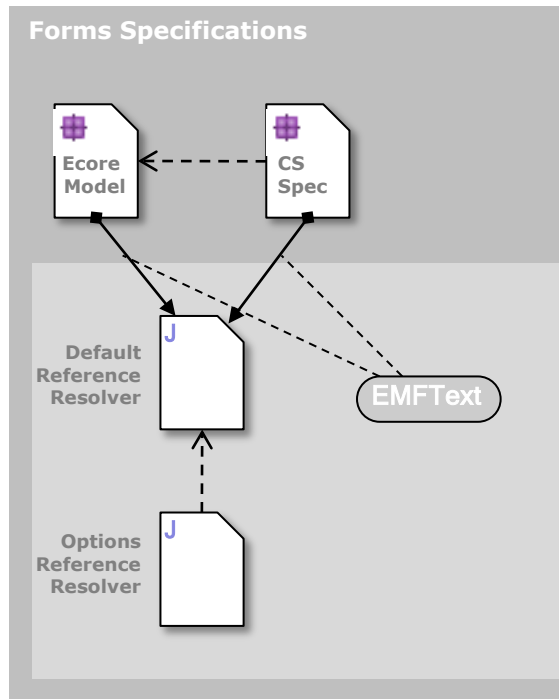
The screenshot shows the JastEMF IDE interface. The main window displays a state machine diagram with states Start, A, B, and End. Transitions are labeled with actions and guards: start [] counter := 0; back-to-A [counter < 11]; go-to-B [] counter := counter + 1; end [counter > 10]. A variable Integer:counter is shown. The Properties window is open for a NormalState, showing a table of properties and values.

Property	Value
Id	0
Label	A
Reachable	Normal State 1, Normal State 0, Final State 5
Successors	Normal State 1

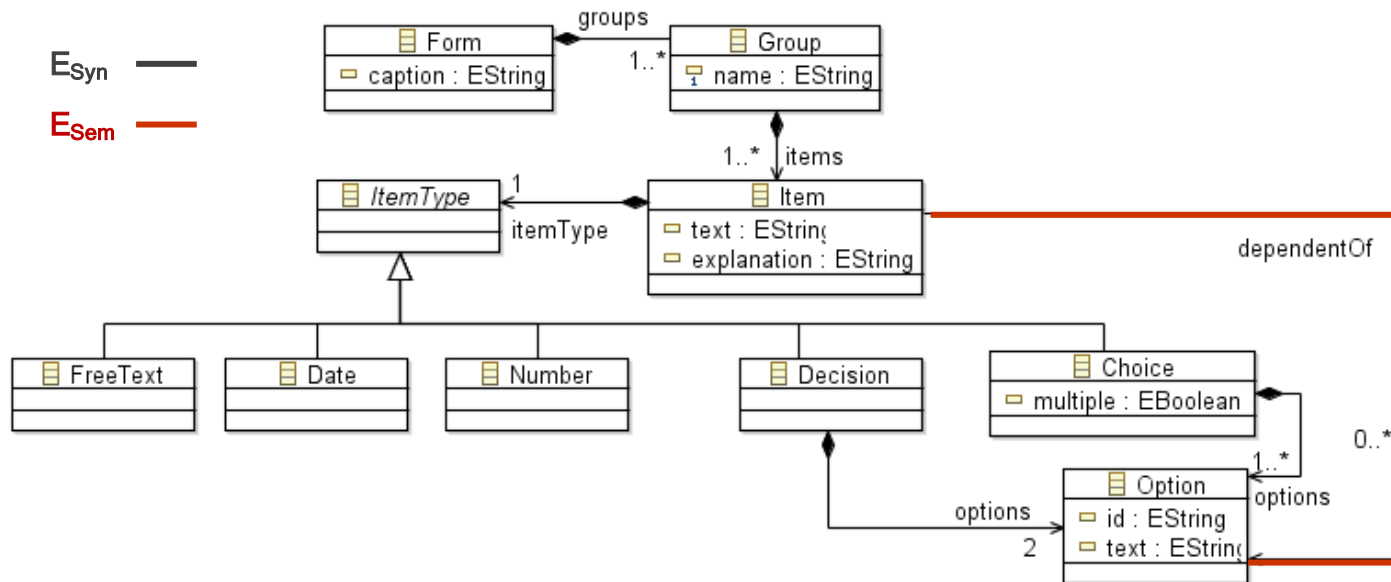
Red arrows indicate the following connections:

- From "compute transition ends from labels" to the transition label "go-to-B [] counter := counter + 1;" in the diagram.
- From "compute closure" to the state "A" in the diagram.
- From "reuse of metamodels and semantics" to the state "A" in the diagram and the "Id" property value "0" in the Properties window.

Example 2: Forms DSL



Example 2: Forms Metamodel



Example 2: Forms as AST Grammar

```
Form ::= <caption:String> groups:Group*;
```

Terminal - EAttribute

```
Item ::= <text:String> <explanation:String> <dependentOfName:String>
       itemType:ItemType;
```

ASTNode - EClass

```
abstract ItemType;
FreeText:ItemType;
Choice:ItemType ::= <multiple:boolean> options:Option*;
Date:ItemType;
Number:ItemType;
Decision:ItemType ::= options:Option*;
```

Nonterminal - EReference
(Containment)

Inheritance

```
Option ::= <id:String> <text:String>;
```

```
Group ::= <name:String> items:Item*;
```

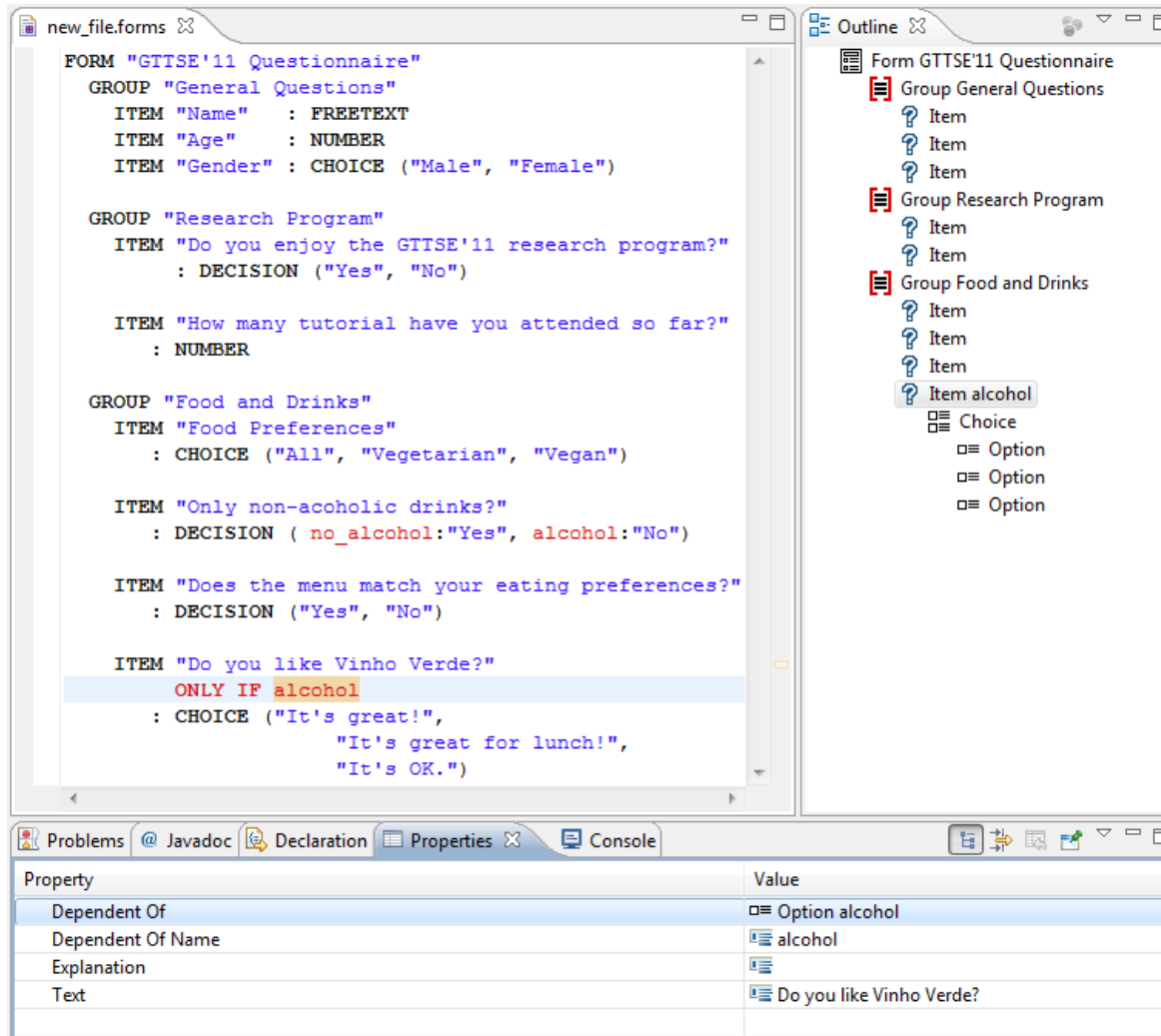
Example 2: Froms Attributes

```
aspect NameAnalysis {
    inh Form ASTNode.form();
    syn EList Item.dependentOf();
    inh EList ASTNode.LookupOption(String optionName);
    coll EList<Option> Form.Options() [new BasicEList()] with add;

    Option contributes this to Form.Options() for form();

    eq Form.getgroups(int index).form() = this;
    eq Item.dependentOf() = LookupOption(getdependentOfName());

    eq Form.getgroups(int index).LookupOption(String optionName){
        EList result = new BasicEList();
        for(Option option:Options()){
            if(optionName.equals(option.getid()))
                result.add(option);
        }
        return result;
    }
}
```



The screenshot shows the JastEMF IDE with a form definition file named 'new_file.forms'. The form is titled 'GTTSE'11 Questionnaire' and contains several groups of questions. The 'Food and Drinks' group includes a question about liking Vinho Verde, which is dependent on an 'alcohol' choice.

```

FORM "GTTSE'11 Questionnaire"
GROUP "General Questions"
  ITEM "Name" : FREETEXT
  ITEM "Age" : NUMBER
  ITEM "Gender" : CHOICE ("Male", "Female")

GROUP "Research Program"
  ITEM "Do you enjoy the GTTSE'11 research program?"
    : DECISION ("Yes", "No")

  ITEM "How many tutorial have you attended so far?"
    : NUMBER

GROUP "Food and Drinks"
  ITEM "Food Preferences"
    : CHOICE ("All", "Vegetarian", "Vegan")

  ITEM "Only non-achoholic drinks?"
    : DECISION ( no_alcohol:"Yes", alcohol:"No")

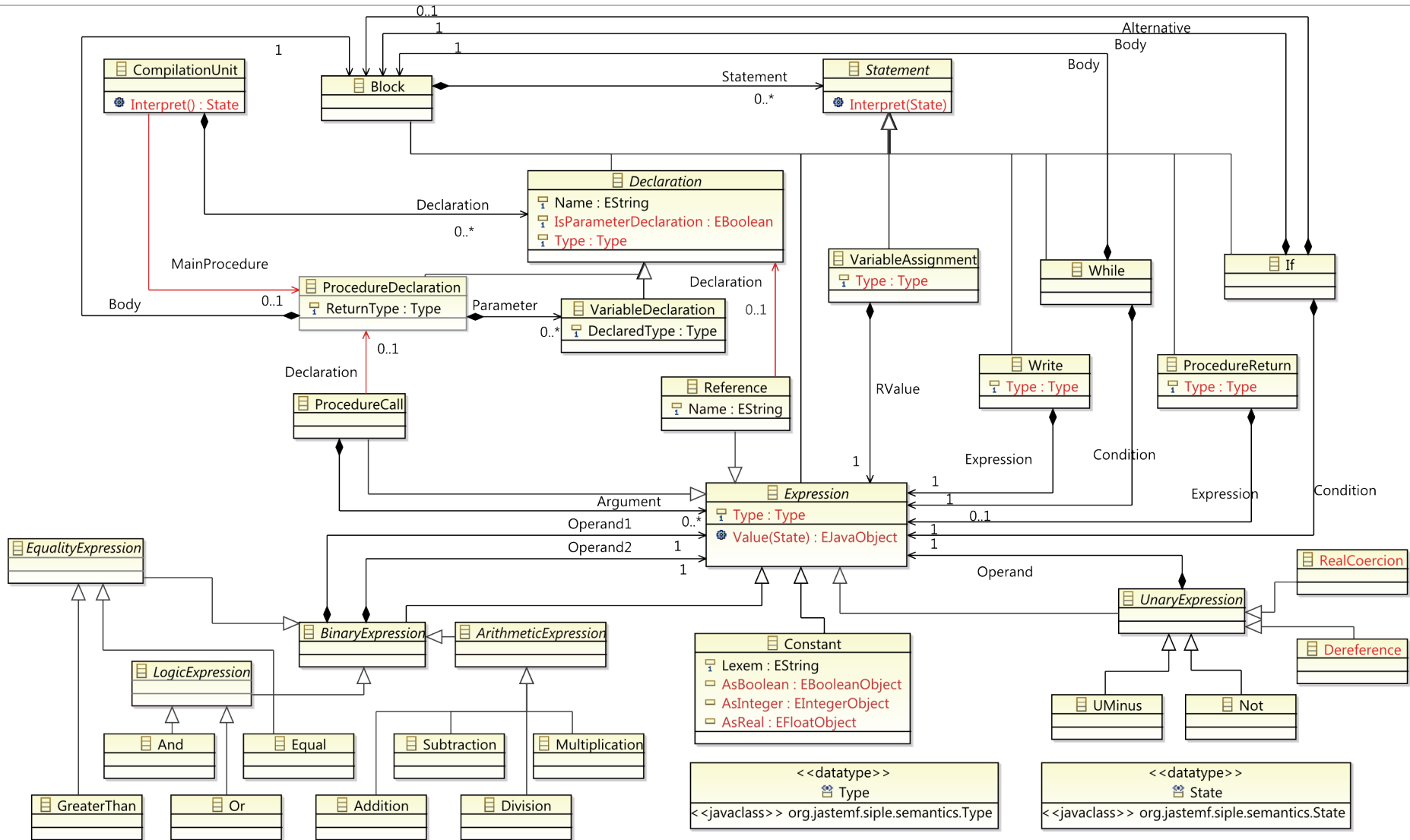
  ITEM "Does the menu match your eating preferences?"
    : DECISION ("Yes", "No")

  ITEM "Do you like Vinho Verde?"
    ONLY IF alcohol
    : CHOICE ("It's great!",
             "It's great for lunch!",
             "It's OK.")
  
```

The Outline view on the right shows the hierarchical structure of the form, including groups like 'General Questions', 'Research Program', and 'Food and Drinks', with their respective items and choices.

The Properties view at the bottom shows the following details for the selected item:

Property	Value
Dependent Of	☐ Option alcohol
Dependent Of Name	☒ alcohol
Explanation	☒
Text	☒ Do you like Vinho Verde?



Example 3: SiPLE Types (Excerpt from Semantic Interface)

```
aspect TypeAnalysis {
    syn Type Declaration.Type();
    syn Type VariableAssignment.Type();
    syn Type ProcedureReturn.Type();
    syn Type Write.Type();
    syn Type Read.Type();
    syn Type Expression.Type();
}

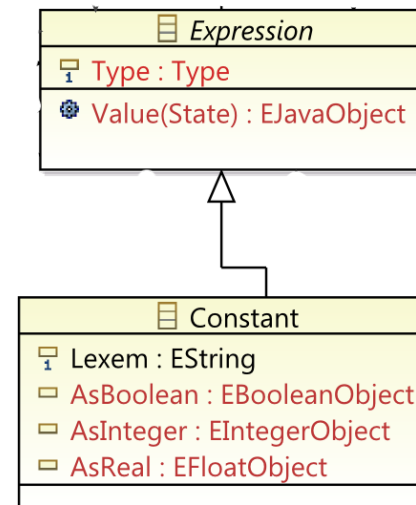
aspect NameAnalysis {
    // Ordinary name space:
    inh LinkedList<Declaration> ASTNode.LookUp(String name);
    syn ProcedureDeclaration CompilationUnit.MainProcedure();
    syn Declaration Reference.Declaration();
}
```


Example 3: SiPLE Types (Excerpt from Definitions)

```

/** Expressions' Type */
eq Constant.Type() {
    if (AsBoolean() != null)
        return Type.Boolean;
    if (AsReal() != null)
        return Type.Real;
    if (AsInteger() != null)
        return Type.Integer;
    return Type.ERROR_TYPE;
}

```



Example 3: SiPLE Types (Excerpt from Definitions)

```

eq LogicExpression.Type() =
    getOperand1().Type().isBoolean() &&
    getOperand2().Type().isBoolean() ?
        Type.Boolean : Type.ERROR_TYPE;
  
```

```

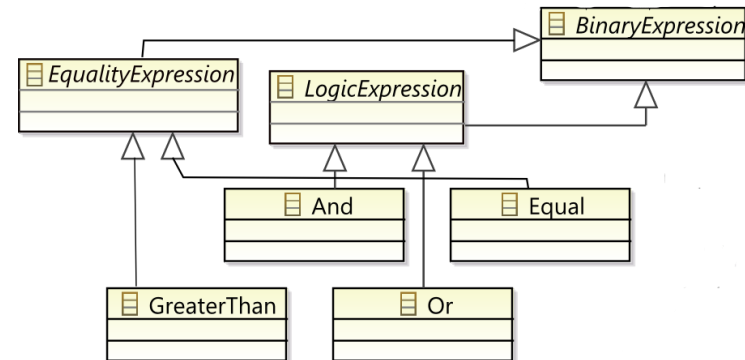
eq EqualityExpression.Type() =
    Type.bequals(getOperand1().Type(), getOperand2().Type()) &&
    (getOperand1().Type().isReal() ||
    getOperand1().Type().isInteger()) ?
        Type.Boolean : Type.ERROR_TYPE;
  
```

```

eq Equal.Type() =
    Type.bequals(getOperand1().Type(), getOperand2().Type()) ?
        Type.Boolean : Type.ERROR_TYPE;
  
```

```

eq ArithmeticExpression.Type() =
    Type.bequals(getOperand1().Type(), getOperand2().Type()) &&
    (getOperand1().Type().isReal() ||
    getOperand1().Type().isInteger()) ?
        getOperand1().Type() : Type.ERROR_TYPE;
  
```



Resource - siple_test/src/test.siple - Eclipse Platform

File Edit Navigate Search Project Run Window Help

Quick Access Resource

Project Explorer

- siple_test
 - src
 - test.siple
 - JRE System Library [JavaSt

test.siple

```

Procedure main() Begin
  x := 0;
  y := 0;
  z := 0;

  Var x:Integer;
  x := 100;

  Procedure writeMain() Begin
    Write x;
    Write y;
    Write z;
  End;

  Procedure l1() Begin
    Var y:Integer;
    y := x + 100;

    Procedure writeL1() Begin
      Write x;
      Write y;
      Write z;
    End;

    Procedure l2() Begin

```

Line: 29

Outline

- Compilation Unit
 - x
 - y
 - z
 - writeGlob
 - Block
 - incGlobVars
 - Block
 - Variable Assignment Integer
 - Variable Assignment Integer
 - Variable Assignment Integer
 - main
 - Block
 - Variable Assignment Integer
 - Variable Assignment Integer
 - Variable Assignment Integer
 - x
 - Variable Assignment Integer
 - writeMain
 - l1
 - Procedure Call Undefined

Tasks Task List Properties

Property	Value
As Boolean	
As Integer	100
As Real	
Lexem	100
Procedure In Context	Procedure Declaration main
Type	Integer

Writable Insert 28 : 9

Literature

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Thank you!