



## Part III: Tool Integration Architectures and Macromodels

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teaching/most](http://st.inf.tu-dresden.de/teaching/most)  
Version 21-0.4, 18.12.21



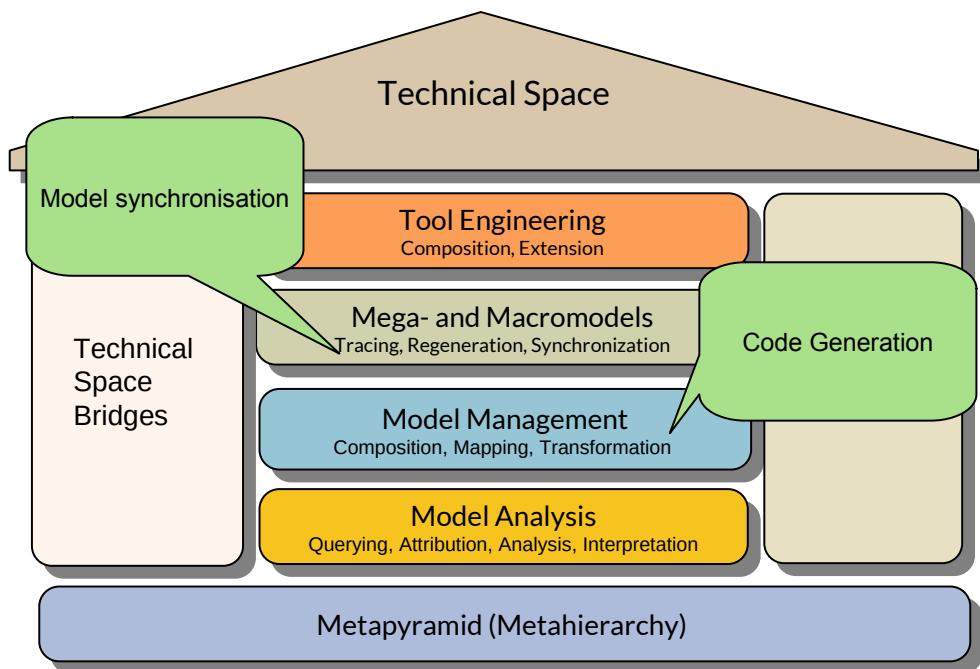
## 30. Model Synchronisation, Code Generation and Round-Trip Engineering for the Consistency of Macromodels

### Code Generation as Apps for RAG

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teaching/most](http://st.inf.tu-dresden.de/teaching/most)

- 0) Tool Integration Architectures
- 1) Single-source principle and macromodel principle
- 2) Code generation techniques
  - Template-based Code generation
  - 3) Re-parsing

## Q10: The House of a Technical Space



## Literature

- ▶ <http://www.codegeneration.net/>
- ▶ [www.programtransformation.org](http://www.programtransformation.org)
- ▶ [http://www.codegeneration.net/tiki-read\\_article.php?articleId=65](http://www.codegeneration.net/tiki-read_article.php?articleId=65)
- ▶ Paul Bassett. Frame-based software engineering. *IEEE Software*, 4(4):9-16, 1987.
  - <http://doi.ieeecomputersociety.org/10.1109/MS.1987.231057>
- ▶ Chris Holmes, Andy Evans. A review of frame technology. University of York, Dept. of Computer Science, 2003  
<ftp://www-users.cs.york.ac.uk/reports/2003/YCS/369/YCS-2003-369.pdf>
- ▶ Daniel Weise and Roger Crew. Programmable syntax macros. In Proceedings of the ACM SIGPLAN '93 Conference on Programming Language Design and Implementation, pages 156-165, Albuquerque, New Mexico, June 23-25, 1993.
- ▶ Optional
  - Völter, Stahl: Model-Driven Software Development, AWL 2005.
  - Falk Hartmann. Safe Template Processing of XML Documents. PhD thesis, Technische Universität Dresden, Fakultät Informatik, July 2011.
    - <http://nbn-resolving.de/urn:nbn:de:bsz:14-qucosa-75342>



## 30.0 Concepts of Tool Integration for Software Factories

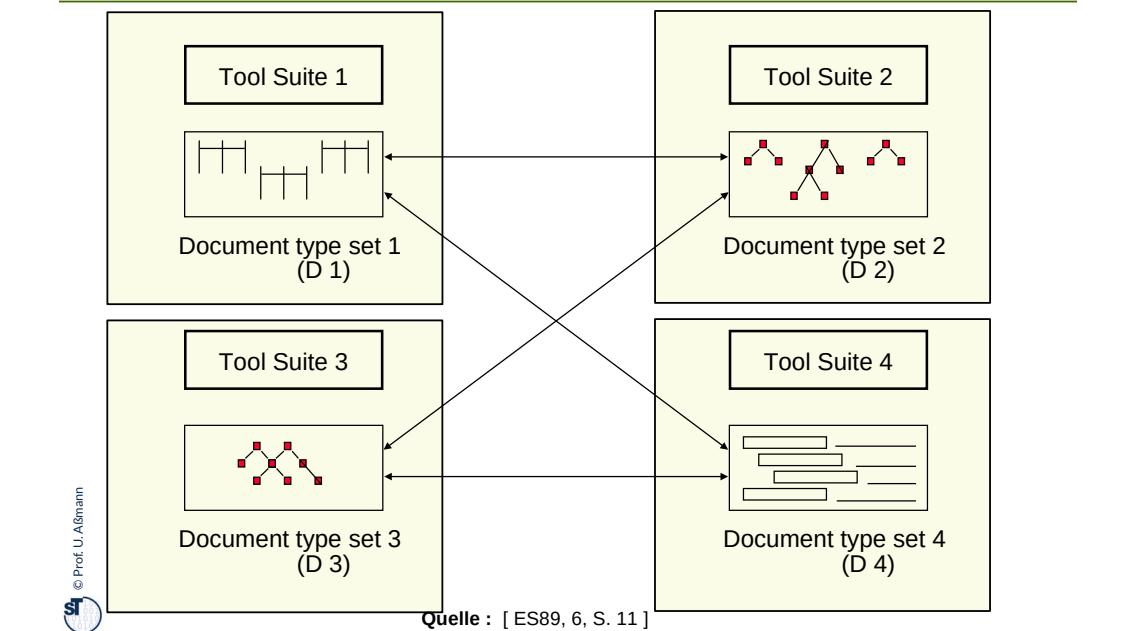
The following integration techniques for standalone tools can be used for

- ▶ Enterprise Application Integration (EAI)
- ▶ Distributed Software Systems
- ▶ Software Factories, MetaCASE tools, IDE
- ▶ Heterogeneous Software Factories



## Integration of Standalone, Persistent Tool Suites by Data Connection

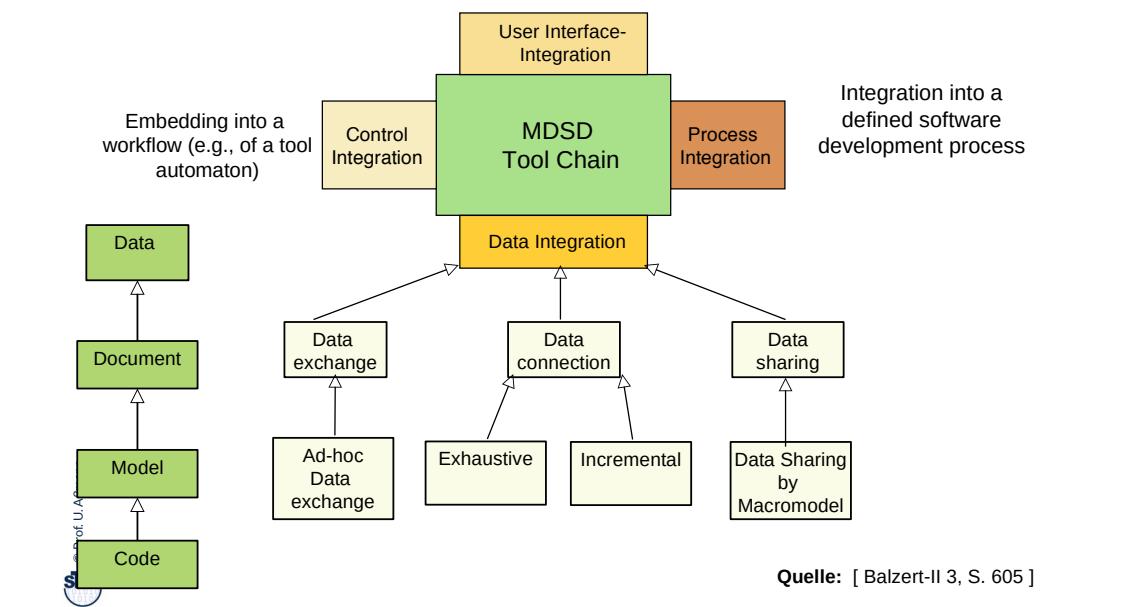
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## Tool Integration in 4 Dimensions

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## 53.2 Data Integration



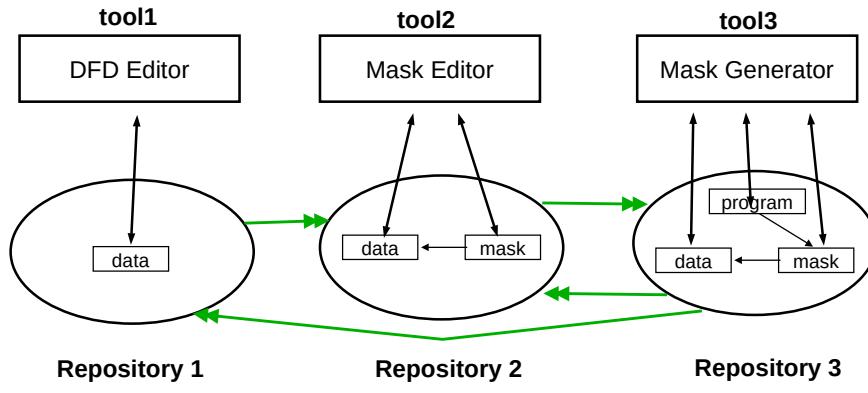
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# Levels of Data Integration

Integration Kind	Schema	Characteristics
Black-box-Integration (weak)	<pre> graph TD     CASE[CASE tool] --&gt; Files[Files]     CASE --&gt; Repo[Repository]     Files &lt;--&gt; Repo     style CASE fill:#e0e0ff     style Files fill:#e0e0ff     style Repo fill:#e0e0ff     </pre>	<ul style="list-style-type: none"> <li>Tool works isolated on own data</li> <li>Repository coordinates data with check in and check out!</li> <li>mostly directory based, e.g., git</li> </ul>
Grey-box-Integration (medium)	<pre> graph TD     CASE1[CASE tool] --&gt; Repo[Repository]     CASE2[CASE tool] --&gt; Repo     CASE3[CASE tool] --&gt; Repo     Repo --&gt; Data1[Data]     Repo --&gt; Data2[Data]     style CASE1 fill:#e0e0ff     style CASE2 fill:#e0e0ff     style CASE3 fill:#e0e0ff     style Repo fill:#e0e0ff     style Data1 fill:#e0e0ff     style Data2 fill:#e0e0ff     </pre>	<ul style="list-style-type: none"> <li>Separate repository, but common access interfaces             <ul style="list-style-type: none"> <li>Access layer to repositories</li> </ul> </li> </ul>
White-box-Integration Data sharing (strong)	<pre> graph TD     CASE1[CASE tool] --&gt; Repo[Repository]     CASE2[CASE tool] --&gt; Repo     Repo --&gt; Data[Data]     Repo --&gt; Schema[Schema]     style CASE1 fill:#e0e0ff     style CASE2 fill:#e0e0ff     style Repo fill:#e0e0ff     style Data fill:#e0e0ff     style Schema fill:#e0e0ff     </pre>	<ul style="list-style-type: none"> <li>Definition of uniform data schema (material metamodels) for all tools</li> <li>Integration of data schemata of tools by             <ul style="list-style-type: none"> <li>Role-based integration</li> <li>Inheritance-based integration</li> <li>dependent on Technical Space</li> </ul> </li> </ul>

## a) Ad-hoc Data Exchange (Simple)

- ▶ No common data, high cost for exchange
- ▶ Exchange with Streams and Data-Flow Architecture (Example: UNIX shell)
  - Queries filter data, transformations change data
  - Data formats are defined in an exchange language
- ▶ Tools are rather independent

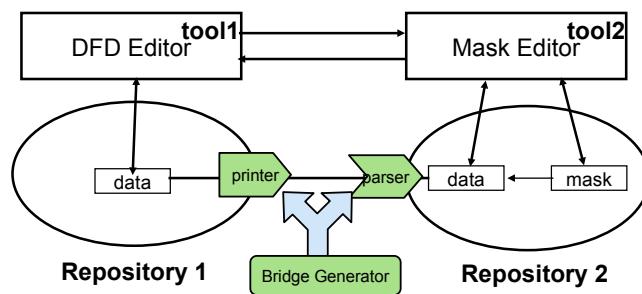


Quelle: nach [ HMF. S. 196 ]

## b) Data Connection with Systematic Data Exchange (Transformation Bridges)

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- ▶ **Data connection (Datenverbindung)** relies on the definition of *mappings between the data schemata (material metamodels)*
- ▶ **Automated data exchange** in standardised exchange formats (**Exchange formats**, such as ASN, XMI, JSON, CDIF, XML)
  - Automation (generation of parsers and printers) relies on mappings between the material metamodels (language mappings)
  - Use as a technical space for the exchange format a link-tree metalanguage (XML, RAG)
- ▶ **Transformation bridge:** a prettyprinter transforms the internal representation of a repository into an external exchange format
  - Parser reads the exchange format again and transforms it into the internal representation of the other repository
  - Query and transformation languages filter the data

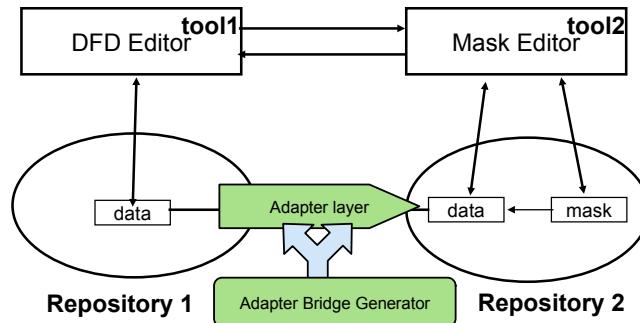


### c) Data Connection with Incremental Data Exchange (Adapter Bridges)

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- ▶ An **Adapter Bridge** is a layer between two repositories allowing for *incremental access* from one to the other
  - Transformation of data incrementally with each access
  - Direct transformation without exchange format



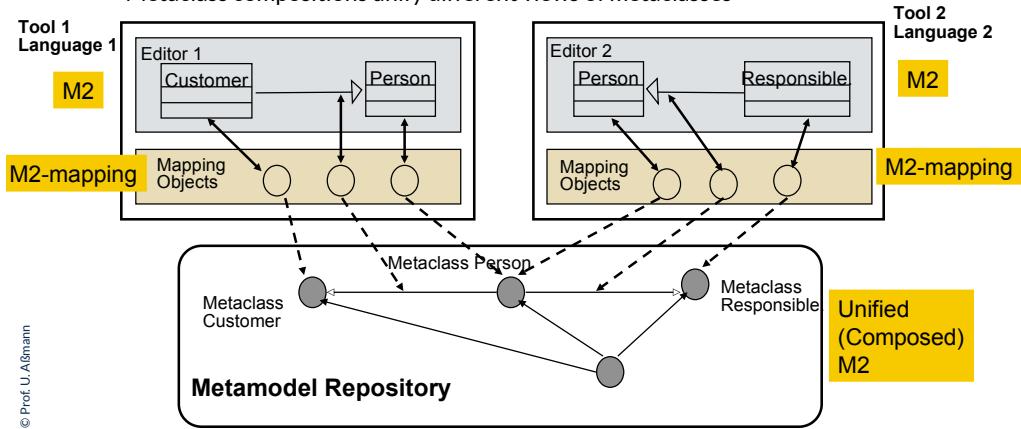
## Current Exchange Formats for Data Connection

An **exchange format** is a standardized format for the exchange of data between tools, also of different vendors

- ▶ Comma-separated values (CSV): simple text-based exchange format for tools on relation, text algebra, and tables (Excel, TeX, ...)
  - No metalanguage but simple table schema <http://tools.ietf.org/html/rfc4180>
  - [http://en.wikipedia.org/wiki/Comma-separated\\_values](http://en.wikipedia.org/wiki/Comma-separated_values)
- ▶ XML Metadata Interchange (XMI) for exchange of UML-diagrams in XML-format
  - Meta Object Facility (MOF) as metalanguage
  - <http://www.omg.com/technology/documents/formal/xmi.htm>
- ▶ JSON hierachic maps
  - TOML, YAML variants are less wordy
- ▶ ASN.1 Standard is a metalanguage based on BNF
  - [http://de.wikipedia.org/wiki/Abstract\\_Syntax\\_Notation\\_One](http://de.wikipedia.org/wiki/Abstract_Syntax_Notation_One)
- ▶ RDF/RDFS Resource Description Format – Models as graphs, stored in elementary triplets  
<http://www.w3c.org>
- ▶ GXL Graph Exchange format: Open Source Format for exchange of graphs
  - <http://www.gupro.de/GXL/>
- ▶ Historic:
  - CASE Tool Data Interchange Format (CDIF) - metalanguage ERD for data definition as well as
  - Data Flow Model, State Event Model, Object Oriented Analysis and Design
  - <http://www.ecma-international.org/publications/files/ECMA-ST/Ecma-270.pdf>

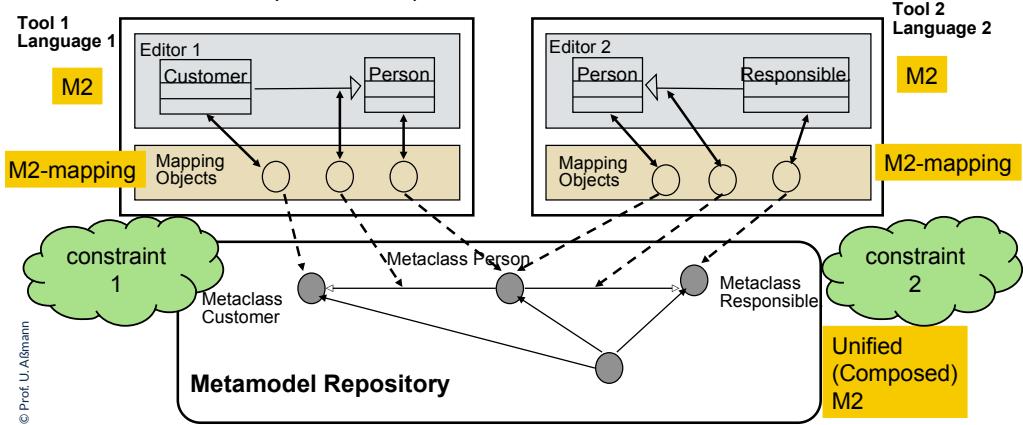
## d) Data Sharing (With Common Repository)

- With data sharing (**Datenteilung**) all tools share common data with a uniform schema (material metamodel)
- Metaclass mappings control the integration of the repositories and metamodels
- Metaclass compositions unify different views of metaclasses



## e) Data Sharing with Macromodels

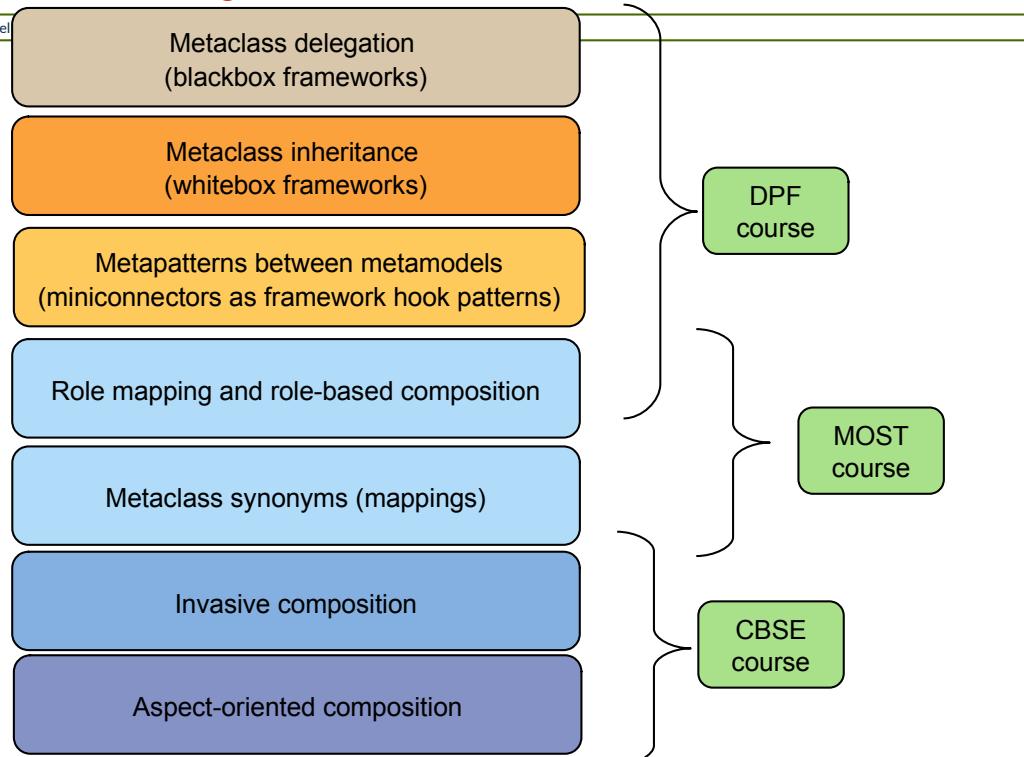
- With **data sharing by macromodels** all tools share common model with a uniform metamodel **underlying wellformedness constraints**
- Metaclass mappings control the integration of the repositories and metamodels
- Metaclass compositions unify different views of metaclasses



Quelle: Platz, D., Kelter, U.: Konsistenzerhaltung von Fensterinhalten in SEU; <http://pi.informatik.uni-siegen.de>

# How to Compose and Relate Metaclasses and Metamodels for Data Sharing?

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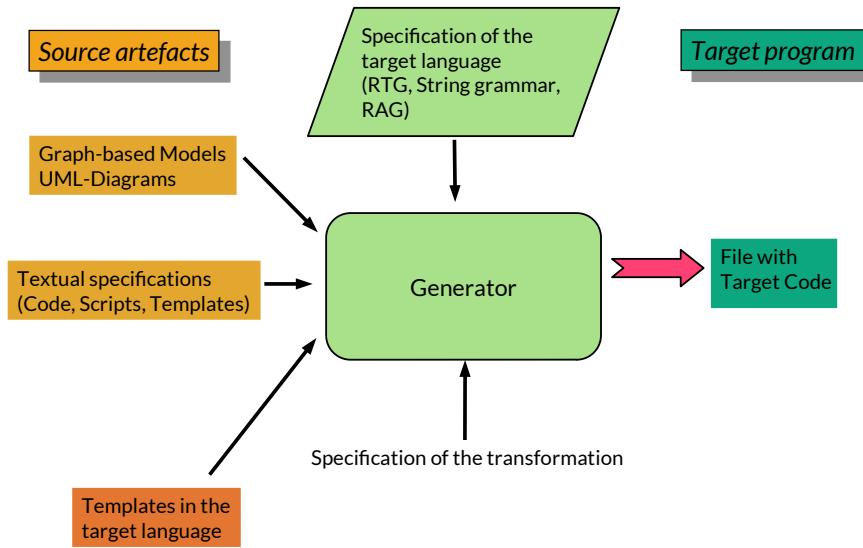
## 30.1 Model2Code Transformation (Code Generation)

Transforming models into code  
(Programmüberführung)

## MDSD-Code-Generators

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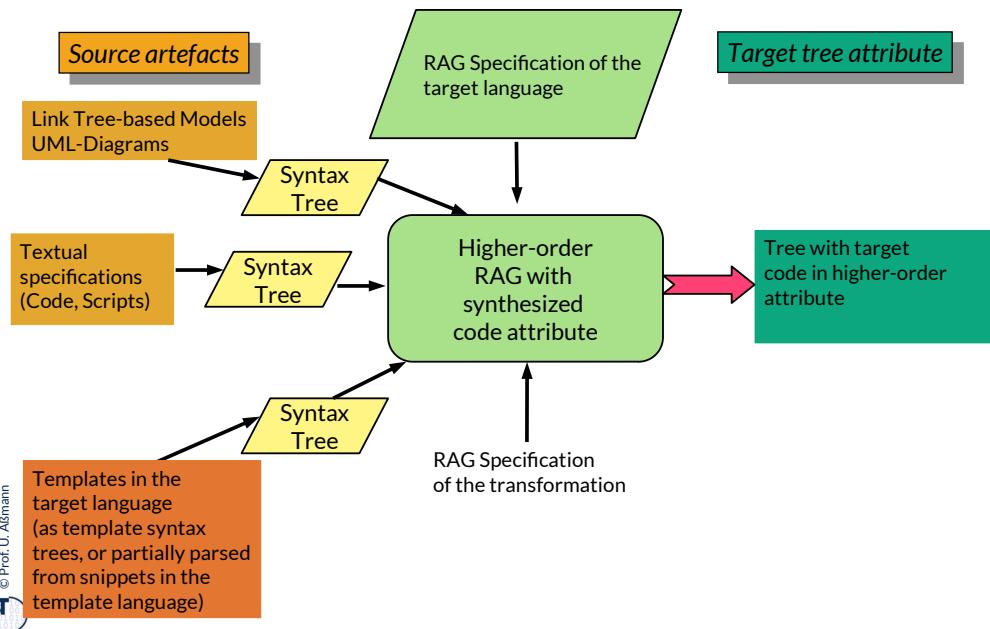
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## MDSD-Code-Generators as Attributors of Syntax Trees

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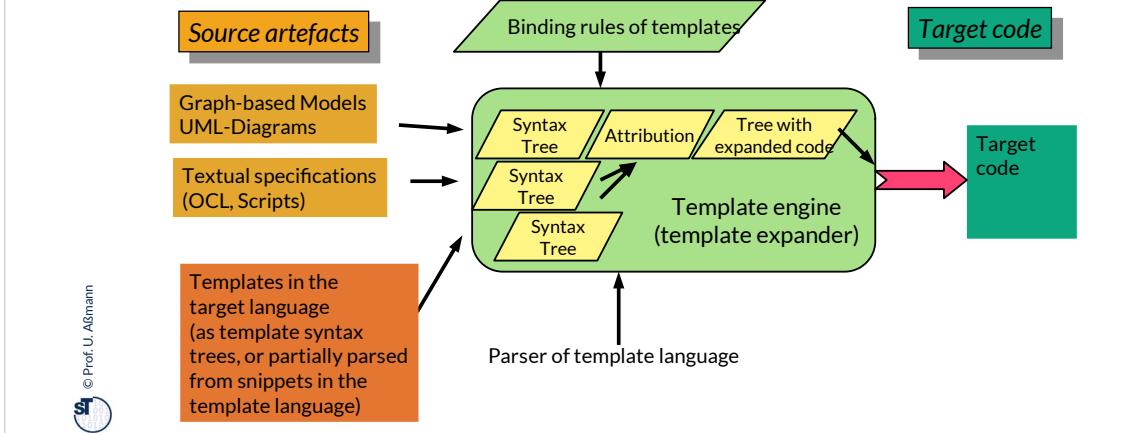


## Different Kinds of Code Generators

- ▶ A **code selector** is a transformation system (on strings, terms, link trees, graphs) covering the input models with rules (**code coverage**) transforming the model elements once
- ▶ A **code scheduler** orders instructions in an optimized manner
  - Code scheduling runs after code selection
- ▶ **Metaprogramming code generators:**
  - A **template expander** generates code by filling code templates with *inset snippets*
  - An **invasive fragment composer (invasive software composition)** composes templates in a typed and wellformed way (→ CBSE)

## MDSD-Code-Generators as Template Expanders

- ▶ A **template engine** hides the tree construction, attribution with code attributes, and pretty-printing under a simple interface. It provides functions
  - `templparse(): String in TemplateLanguage → Tree`
  - `pparse(): String in BaseLanguage → Tree`
- ▶ Template engines are *apps* of higher-order RAGs





### 30.1.2 Code Generation in RAGs

- ▶ With higher-order (tree-generation) attributes and special functions
  - partial parsing
  - template expansion



## Code Generation with RAGs

- ▶ Attribution functions may generate code syntax trees
- ▶ Suppose a *partial parse function* pparse(): String->LinkTree

```
eq Constant.Code() {
    if (AsBoolean())
        if (AsValue() == 1)
            return pparse("(boolean)1");
        else if (AsValue() == 0)
            return pparse("(boolean)0");
        else return EmptyTree;
    else {
        if (AsValue() == 1)
            return pparse("new Integer(1)");
        else if (AsValue() == 0)
            return pparse("new Integer(0)");
        else
            return pparse("new Integer("+AsValue()+")");
    }
}
```

## Template-Based Code Generation with RAGs

- ▶ Attribution functions may expand code templates to code trees
- ▶ Done with the ***template processing function***  
`templparse(): String, List(ID)->LinkTree` that expands variable names into attribution functions, e.g., `TypeParameterName` → `TypeParameterName()`
- ▶ `templparse()` is called a *template processor*, `String` is of a *template Language*

```
eq GenericClassInstantiation.Code() {
    return templparse(
        "public class GenClass$TypeParameterName$ extends Object {
            private int myId;
            public GenClass$TypeParameterName$() { // constructor
            }
            public int getId() { return myId; }
        }"
        , List(pparse("Person"))
    );
}
```

## Template-Based Code Generation with RAGs

- ▶ The **template processing function** can be made generic in terms of grammars:  
templparseGeneric(): CSGrammar, RTG, String, List(ID)->LinkTree that expands variable names into attribution functions, e.g., TypeParameterName → TypeParameterName()
- ▶ templparse() is called a *template processor*, String is of a *template language*

```
eq GenericClassInstantiation.Code() {
    CSGrammar CSAceleo;
    RTGrammar RTGAceleo;
    return templparseGeneric(CSAceleo, RTGAceleo,
        "public class GenClass$TypeParameterName$ extends Object {
            private int myId;
            public GenClass$TypeParameterName$() { // constructor
            }
            public int getId() { return myId; }
        }"
        , List(pparse("Person"))
    );
}
```



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### 30.1.3 Single-Source Principle and Macromodels



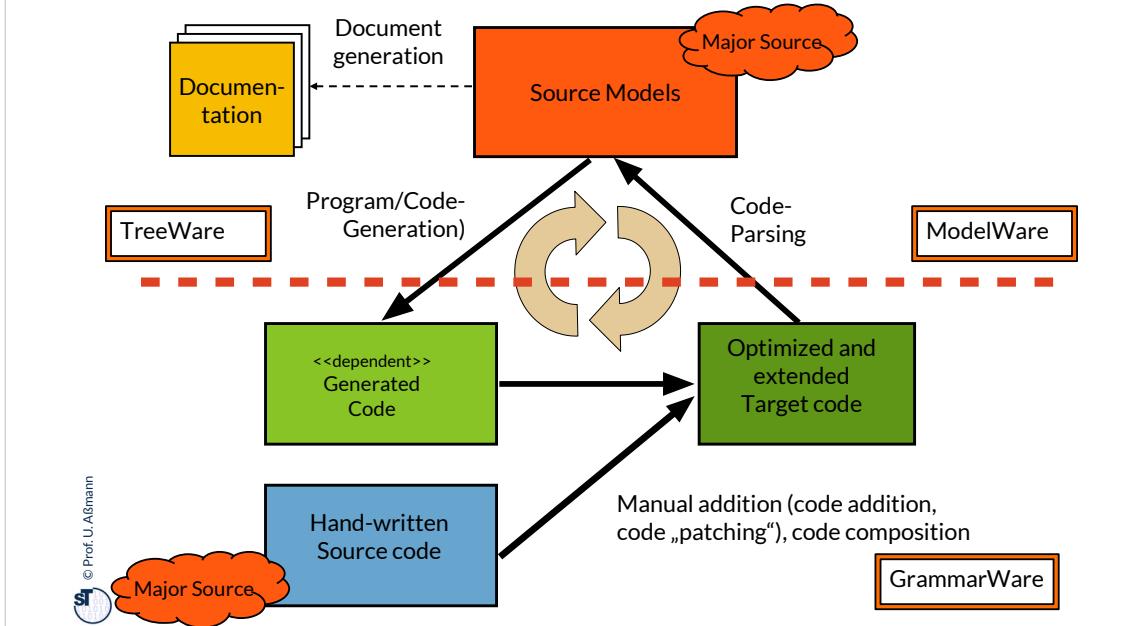
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## Single-Source-Principle, Major-Source, Code Addition, and Round-Trip Engineering

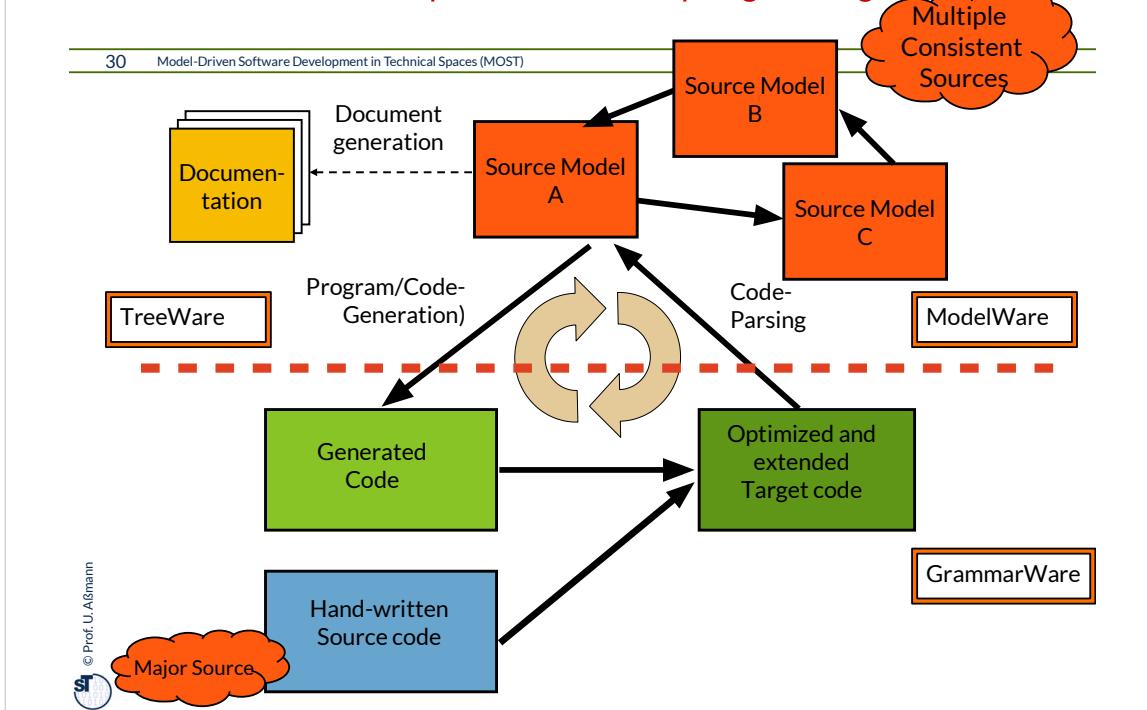
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## Macromodel Principle and Round-Trip Engineering

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## Single-Source Principle, Major-Source Principle and Macromodel Principle

- ▶ A **Single-Source-Technology** uses a *single source* (one set of *major-source* models), from which code, tests (derived models) are derived, with automatic synchronisation and consistency
- ▶ A **Macromodel Technology** is a single-source technology with automatic synchronisation and consistency between ALL (major-source) models, code, tests, and documentation (all models of a multi-model)
- ▶ In a macromodel, there are always ***derived models***
  - Generated code (this chapter)
  - Generated documentation (Chapter on documentation)
  - Generated test suites and data

## Synchronization is Used in Round-Trip Engineering

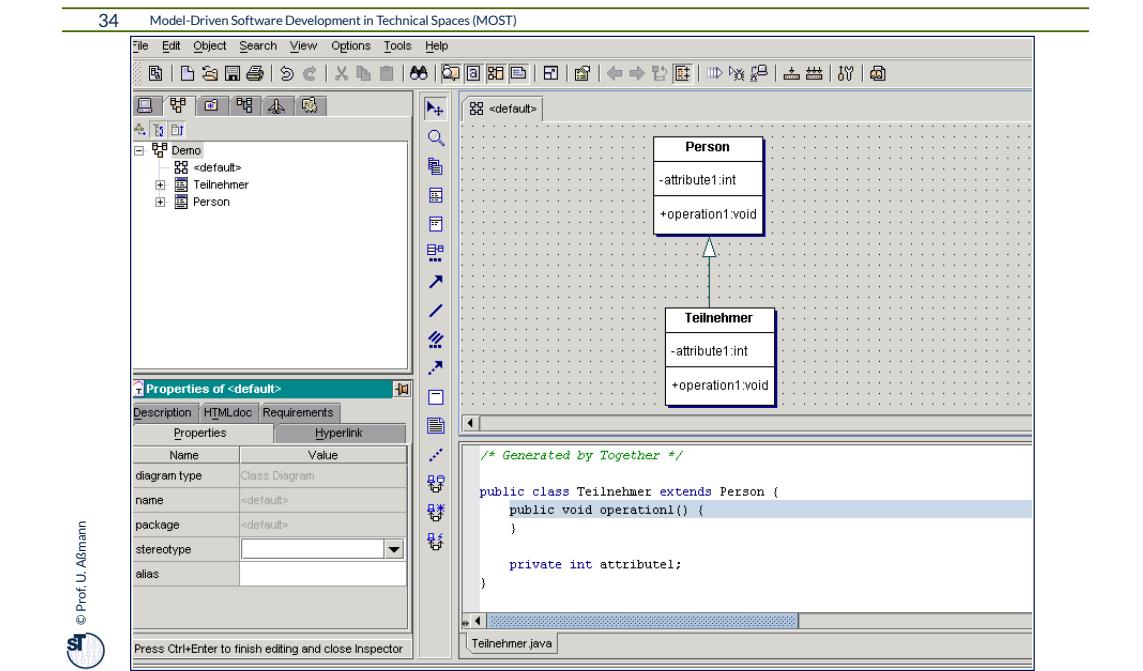
- ▶ Technically, the Single-Source-Principle and the Macromodel principle needs **Round-Trip-Engineering (RTE)** between ModelWare and GrammarWare, to achieve
  - **Model-to-code synchronisation** with
    - **Codegeneration** into several programming languages
    - **Template-based codegeneration** inserts code snippets into code templates
    - **Code reparsing** of the changed source code into models
  - **Model-to-model synchronization** (later) with
    - **Bidirectional transformations** (e.g., with Triple-Graph-Grammars, TGG)
    - **View based transformations** (with e.g., with Single-Underlying Model, SUM)

## Example: Round-Trip Engineering in Together (P. Coad, Borland)

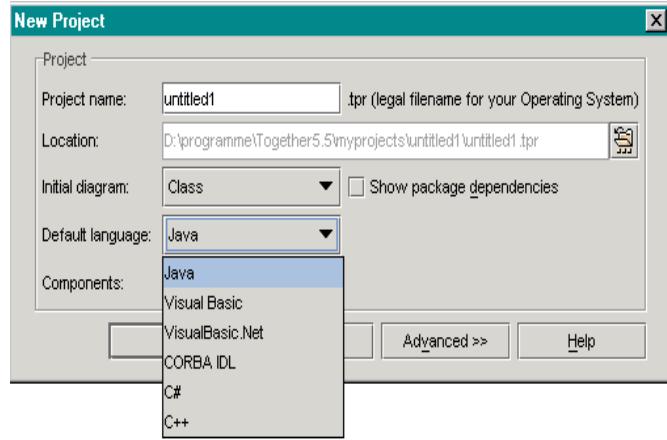
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- ▶ In 1997, the CASE tool **Together** was the first to provide a Single-Source-Technology with automatic synchronisation and consistency between UML model, code and documentation
- ▶ Supported Programming Languages: Java, Visual Basic, VisualBasic.Net, CORBA IDL, C++, C#
  - Synchronisation by reparsing of generated, modified and extended code
- ▶ Round-trip Engineering:
  - Changes of class diagrams will be transformed to code
  - Changes of code reparsed to class diagrams
  - Reverse Engineering of entire projects

## Together Screenshot



## Code Generation in Different Languages in Together



- Supports roles: Business Modeler, Designer, Developer and Programmer
- Appropriate views can be configured
- Code template based code generation



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## 30.2 Technologies for Model-2-Code Generation and Synchronization

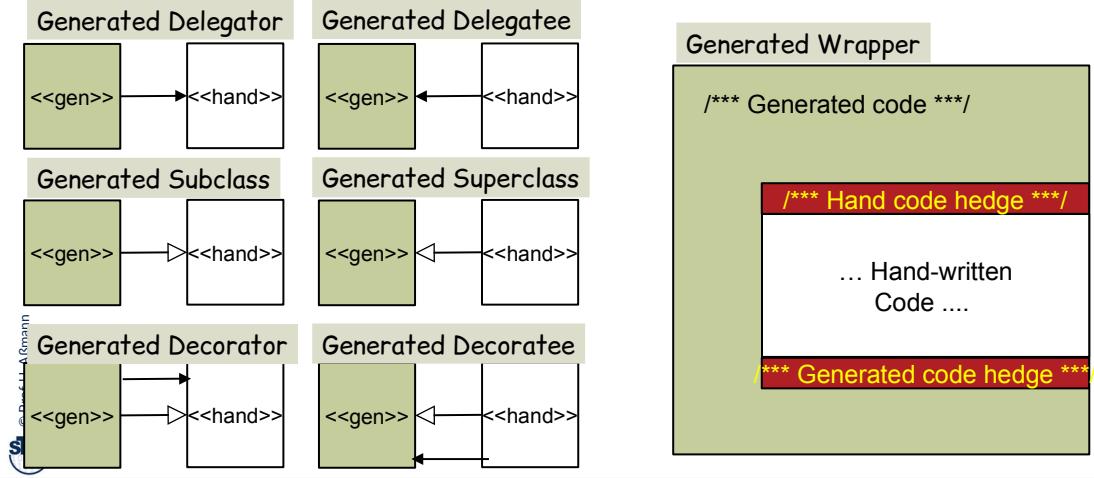


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# Composition of Separated Hand-Written and Generated Code

- ▶ In separate files: Coupling by implementation pattern [Völter/Stahl]
- ▶ Use class composition like delegation, TemplateMethod, Composite, Decorator, etc
- ▶ Synchronization is easy: do not touch generatees
- ▶ In one file: Coupling with **hedges** (**Trennmarkierung**)
- ▶ Synchronization should stay out of hedged areas



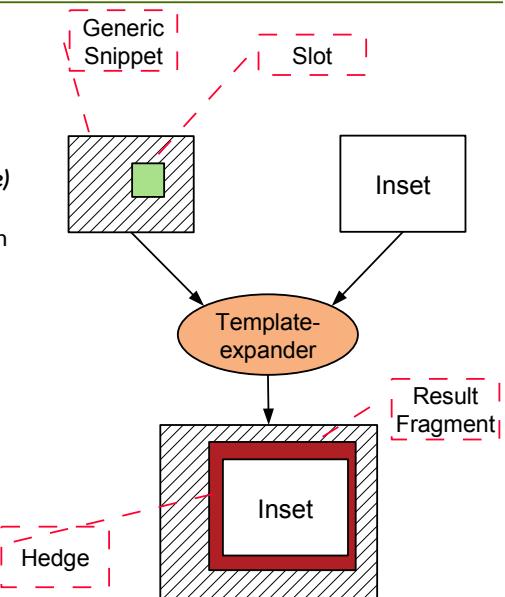
## Composition of Generated and Hand Written Code in an RAG

- ▶ Fine-grained glueing possible
- ▶ Wrapping hedges around synthesized snippets

```
eq Procedure.Code() {
    return Head.Code()+
        “/** HEDGE2Gen BEGIN */“+
        GeneratedBody.Code()+
        “/** HEDGE2Gen END */“;
}
eq Head.Code() {
    return pparse(“public „+Head.name+“()”);
}
eq GeneratedBody.Code() {
    return Body.Code();
}
```

## Snippet and Template Programming with RAG

- ▶ A **fragment (snippet)** is a incomplete sentence of a language, derived from a nonterminal of the grammar, or described by a metaclass
- ▶ A **generic fragment (template, form, frame)** is a fragment with slots (holes, code parameters, variation points), which can be bound (filled, expanded) with an **Inset** fragment to a result fragment
- ▶ A **extensible fragment** is a fragment with hooks (extension points), which can be extended to a fragment
- ▶ **Generic programming** is programming with generic fragments (templates).
- ▶ **Invasive programming** is programming with generic and extensible fragments (templates with hooks)
- ▶ → CBSE course





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## 30.2.1 Template-based Code Generation (Schablonenbasierte Programmüberführung)



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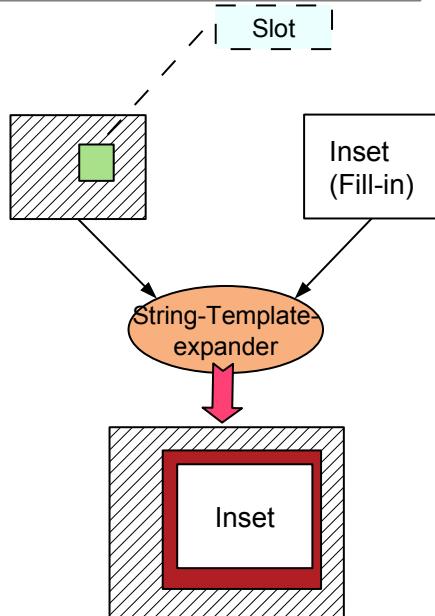
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## Template Expansion by Composition of Insets and Hedging

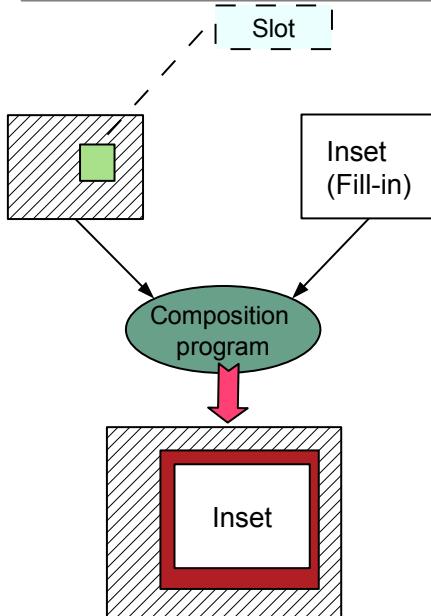
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### Coupling by string expansion



### Coupling by composition program



## Slots are Marked by Hedges

- ▶ **Hedges** are delimiters that do not occur in the base nor in the slot language
- ▶ **Slot hedges** are template2slot hedges marking the transition from the code language to the slot language
- ▶ **Inset hedges** are metaprogramming2code hedges marking the transition from the metaprogramming language to the code language

```
// code slot hedges << >>
template (superclass:CLASS, t:TYPE) {
    class Worker extends <<superclass>> {
        <<t>> attr = new <<t>>();
        <<t>> getAttr();
        void setAttr(<<t>> a);
    }
}
```

# Tools for Untyped Template Expansion

- ▶ **Frame processing** was invented in [P. Bassett] as an *untyped string template expansion technology*, universal for all textual languages [Holmes/Evans]
  - Frame processing is the main technology for web engineering today: it organizes reuse of page templates
  - The original frame processor used \$ as a hedge symbol for slots (slot variables)
- ▶ **Macro processing** is not much different
  - Because only slot variables hold insets, macro parameters correspond to slot variables
- ▶ **XML template processing** engine XVCL [Jarzabek] is an XML-controlled frame processor
  - <http://sourceforge.net/projects/fxvcl/files/XVCL%20Specification/Version%202.10/>
- ▶ **String template engines** in use today
  - Apache Velocity <http://velocity.apache.org/>
  - Parr's template engine StringTemplate
  - Jenerator for Java <http://www.voelter.de/data/pub/jeneratorPaper.pdf>
  - Acceleo <https://www.eclipse.org/acceleo/>

## Velocity String Template Language

- ▶ Velocity Template Language (VTL) is a frame processing language with
  - metaprograms in slots, written in a **slot language (blue)**
- ▶ {#, \$} are slot hedges
- ▶ < (from XML) is the inset hedge

```
<html>
<body>
#set( $foo = "Velocity" )
Hello $foo World!
</body>
</html>
```

```
<HTML>
<BODY>
Hello $customer.Name!
<table>
#foreach($mud in $mudsOnSpecial)
  #if
    ($customer.hasPurchased($mud))
      <tr>
        <td>
          $logger.getPromo( $mud )
        </td>
      </tr>
    #end
#end
</table>
```

<http://velocity.apache.org/engine/releases/velocity-1.7>

## Velocity Template Language

- ▶ Velocity Template Language (VTL) is a simple scripting language in the spirit of TCL
- ▶ It has control structures (if, switch, foreach), assignments (set), and macros
- ▶ Similar: Acceleo (in exercises)

```
#macro( inner $foo )
    inner : $foo
#end

#macro( outer $foo )
    #set($bar = "outerlala")
    outer : $foo
#end

#set($bar = 'calltimelala')
#outer( "#inner($bar)" )
```

**Problem:** the result of string template expansion  
may not be syntactically correct, nor well-formed, target language  
(error-prone)

## Typed Template Expansion

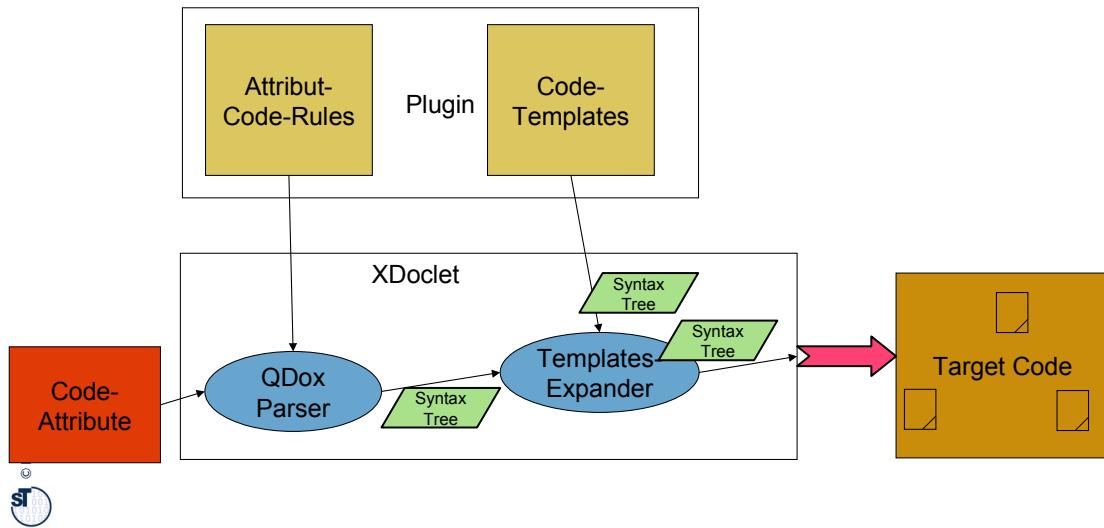
- ▶ Metamodel-controlled template engines
  - EMF: Xtend and Xpand scripting languages
  - XML slot markup language
  - Acceleo code generating system (see exercises)
- ▶ Invasive Software Composition provides fully typed and wellformed template expansion (see CBSE course)
  - Typed template expansion **and** -extension **and** weaving
  - Can be instantiated for arbitrary languages
  - <http://www.the-compost-system.org> (obsolete now)
  - <http://www.reuseware.org>
  - <https://bitbucket.org/svenkarol/skat/wiki/Home>

## Semantic Macros

- ▶ **Semantic Macros** are metaprogramming procedures which are typed parameters and results.
  - A semantic macro is compiled to a fragment tree
    - which is instantiated by fragment parameters, type-checked on the metamodel, and copied to the instantiation spot
  - They allow for type-safe static metaprogramming.
  - In an higher-order RAG, a semantic macro can be instantiated in a higher- order attribute
- ▶ Examples:
  - Scheme
  - Scala <http://scalamacros.org/>
  - <http://docs.scala-lang.org/overviews/macros/overview.html>

## Xdoclet (xdoclet.sf.net) for Metadata-Based Code Generation

- ▶ Xdoclet transforms attributes (metadata) into helper code (aka boilerplate code)
  - With template-based code generation
  - Metadata attributes *trigger* the filling of templates, used from a library



## Slot Markup Languages

- ▶ A **slot markup language** is a special template language for *any* XML dialect
- ▶ The slot language is represented as an XML dialect itself (XSD schema) [Hartmann]
  - Uniform syntax for templates
  - XML tools are usable
  - Filling templates is
    - type-safe
    - and wellformed, because OCL constraints can be defined that are checked

# The End

- ▶ Why is code generation a good application for RAG?
- ▶ How would you generate code with Xcerpt?
- ▶ Explain the difference of the code generation patterns GeneratedDelegatee, GeneratedDelegator, GeneratedSuperClass, GeneratedSubclass!
- ▶ Why does code generation most often use synthesized attributes?
- ▶ What is the difference of a metadata attribute (annotation), and an attribute in an RAG?
- ▶ Why are template engines apps for RAGs?
- ▶ Think about GOTO statements in machine code, or in C programs.
  - How would you represent them in an RAG?
  - Why are AG not really appropriate for representing GOTOS?



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## 30.A.1. Code Modification and Reparsing (Codemodifikation und -rückführung)



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concept  
Exzellenz aus  
Wissenschaft  
und Kultur

Model-Driven Software Development in Technical Spaces (MOST) © Prof. U. Aßmann

## Example of Code Reparsing Technique

- ▶ Code-Reparsing in Fujaba:  
[http://www.fokus.fraunhofer.de/en/fokus\\_events/motion/ecmda2008/\\_docs/  
rs01\\_t03\\_ManuelBork\\_EMCD2008\\_slides.pdf](http://www.fokus.fraunhofer.de/en/fokus_events/motion/ecmda2008/_docs/rs01_t03_ManuelBork_EMCD2008_slides.pdf)
- ▶ Parallel Parsing of Template and Generated Code, with comparison to resolve  
indeterministic situations of re-parsing