TECHNISCHE UNIVERSITÄT DRESDEN

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Reference Attribute Grammars with JastAdd

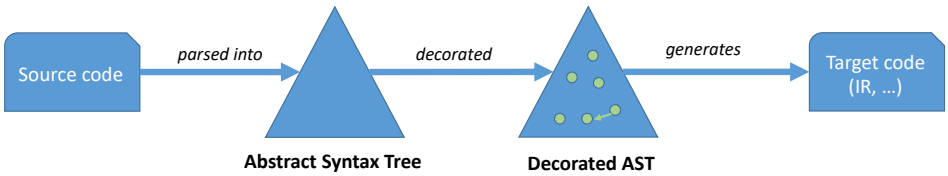
Dresden, January 14, 2021

Attribute Grammars (AGs)

Attributes

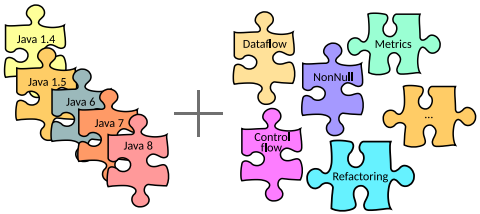
- compute *derived* properties of nodes in abstract syntax tree
- proposed by Donald Knuth in 1968
- *references* **(RAGs)** simplify navigation in attribute definitions

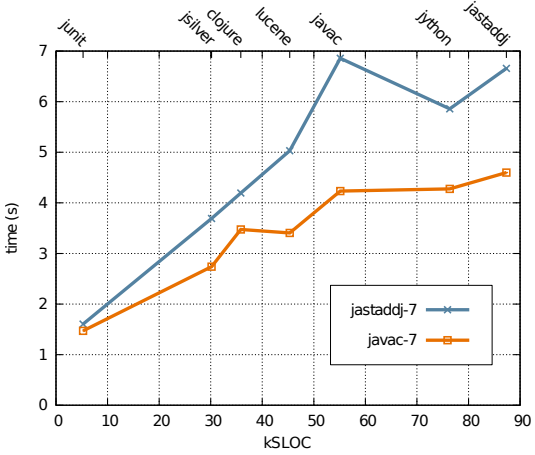
Today's Focus: JastAdd RAG system



JastAdd Applications

ExtendJ



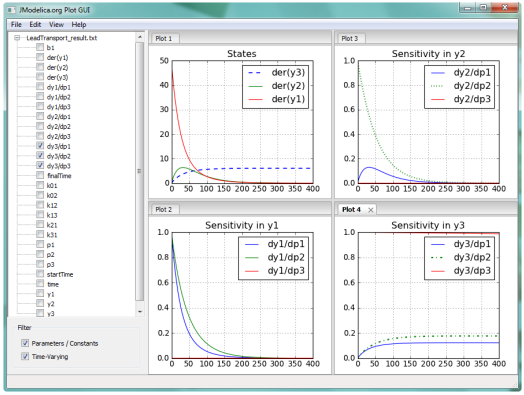


— Java 8 compiler with many extensions

— performance compared to OpenJDK

JastAdd Applications
JModelica

- open source modelica compiler
- maintained and used by company Modelon

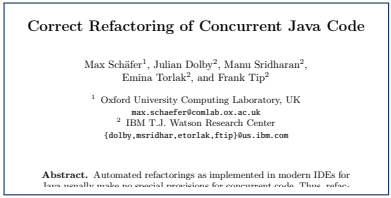


JModelica User Guide, jmodelica.org

Refactoring and JastAdd

Previous work by Max Schäfer

- papers at ECOOP, OOPSLA, ICSE, …
- dissertation
- JastAdd refactoring tool JRRT
  `https://code.google.com/archive/p/jrrt/`



Correct Refactoring of Concurrent Java Code

Max Schäfer[1], Julian Dolby[2], Manu Sridharan[2], Emina Torlak[2], and Frank Tip[2]

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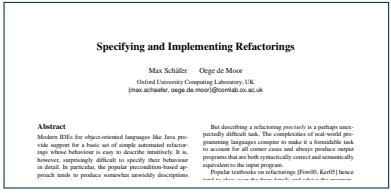
Abstract. Automated refactorings as implemented in modern IDEs for

[Schäfer et al., 2010]



Specifying and Implementing Refactorings

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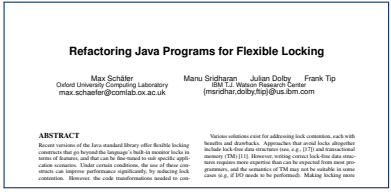
Abstract

Modern IDEs for object-oriented languages like Java provide support for a basic set of simple automated refactorings whose behaviour is easy to describe intuitively. It is, however, surprisingly difficult to specify their behaviour in detail. In particular, the popular precondition-based approach tends to produce somewhat unwieldy descriptions

But describing a refactoring *precisely* is a perhaps unexpectedly difficult task. The complexities of real-world programming languages conspire to make it a formidable task to account for all corner cases and always produce output programs that are both syntactically correct and semantically equivalent to the input program.

Popular textbooks on refactorings [Fow00, Ker05] hence

[Schaefer and de Moor, 2010]



Refactoring Java Programs for Flexible Locking

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ABSTRACT

Recent versions of the Java standard library offer flexible locking constructs that go beyond the language's built-in monitor locks in terms of features, and that can be fine-tuned to suit specific application scenarios. Under certain conditions, the use of these constructs can improve performance significantly, by reducing lock contention. However, the code transformations needed to con-

Various solutions exist for addressing lock contention, each with benefits and drawbacks. Approaches that avoid locks altogether include lock-free data structures (see, e.g., [17]) and transactional memory (TM) [13]. However, writing correct lock-free data structures requires more expertise than can be expected from most programmers, and the semantics of TM may not be suitable in some cases (e.g. if I/O needs to be performed). Making locking more

[Schafer et al., 2011]

JastAdd Exercise 1

Presentation

— a *practical* introduction to RAGs
— the JastAdd compiler and its infrastructure

Homework

— simple expression language
— construction of a small grammar
— writing some attributes

Reference Attribute Grammars

An introduction to

- grammar specification
- syntax trees
- attribute specification

Reference Attribute Grammars

An introduction to

- grammar specification
- syntax trees
- attribute specification

⚠ Disclaimer

- focus on the JastAdd understanding
- no parsing, only syntax trees

JastAdd Grammar

Elements:

Nonterminals	`A B SomeName`
Terminals/Tokens	`<X:int> <Y>` (default type `String`)

Production rules:

Child nodes	`A ::= C First:B Second:B;`
List/optional children	`B ::= C* [MyD:D];`
Terminals	`C ::= <TerminalSymbol:String>;`
Abstract nonterminals	`abstract E ::= <Name>;`
Inheritance	`F : E ::= <Value:int>;` `G : E ::= H <Value:float>;`
Empty productions	`H ::= /* right side can be empty! */;` `H; // also valid`

JastAdd Grammar

Generated interface for nonterminals:

regular nonterminal	`class A { /* */ }`
abstract nonterminal	`abstract class E { /* */ }`

Generated child accessors (within nonterminal class):

unnamed child	`public C getC() { /* */ }`
named child	`public C getMyChild() { /* */ }`
list children	`public C getC(int index) { /* */ }`
optional child	`public boolean hasMyD() { /* */ }`
terminal	`public String getTerminalSymbol() { /* */ }`

JastAdd Attributes

Attributes

- proposed by Donald Knuth [Knuth, 1968]
- computed properties of tree
- side-effect free
- *declaration* and *definition*
- different types with different information flow

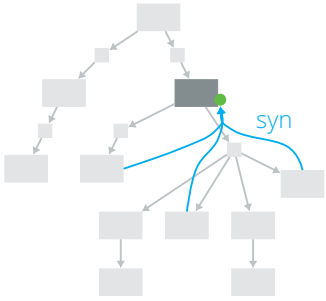
JastAdd Attributes
Synthesized Attributes

Information from the subtree: *synthesized* attributes

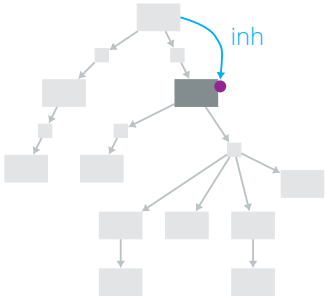
- Must be defined for declared type
- If type is abstract for all non-abstract sub-types
- Example:

```
syn boolean TrackElement.isSegment();

// attribute equations
eq Switch.isSegment()  { return false; }
eq Segment.isSegment() { return true;  }
```



JastAdd Attributes
Inherited Attributes



Information from parent: *inherited* attributes

- Must be defined on an ancestor
- Example:

```
inh Region Element.containingRegion();

// attribute equation
eq Region.getElement(int index).containingRegion() = this;
```

JastAdd Attributes
Inherited Attributes

Information from parent: *inherited* attributes

- Must be defined on an ancestor
- Example:

```
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// attribute equation
eq Region.getElement(int index).containingRegion() = this;
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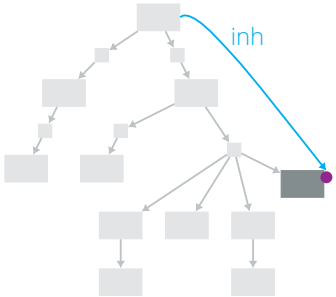
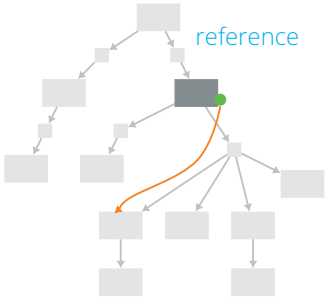


JastAdd Attributes
Reference Attributes



Existing nodes as attribute values

- can be any type of attribute (syn, inh, coll)
- Example:

```
inh Region Element.containingRegion();

// attribute equation
eq Region.getElement(int index).containingRegion() = this;
```

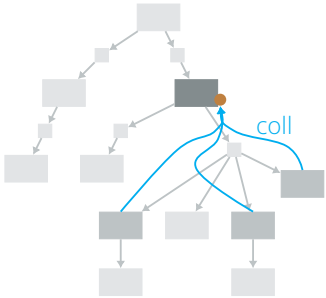
JastAdd Attributes
Collection Attributes

Collecting information: *collection* attributes

- Must be defined for declared type
- If type is abstract for all non-abstract sub-types
- Example:

```
coll Set<Element> Region.coolElements() [new HashSet];

// contribution to collection
Element contributes this
  when isCool()
  to Region.coolElements();
```



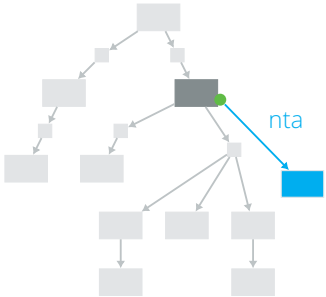
JastAdd Attributes
Nonterminal Attributes

Building new subtrees: *nonterminal* attributes

- Also: *higher order attribute*
- Subtrees must be **new** objects!
- Example:

```
// grammar excerpt
A ::= /* ... */;
B ::= <Name:String>;
// declaration
syn nta B A.getB();

// attribute equation
eq A.getB() {
  B b = new B();
  b.setName("Boaty McBoatface");
  return b;
}
```



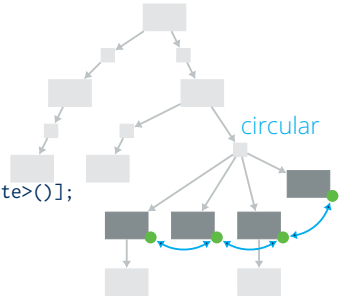
JastAdd Attributes
Circular Attributes

Fix-Point Computation: *circular* attributes

- can call itself
- computed iteratively
- example:

```
syn Set<State> State.reachable() circular [new HashSet<State>()];

eq State.reachable() {
  HashSet<State> result = new HashSet<State>();
  for (State s : successors()) {
    result.add(s);
    result.addAll(s.reachable());
  }
  return result;
}
```



Attributes in JastAdd

synthesized: information from subtree

inherited: information from parents

reference: any kind of attribute can be reference; points to other nonterminal

collection: information from nodes of certain type

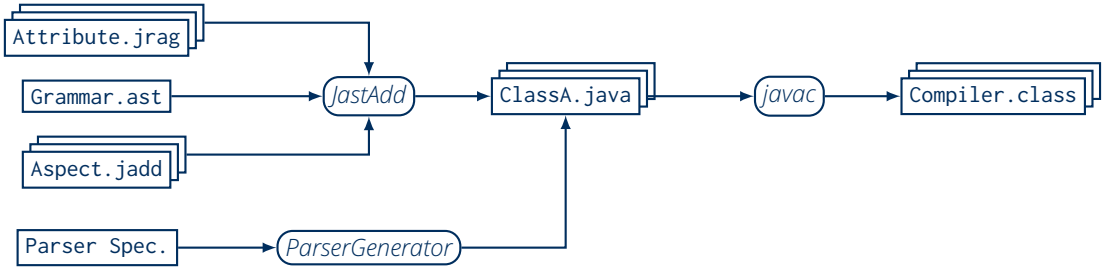
nonterminal: can be synthesized or inherited; compute new subtrees

circular: any kind of attribute can be circular; iterative fixpoint computation

The JastAdd System

RAG to Java Code Generation

- nonterminals → classes
- attributes → methods
- additional magic



Other JastAdd Features

Attribute Evaluation

- caching of attribute values
  - attribute values are memoized
  - configurable on per-equation level
- incremental attribute evaluation
  - dynamic attribute dependency graph

Aspect-Oriented Programming Features

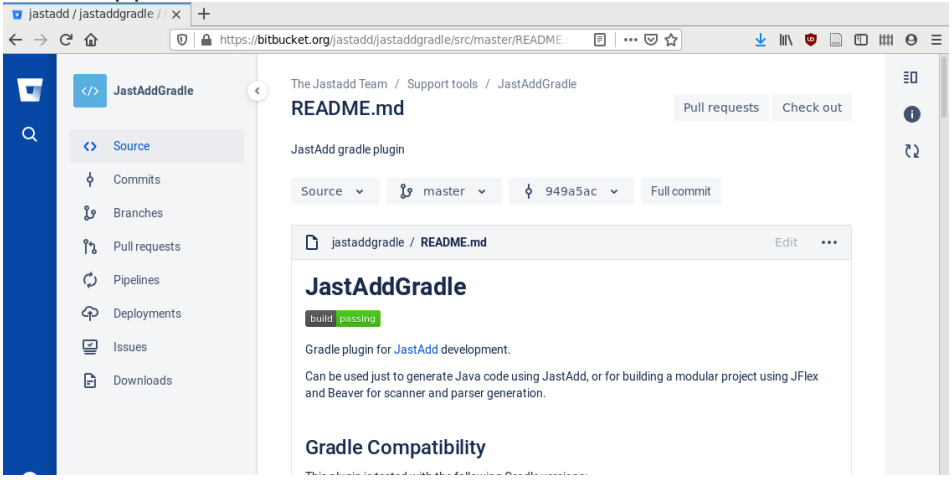
- additional methods can be woven into classes
- methods and attributes can be refined

Other nice features

- debugging and tracing support

Build Tools: JastAddGradle

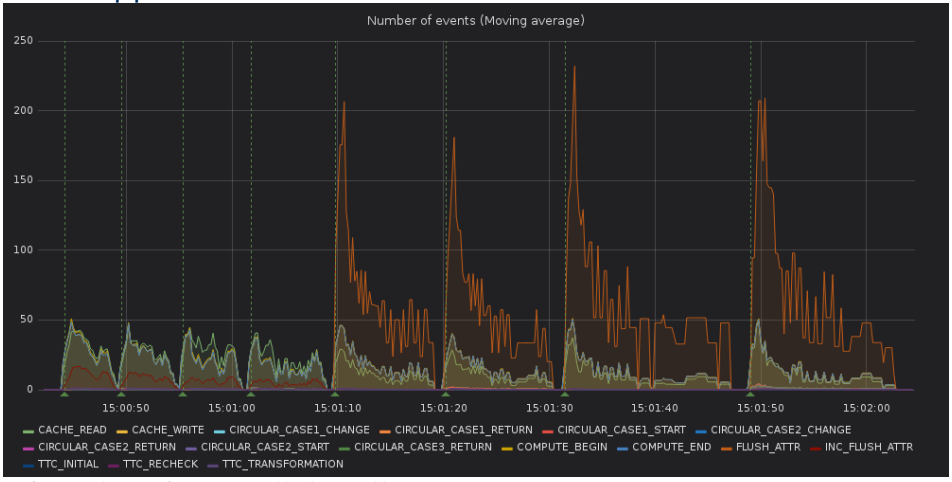
Tool Support



JastAdd Gradle plugin

Tracing API

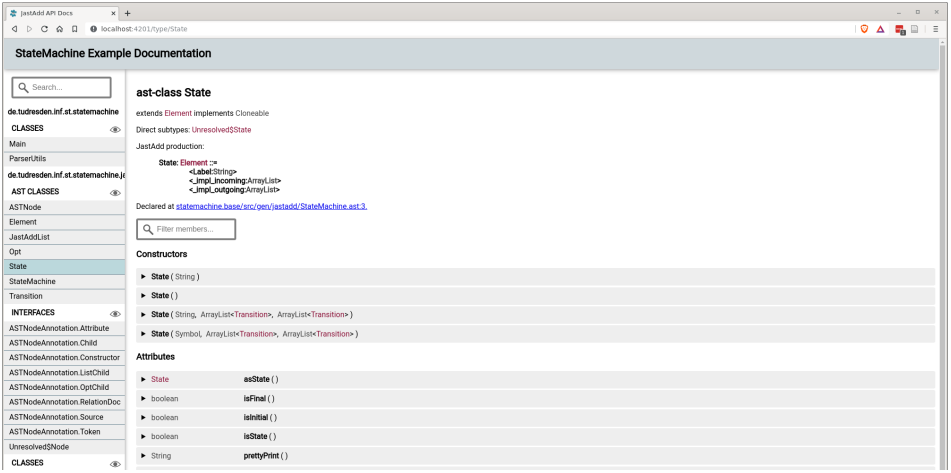
Tool Support



Grafana visualization of events created by the JastAdd tracing API

Documentation Generation: RAGdoc

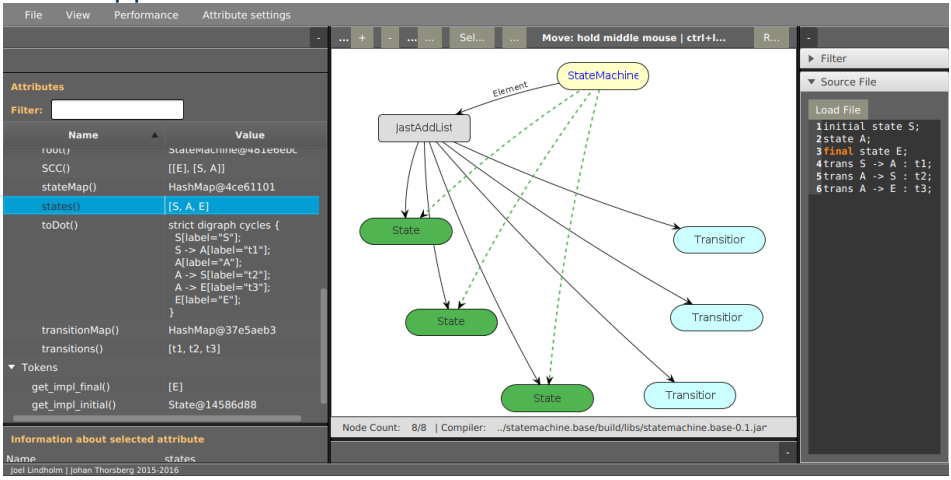
Tool Support



RAGdoc documentation including links to source code

Visualization and Debugging: DrAST

Tool Support



DrAST visualization with computed attributes

JastAdd
Important Information

JastAdd
- Website with reference manual and bibliography `www.jastadd.org`
- Source code `https://bitbucket.org/jastadd/jastadd2`

Build tool support
- gradle/maven/...packages:
  - `org.jastadd:jastadd`, `org.jastadd:jastaddparser`, `org.jastadd:jastaddgradle`
- gradle plugin: `https://bitbucket.org/jastadd/jastaddgradle/`

DrAST
- code and doc: `https://bitbucket.org/jastadd/drast/`

RagDoc
- code and doc: `bitbucket.org/extendj/ragdoc-builder/` , `bitbucket.org/extendj/ragdoc-view/`

Questions so far?

References I

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Mathematical systems theory, 2(2):127–145.

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