



## IV. The Technical Space Graphware

### 40. Flat Analysis in Graphware: Graph Querying, Metrics, Reachability Analysis and Megamodel Dependency Analysis

Prof. Dr. U. Aßmann  
Technische Universität Dresden  
Institut für Software- und  
Multimediatechnik  
<http://st.inf.tu-dresden.de>  
Version 21-1.2, 29.01.22

- 1) Graph-Based DDL and CDL
  - 1) Relational Schema
  - 2) Entity-Relationship Diagrams
  - 3) MOF and ERD
- 2) Graph Query Languages
  - 1) QL
  - 2) Metrics with QL
  - 3) Lifting Info to the Macromodel Level
  - 4) Macromodel Dependency Analysis
  - 5) Other graph query languages

In a tree, there are paths bottom-up and top-down.

In a graph, there are paths everywhere.

If one removes the color from links in a link tree, graphs result.

# Obligatory Literature

- ▶ [http://en.wikipedia.org/wiki/List\\_of\\_UML\\_tools](http://en.wikipedia.org/wiki/List_of_UML_tools)
- ▶ [http://en.wikipedia.org/wiki/Entity-relationship\\_model](http://en.wikipedia.org/wiki/Entity-relationship_model)
- ▶ <http://www.utexas.edu/its/archive/windows/database/datamodeling/index.html>
- ▶ [deMoor] Oege de Moor, Mathieu Verbaere, Elnar Hajiyev, Pavel Avgustinov, Torbjorn Ekman, Neil Ongkingco, Damien Sereni, Julian Tibble, "Keynote Address: .QL for Source Code Analysis", SCAM, 2007, 2013 IEEE 13th International Working Conference on Source Code Analysis and Manipulation (SCAM), pp. 3-16, doi:10.1109/SCAM.2007.31
- ▶ CodeQL is free now (via github): <https://github.com/github/codeql>
- ▶ <https://semml.com/codeql>, <https://help.semml.com/QL/learn-ql/>
- ▶ <https://help.semml.com/QL/learn-ql/java/ql-for-java.html>
- ▶ Language handbook <https://help.semml.com/QL/ql-handbook/index.html>
  - Specification <https://help.semml.com/QL/ql-spec/language.html>
- ▶ Thief detective game: <https://help.semml.com/QL/learn-ql/beginner/find-thief-1.html>
- ▶ Industrial case studies: <https://semml.com/case-studies>
- ▶ Community-driven security analysis:
  - Github repo of LGTM examples <https://github.com/Semml/ql>  
<https://securitylab.github.com/tools/codeql> <https://lgtm.com/help/lgtm/about-lgtm>
  - Query console <https://lgtm.com/query>
  - <https://lgtm.com/help/lgtm/console/ql-java-basic-example>

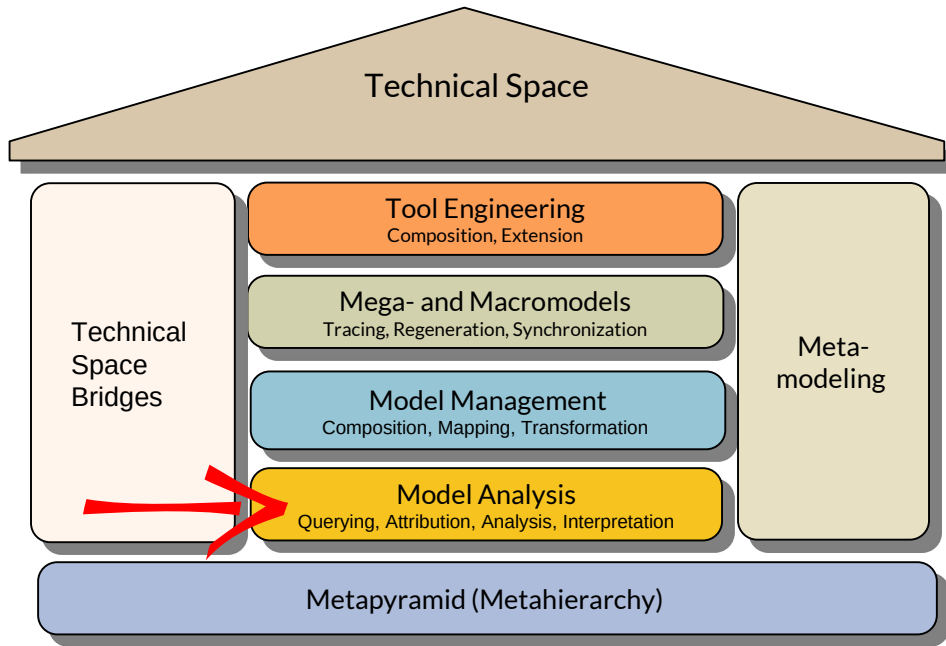


# References

- ▶ [Chen] P. P.-S. Chen. The entity-relationship model - towards a unified view of data. Transactions on Database Systems, 1(1):9-36, 1976
- ▶ A Comparison of ATL and Story-Driven Modeling (Fujaba-style GRS)

[http://www.es.tu-darmstadt.de/fileadmin/download/publications/spatzina/PP\\_AGTIVE\\_2011.pdf](http://www.es.tu-darmstadt.de/fileadmin/download/publications/spatzina/PP_AGTIVE_2011.pdf)

# Q10: The House of a Technical Space



# Q11: Overview of Technical Spaces in the Classical Metahierarchy

	Gramm arware (String s)	Text- ware	Table-ware		Treeware (trees)	Link-Tree- ware		Graph ware/ Model ware			Role- Ware	CROM- Ware	Ontology -ware
	Strings	Text	Text- Table	Relationa l Algebra	NF2	XML	Link trees	MOF	Eclipse	CDI F	MetaEdit+	Context- role graphs	OWL-Ware
M 3	EBNF	EBNF		CWM (common warehouse model)	NF2- language	XSD	JastAd d, Silver	MOF	Ecore, EMOF	ERD	GOPPR	CROM	RDFS OWL
M 2	Gramma r of a language	Gramm ar with line delimit ers	csv- head er	Relation al Schema	NF2- Schema	XML Schema , e.g. xhtml	Specific RAG	UML- CD, -SC, OCL	UML, many others	CDI F- lang uage s	UML, many others	CROM	HTML XML MOF UML DSL
M 1	String, Program	Text in lines	csv Table	Relation s	NF2-tree relation	XML- Docum ents	Link- Syntax- Trees	Classes, Progra ms	Classes, Progra ms	CDI F- Mod els	Classes, Programs	CROM models	Facts (T- Box)
M 0	Objects	Sequenc es of lines	Sequenc es of rows	Sets of tuples	trees	dynam ic semanti cs in browse r		Object nets	Hierarch ical graphs	Objec t nets	Object nets	Context- Object-Role Nets	A-Box (RDF- Graphs)

today

# From Syntax Trees to Syntax Graphs

- ▶ In the TS Graphware, the secondary relations of link trees become *primary relations*, i.e., we treat *real graphs*
- ▶ Abstract syntax trees (AST) change to Abstract Syntax Graphs (ASG)
- ▶ Attributed link trees (ALT) change to Attributed Program Graphs (APG)

# Flat and Deep Model and Code Analysis

- ▶ DQL answer questions about the materials in a repository or in a stream
  - Analytics for one document alone (metrics, “Business Intelligence”)
  - Filtering of a stream
  - Combining input streams

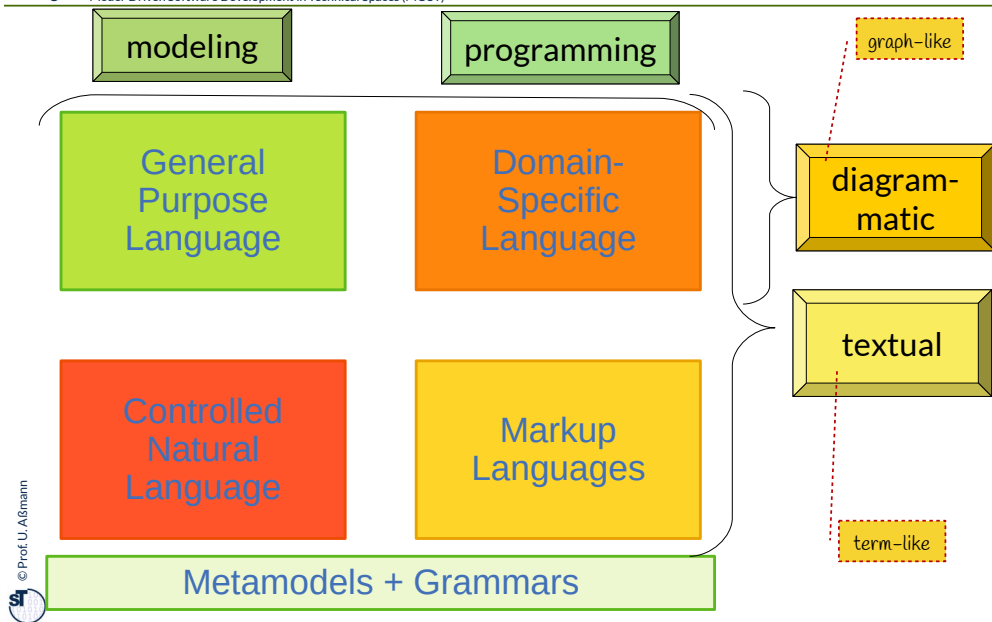
CQL do the same for programs and models:

- ▶ **Flat model analysis** asks questions on
  - the direct context of a model element (context-free queries, pattern matching)
  - the global knowledge about a model element
  - **Software metrics:** counting objects, relationships, dependencies
  - **Inter-model dependencies** between models in a megamodel
- ▶ **Deep model analysis** (value flow analysis, data-flow analysis, inter-procedural analysis, inter-component analysis) respects the main structure of a model and asks the question
  - whether certain parts of a model are reachable from each other (connected)
  - what is the context of a model element in a structured environment (abstract syntax tree, control flow graph, value flow graph, dependency graph)
  - where do attributes flow (in an attribution)

## Q16: Languages in Software Factories

8

Model-Driven Software Development in Technical Spaces (MOST)







## 40.1 DDL in the Graph-Based Technical Spaces

## 40.1.1 Technical Space RelationWare with DDL Relational Schema

Relational Algebra works with *typed relations*

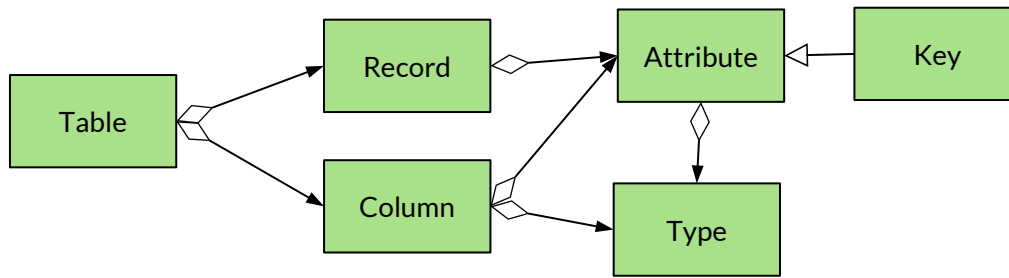


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# Technical Space Relational Algebra mit Metalanguage Relational Schema

11 Model-Driven Software Development in Technical Spaces (MOST)

- ▶ Relational Algebra (Codd) works on tables of tuples with attributes
  - See courses on databases



Relational Schema  
Metamodel

Key	FirstName	Surname	Street	Town
@1	Uwe	Aßmann	Bakerstreet 5	New York
@2	Frank	Miller	Northstreet 9	Pittsburgh
@3	Mary	Baker	Magdalenstr eet	Oxford

RS muss metamodelliert werden.  
RS ist Teil von ERD von eMOF

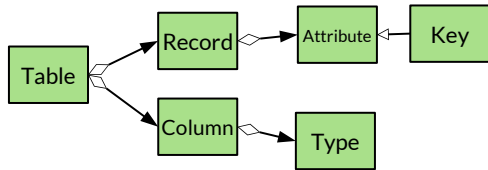
## 40.1.2 Excursion: Textual Notation for Graphs

Relational Algebra works with *typed relations*



# Textual Notation for Graphs and Diagrams

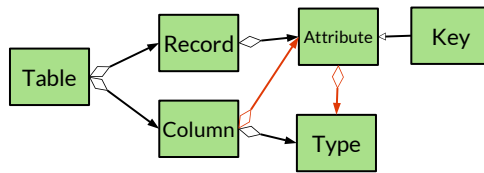
- ▶ A hierarchic structure (tree or term) can be expressed in *term-like syntax*:



```
// without edges
Table [ Record [ Attribute [ Key ] ],
      Column [ Type ]
]
```

```
// with edges
Table [ has [ Record [ has [ Attribute [
    subclasses [ Key ] ] ] ] ],
      has [ Column [ has [ Type ] ] ]
]
```

- ▶ A real graph or diagram can be split into terms and joined by conjunction (*spanning tree decomposition*)

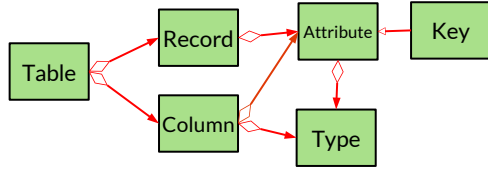


```
// without edges
Table [ Record [ Attribute [ Key ] ],
      Column [ Type ] ]
AND Column [ Attribute ]
AND Attribute [ Type ]
```

```
// with edges
Table [ has [ Record [ has [ Attribute [
    subclasses [ Key ] ] ] ] ],
      has [ Column [ has [ Type ] ] ] ]
AND Column [ has [ Attribute ] ]
AND Attribute [ has [ Type ] ]
```

# Textual Notation for Graphs and Diagrams


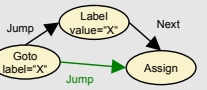
- ▶ A real graph or diagram can be split into flat terms (triples) and joined by conjunction (**edge decomposition, triple decomposition**)
- ▶ Most query and transformation languages in Graphware use either
  - spanning tree decomposition
  - edge decomposition.
- ▶ Ontology languages (such as OWL and RDFS) use triple decomposition



```
// with edges  
has[Table, Record] AND has[Table, Column]  
AND has[Record, Attribute] AND subclasses[Attribute, Key]  
AND has[Column [Type]] AND has[Column, Attribute]  
AND has[Attribute, Type]
```

# Different Notations for Node-Edge Patterns in Edge Decomposition

- ▶ In edge decomposition of query graphs, for notation of edges (and predicates), textual as well as graphical notations exist

	Datalog Prolog	Graphic (Optimix, EARS)	Textual graphics (TgreQL, GrGen)	Juxtaposition	Object-oriented (.QL)
edges	$e(N,M)$		$-N-e-M->$ $N-e->M$	$N e M$	$N.e(M)$
recursion	$r(N,M) :-$ $e(N,Z),$ $r(Z,M)$		$N-e^*->M$	$N e^* M$	$N.e^*(M)$

## 40.1.2 Technical Space ER-Ware with DDL Entity-Relationship-Diagrams (ERD)

A Simple DDL/CDL with Mapping to the  
Relational Algebra

Relations and Entities (without inheritance)

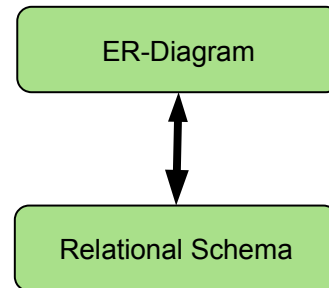


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## Modeling with Entity-Relationship-Diagrams (ERD)

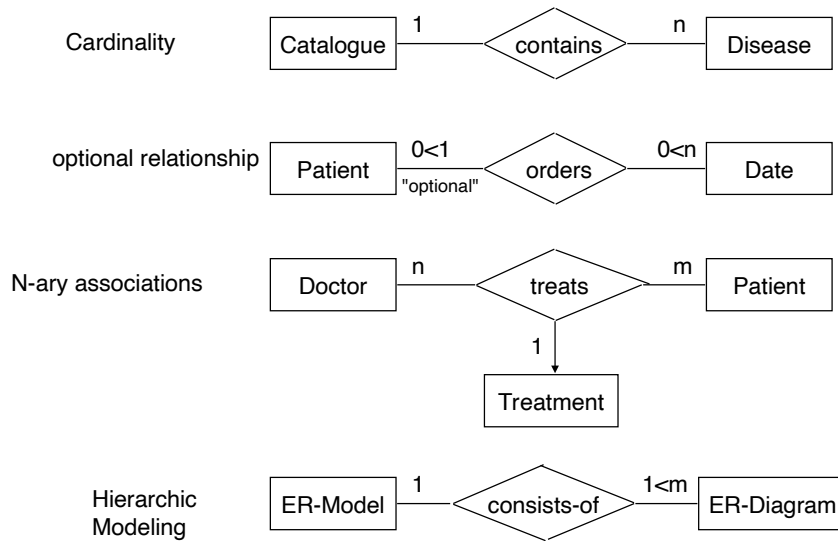
- ▶ ERD can be mapped easily to relational schema (with an invertible 1:n-mapping, **ER-RS-mapping**)
  - Entities form special relations with “identifier” (key, surrogate)
  - ER-diagrams can be stored easily in databases (simple persistence)
- ▶ ERD is often used as CDL in larger integrated development environments (simple persistence of code and models)



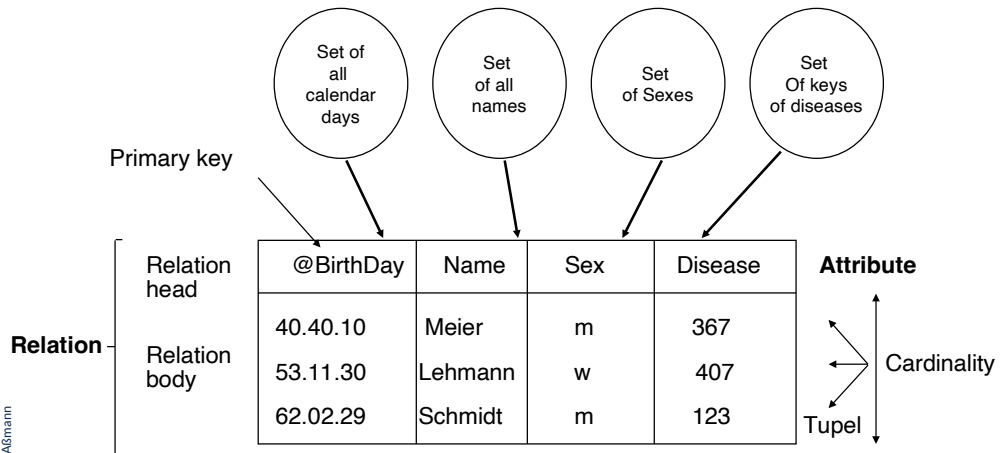
Diskussion  
Abbildung Vererbung  
Hibernate??

## ERD-Relationships in Chen-Notation (unlike UML)

- ▶ All "entities" (classes) are represented as "entity-"tables



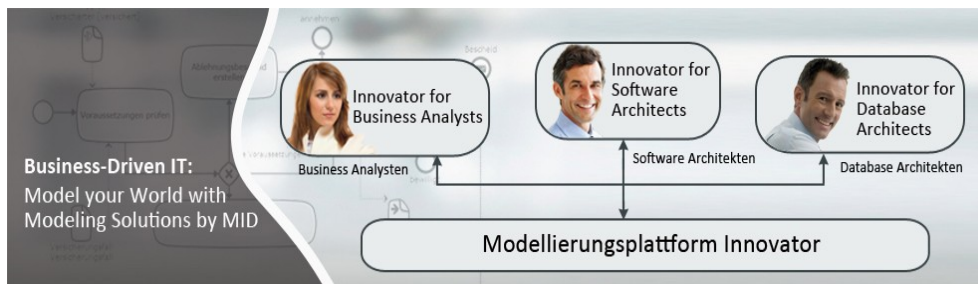
## Mapping of Entity Type "Patient" to the Relational Schema



Geburtstag ist kein Primärschlüssel..

## Importance of ERD

- ▶ ERD is the “better” relational schema, because it treats objects (entities)
  - Often used for data dictionaries in information systems
- ▶ ERD, however, does not support inheritance
  - Applications can easier be verified, e.g., for embedded or safety-critical systems
- ▶ Typical Tool: MID Innovator for database architects:



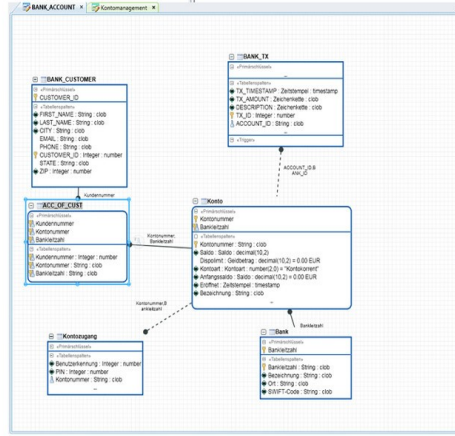
<http://www.mid.de/index.php?id=541>

[http://www.mid.de/uploads/pics/Banner\\_Modellierungsplattform\\_03.jpg](http://www.mid.de/uploads/pics/Banner_Modellierungsplattform_03.jpg)

# Mapping ERD to RS in MID

<http://www.mid.de/typo3temp/pics/f0df65b8a2.jpg>

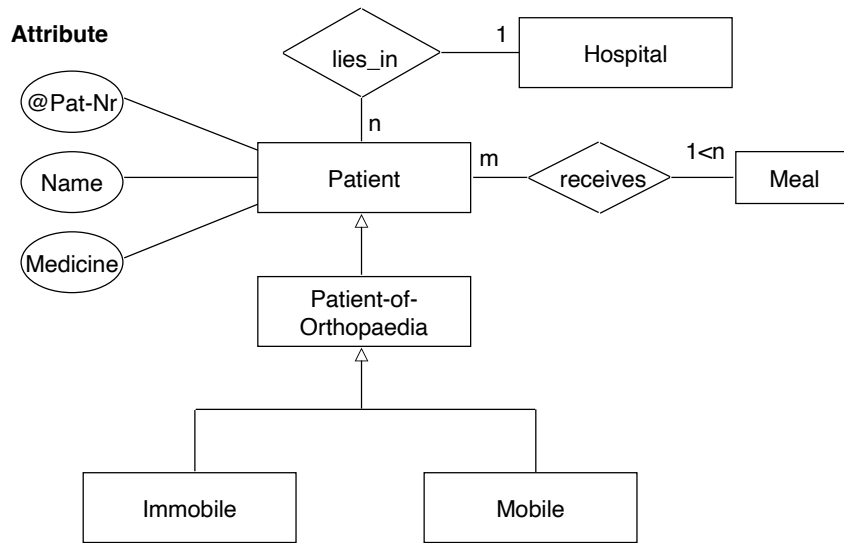
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The screenshot shows a 'Mapping' tool window with the following details:

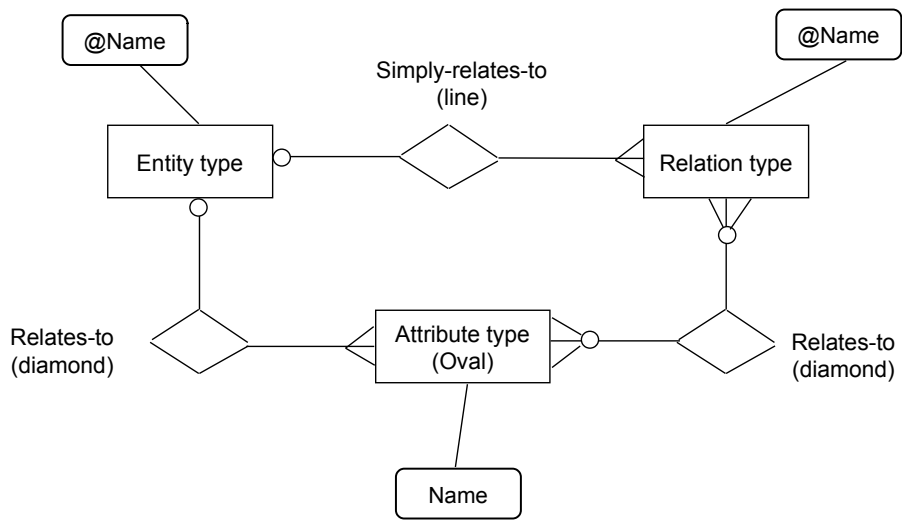
- Source Model:** Oracle-DB-Schema aus ER-Modell ableiten
- Target Model:** Oracle
- Source Model Elements (Auswahlmodelle):** Überweisung, Transaktion, Kunde, Kontozugang, Konto, gehört, Dauerauftrag, Buchung, Barauszahlung, Bank, Kontoauszug.
- Target Model Elements (Zielmodell):** Oracle, DefaultCatalog, DefaultSchema, Bank, BANK\_SHEDL\_TX, TX\_INTERVAL, TX\_BEGINDATE, TX\_ID, erKey, BANK\_TX, Buchung, dbKey, erKey, erKey, Buchung, BANK\_CUSTOMER, ACC\_OF\_CUST, Kontoauszug.

## Extended ERD (EERD) Uses Inheritance Example: Patient Record



Schlechtes Beispiel, da mit Vererbung

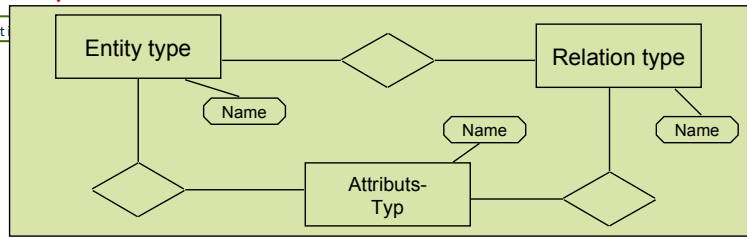
# The Metamodel of ERD in ERD (lifted ERD Metamodel)



# Metahierarchy with ERD as Metalanguage (lifted metamodel)

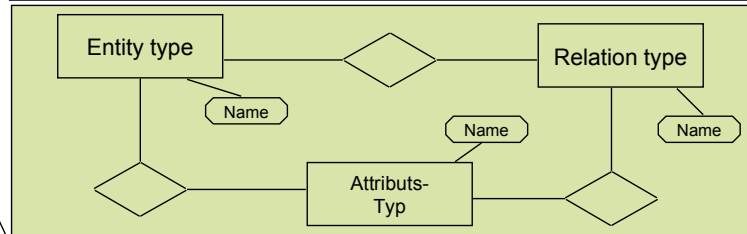
**M3**

**Metametamodel**



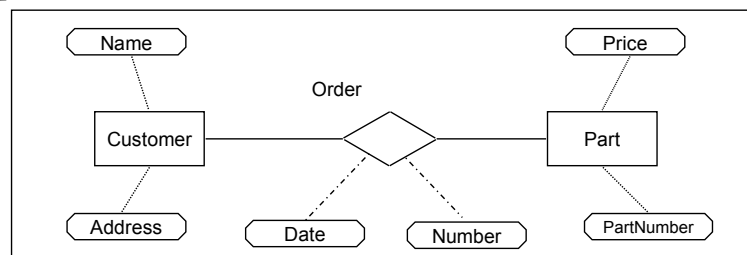
**M2**

**Metamodels**



**M1**

**Models**



Sollte mit zweiter Folie auf M3 kontrastiert werden

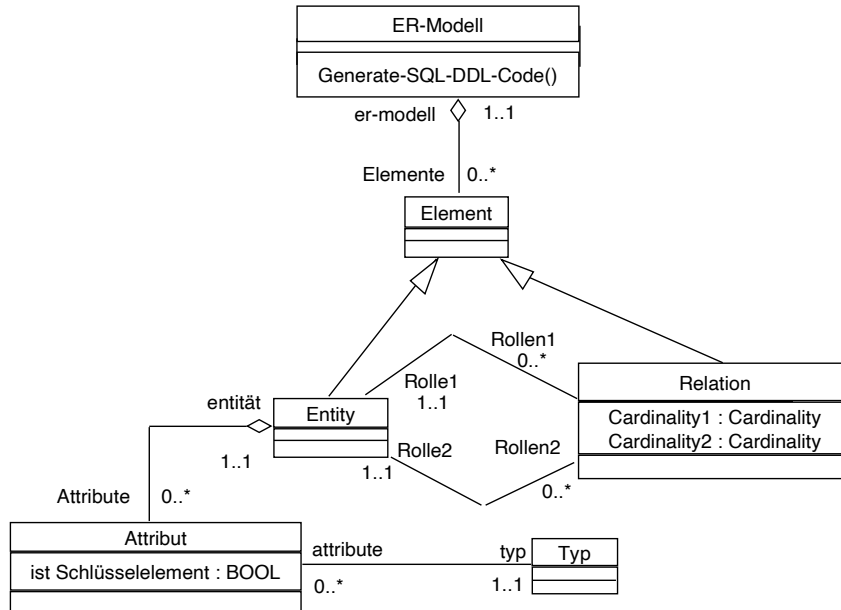
Was bedeutet Kreis?

M3/M2 embedding erklären für DDL



# MOF is ERD with Inheritance

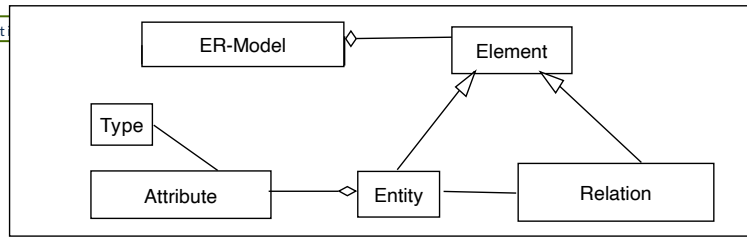
## Meta-Modell of Entity-Relationship-Diagramms (in MOF)



# Metahierarchy with MOF as Metalanguage (non-lifted)

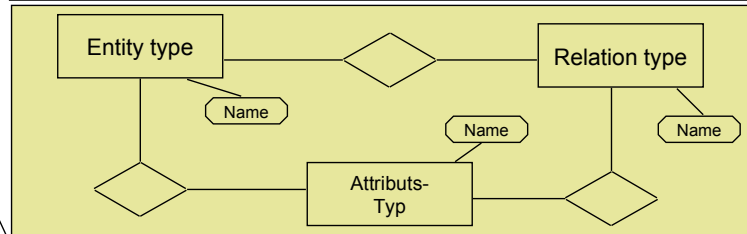
**M3**

**Metametamodel**



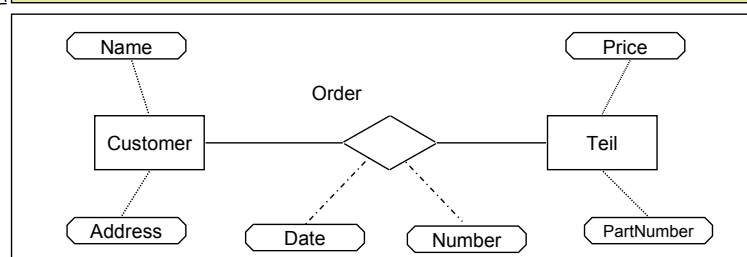
**M2**

**Metamodels**



**M1**

**Models**



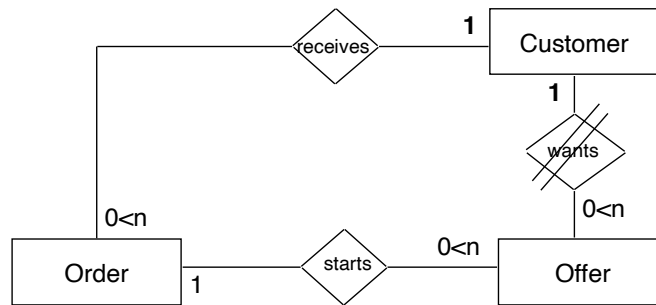
Sollte mit zweiter Folie auf M3 kontrastiert werden

Was bedeutet Kreis?

M3/M2 embedding erklären für DDL

## Consistency Constraints in ERD Models

- ▶ An ERD can contain integrity constraints (consistency constraints)
- ▶ Ex.: **Cycle-freedom constraint**: Check: find cycles in the graph of a ER diagram
- ▶ Correct by
  - cutting a cycle at the least important position (human intervention)
  - Finding a spanning tree and cutting all other edges
- ▶ Instead of cutting, edges can be made secondary links (then we have link trees)



Viel zu detailliert

## Other Consistency Constraints of ER-Models

- ▶ **Range checks** for attributes
- ▶ **Key dependencies (functional dependencies):**
  - Uniqueness of attribute values: An attribute K of a relation R is a key candidate, if only one tuple has the same value of K
  - Key minimality: Is the attribute K compound, no component can be removed to loose the key condition.
  - Primary key serves for identification of a tuple (“entity check”)
  - Secondary keys: other keys
  - Foreign key reference (primary key reference): A foreign key (link) is referencing a tuple in another relation by its primary key
- ▶ **Referential Integrity**
  - The model does not contain undefined foreign keys (links)
  - i.e., all names (links) can be resolved by name analysis

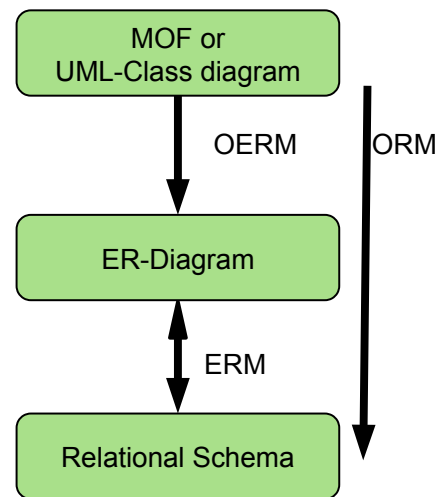
Allgemeiner fassen, mehr Beispiele.

## 40.1.3 MOF as Extended ERD



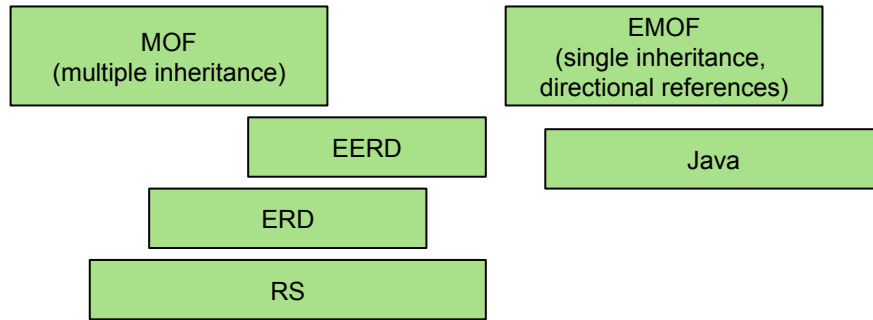
# Data Modeling for Information Systems (Object-Relational Mapping, ORM) with UML-CD, ERD and RS

- ▶ For persistence, objects should be stored with an object-relational mapping to a database (OR-Mapping)
- ▶ OERM-Mapping of class diagrams to ERD is (unfortunately) indeterministic
  - Inheritance mapping
  - Identification of keys (primary, secondary, foreign)
  - Resolution of multiple inheritance by copying
  - Cannot be inverted automatically
- ▶ Between ERD und RS exists a *deterministic, bidirectional* mapping (ER-Mapping) by which the data models can be synchronized (restored without information loss)



# The Difference of ERD, MOF and EMOF

- ▶ MOF extends ERD with multiple inheritance and method signatures
- ▶ However, MOF must be mapped down to Java
  - Inheritance
  - Bidirectional associations
- ▶ EMOF has only directed references, no bidirectional associations
  - Only simple inheritance
- ▶ EMOF can directly be mapped down to Java, C++, or C#



## 40.2 Flat Model Analysis with Graph Query Languages (Graph QL)

DQL – Data Query Languages

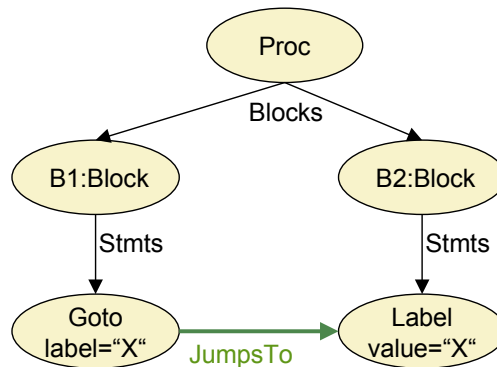
CQL – Code Query Languages



# Graph Pattern Matching of Non-Tree Patterns

- ▶ **Graph pattern matching** works by mapping a graph pattern (graphlet) to the manipulated graph.
- ▶ Ex.: Linking gotos and Block-entry statements to build up the control-flow graph

```
-- Datalog notation (edge decomposition):  
JumpsTo(Goto,Label) :-  
  Blocks(Proc,B1:Block),  
  Blocks(Proc,B2:Block),  
  Stmts(B1,Goto),Stmts(B2,Label),  
  Goto.label==X, Label.value==X.  
  
-- Optimix notation with if-then rules  
(edge decomposition):  
If  Blocks(Proc,B1:Block),  
    Blocks(Proc,B2:Block),  
    Stmts(B1,Goto),Stmts(B2,Label),  
    Goto.label==X, Label.value==X  
Then  
  JumpsTo(Goto,Label).  
  
- regular expression notation (TGreQL):  
JumpsTo := Proc.Blocks.Stmts.Goto.label(X)  
        AND Prod.Blocks.Stmts.Label.value(X)
```



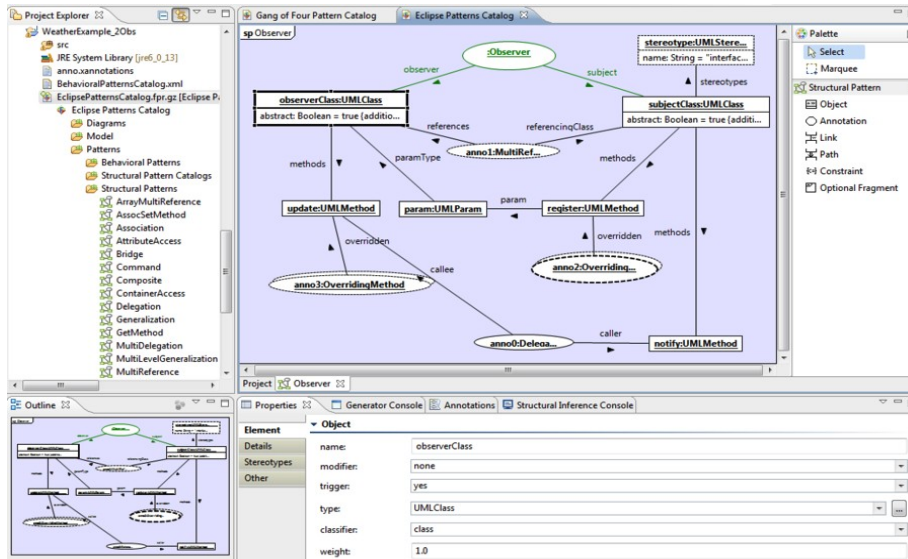


## 41.1. Introduction to Diagrammatic Storyboard Rule Notation for Graph Rewriting

Coloring for rules originally introduced by Fujaba  
[www.fujaba.de](http://www.fujaba.de) (tool now unsupported)



- ▶ Fujaba is a MetaCASE-tool based on GRS with home-grown metalanguage and metamodel
- ▶ Basic technology: graph pattern matching and rewriting



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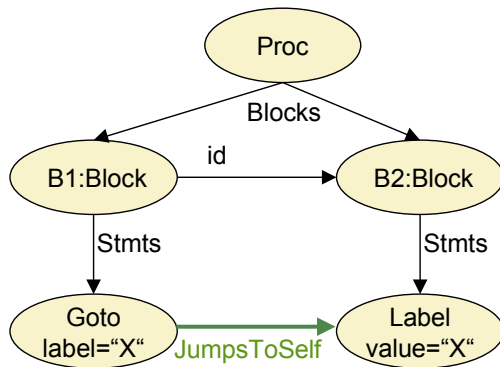
<http://www.fujaba.de/typo3temp/pics/604c5c6c9e.png>

# Pattern Matching of Non-Tree Patterns

- ▶ **Flat analysis** does not interpret the program while analysing
- ▶ **Deep analysis** interprets the primary graph (ASG) to use the program semantics

- ▶ Query: **Which blocks jump to themselves?**

```
-- Datalog notation (edge decomposition):  
JumpsToSelf(Goto, Label) :-  
  Blocks(Proc, B1:Block),  
  Blocks(Proc, B2:Block), id(B1, B2)  
  Stmts(B1, Goto), Stmts(B2, Label),  
  Goto.label==X, Label.value==X.
```

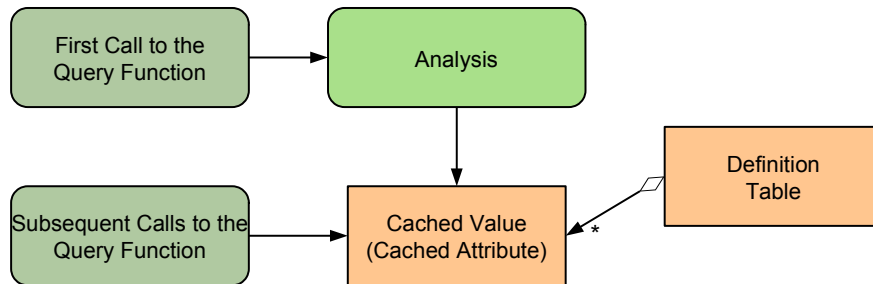


From the metamodel, we can define **access**, **helper**, **query** and **attribution functions**, functions to access, query **attributes** or **neighbors**:

- ▶ **(Local) Attribute access functions:**
  - ModelElement.hasName()
  - ModelElement.getDeclaringType()
- ▶ **Neighbor access functions** (via references):
  - Class.getPackage(): for neighbor Package
  - Class.getUpperClass(): get the direct upper class
  - Class.getDeclaresMethod(): for contained Method
- ▶ **Query functions** looking up information in the abstract syntax graph (ASG) or model:
  - Expr.getUsedTypes(): search all types which are used in Expr (type analysis, type resolution)
  - Name.getType(): search the type object to the Name
  - Name.getMeaning(): search the definition of the Name
  - Stmt.getProcedure(): search out to find the procedure of the Stmt
- ▶ **Pattern match functions** assemble all matching redexes of a pattern
  - findRedexes (Pattern) → Redexes

## Name and Type Analysis: Caching a Query Function

- ▶ Some values of query functions change never, once they have been determined
  - The values can be cached
- ▶ **Attribute caching** is a mechanism to cache semantic attributes in an ASG or model for faster access
- ▶ A **definition table (often called symbol table)** is a set of cached attributes.





## 40.2.1 QL and CodeQL – Relational Queries on Source Code in Technical Space Java

QL uses edge decomposition (Datalog style) to express graph queries

Courtesy to Florian Heidenreich and  
<http://semml.com> (Semml now part of Github)



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- ▶ **QL** is an object-oriented query language in the spirit of SQL and Datalog
  - Developed in the group of Prof. Oege de Moor (Oxford)
  - Marketed by Semmler.com
  - In 2019 bought by github
- ▶ Queries, metrics, visualizations are supported
  - Repositories with Java and Objective-C code
  - Works also now on C/C++
- ▶ Metamodel is EMOF-like (single inheritance, references)
  - Classes, Methods, Blocks are interpreted as basic **sets** of objects, **relational tables** (sets of tuples over member entries), resp. **Predicates** (telling whether a tuple exists)
    - **from** Class *c*, Methods
  - Definition and use of access functions:
    - Class.getDeclaredMethod(): for neighbor Method
    - Class.getPackage(): for neighbor Package
    - ModelElement.hasName(): get the Name
    - ModelElement.getDeclaringType(): get the Type



## Query form: Extended Where- Select Clauses

- ▶ Expressions like in Xcerpt and SQL:
  - FROM <classes> WHERE <conditions> SELECT <variables>

**FROM** ..base sets..

**WHERE**

..and/or/not predicate list..

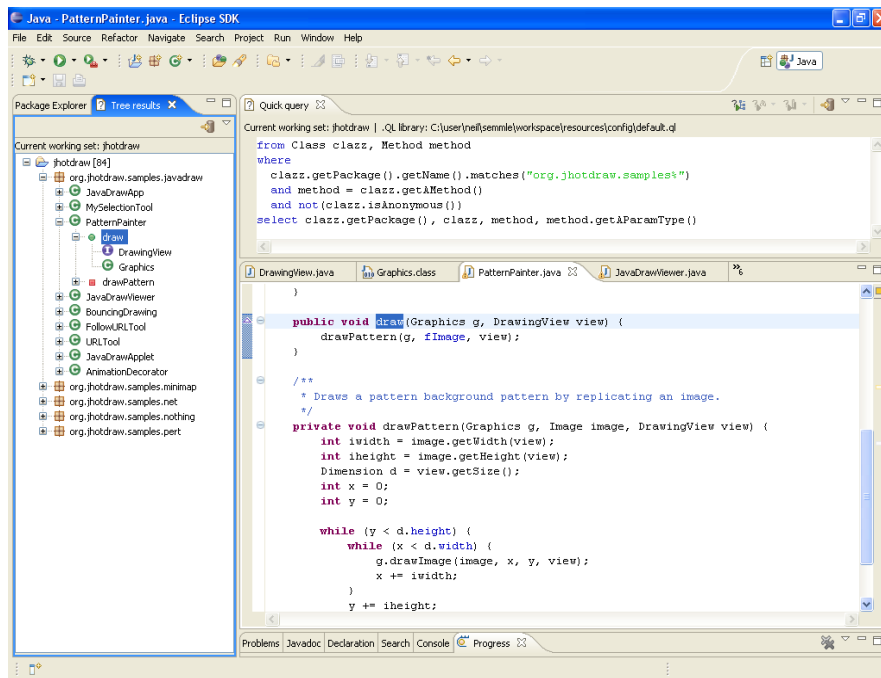
..call of helper functions..

..call of predicates..

..check on equalities, inequalities..

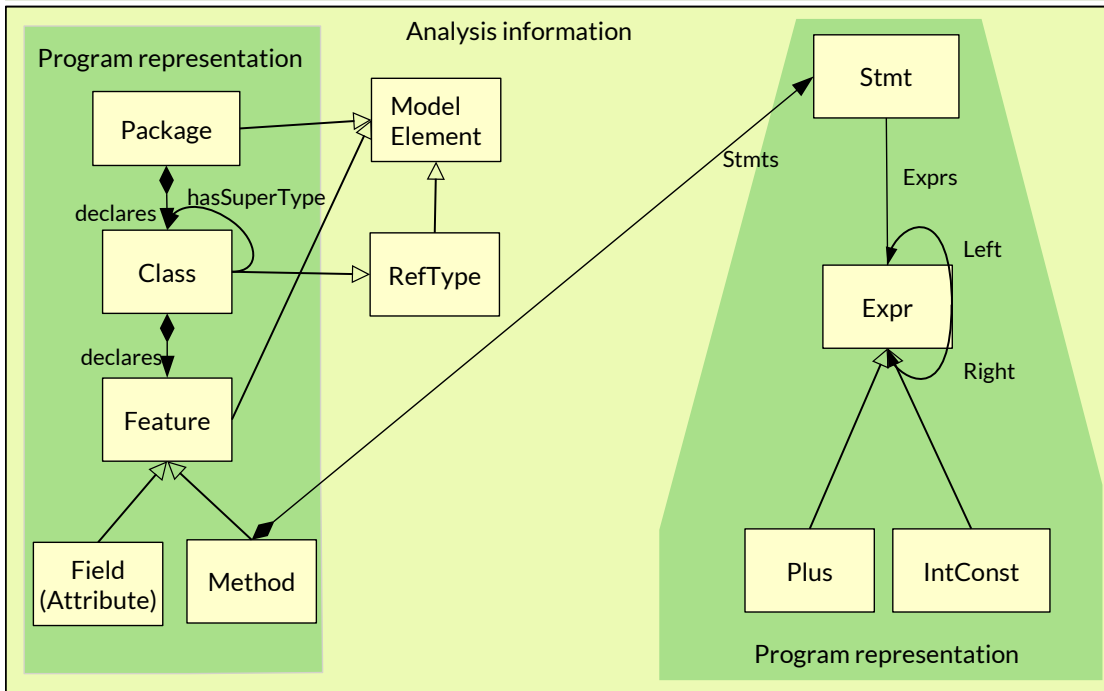
**SELECT** variable list

# Code Display





# A Simple Model (Schema) of Semmlle-Java-DDL in EMOF



- ▶ Query examples:
  - Select Statements on classes, methods, statements, expressions
- ▶ Language features:
  - Queries embedded in classes, shareable with inheritance
  - User defined query classes
  - Local Variables in queries
  - Non-deterministic methods returning *sets* and *streams*
  - Casts
  - Chaining
  - Lifting-queries
- ▶ Metric examples:
  - Aggregation functions
  - SLOC metrics
  - #Methods
- ▶



## Select Statements (1)

- ▶ The where-clause uses edge decomposition of a query graph
- ▶ Example:
- ▶ Find all classes `c` implementing `compareTo`, but do not overwrite `equals`
- ▶ Find their packages
- ▶ Return tuples of package and class

```
from Class c
where
  c.declaresMethod("compareTo")
  and not (c.declaresMethod("equals"))
select
  c.getPackage(), c
```



## Select Statements (2)

- ▶ Find all **main**-methods declared in a package ending with „demo“
- ▶ Return tuples (package, declaring type, method)
- ▶ Also called **pattern matching**

```
from Method m
where
  m.hasName("main")
  and m.getDeclaringType().getPackage().getName().matches("%demo")
select
  m.getDeclaringType().getPackage(),
  m.getDeclaringType(),
  m
```



## Definition of New Functions and Predicates

- ▶ Definition of new **query functions** by declaring query functions/methods in a class (note: this is similar to attributions in JastAdd)
  - Remark: Methods may be indeterministic, i.e., return collections of objects

```
class Classinfo {
  Method ndMethod(Class c) {
    c.declaresMethod("sumUpBill")
  }
}
```

- ▶ Definition of new **predicates** as methods in a class, using a domain-specific language language extension of Java
- ▶ Testing on or-conditions:

```
predicate isJDKMethod (Method m) {
  m.hasName("equals")
  or m.hasName("hashCode")
  or m.hasName("toString")
  or m.hasName("clone")
}
```



## Definition of New Predicates

- ▶ Use of Kleene Star for transitive closure on predicates/edges
- ▶ The Kleene star expands the relation *transitively* (*transitive closure*)
- ▶ Here, `hasSupertype` is deeply searched:

```
predicate upperClass(RefType down, RefType up) {  
    down.hasSupertype (up)  
}
```

- ▶ Reachability in control-flow graph over statements

```
predicate control flowReach(Stmt rst, Stmt reachable) {  
    rst.successor (reachable)  
}
```

## Definition of New Predicates

- ▶ Complicated, composed path expressions become possible
- ▶ Query: *Check for a middle class in the inheritance hierarchy:*

```
predicate inTheMiddle(RefType down, RefType middle, RefType up) {  
  down.hasSupertype (middle) and  
  middle.hasSupertype (up)  
}
```



## Local Variables in Queries

Query: Find all methods calling `System.exit(...)`

Sysexit is a local variable

```
from Method m, Method sysexit, Class system
where
  system.hasQualifiedName("java.lang", "System")
  and sysexit.hasName("exit")
  and sysexit.getDeclaringType() = system
  and m.getACall() = sysexit
select m
```



## The Use of Non-deterministic Methods

- ▶ Query: **Synthesize a call graph between the methods of two packages**
  - Call graph is returned as a set of tuples of (caller, callee)
- ▶ getARefType and getACallable are indeterministic, i.e., return collections of objects

```
from Package caller, Package callee
where caller.getARefType().getACallable().calls(
    callee.getARefType().getACallable())
    and caller.fromSource()
    and callee.fromSource()
    and caller != callee
select caller, callee
```

## Chaining (Multiple Source - Multiple Target Graph Reachability Problem, MSMT)

- ▶ MSMT problems connect a set of source nodes with a set of target nodes (reachability)

Query: Find all Pairs (s,t) such that

- ▶ t is a direct superclass of s
- ▶ s is superclass of `org.jfree.data.gantt.TaskSeriesCollection`
- ▶ t is superclass of s
- ▶ and t is not `java.lang.Object`

```
from RefType tsc, RefType s, RefType t
where
  tsc.hasQualifiedName("org.jfree.data.gantt", "TaskSeriesCollection")
  and s.hasSubtype*(tsc)
  and t.hasSubtype(s)
  and not(t.hasName("java.lang.Object"))
select s,t
```



## QL-Query Classes (Dynamic Classes/Sets)

- ▶ **Query classes** in QL are sets described by special predicates and nested other predicates
  - They define “synthetic” objects and “truths” about the model
  - Their constructors define restrictions of metaclasses

```
// de nition of a query class as subclass of a metaclass
```

```
class VisibleInstanceField extends Field {
```

```
  VisibleInstanceField() {
```

```
    not (this.hasModi er("private")) and
```

```
    not (this.hasModi er("static"))
```

```
  }
```

```
  predicate readExternally() {
```

```
    exists (FieldRead fr |
```

```
      fr.getField()==this and
```

```
      fr.getSite().getDeclaringType()
```

```
        != this.getDeclaringType()
```

```
    }
```

```
  }
```

© Prof. U. Alsmann

```
// use of a query class
```

```
from VisibleInstanceField vif
```

```
where vif.fromSource() and not
```

```
  (vif.readExternally())
```

```
select vif.getDeclaringType().getPackage(),
```

```
  vif.getDeclaringType(),
```

```
  vif
```



## 40.2.2 Metrics with QL



# Aggregation Functions for Computing Metrics

- ▶ Compute the average number of methods per type and package
  - Other aggregation functions: count, sum, max, min, avg
- ▶ Employs „Eindhoven Quantifier Notation“ (Dijkstra et al.)
  - $C \mid \langle \text{predicate} \rangle$
- ▶ Query: „Compute the average number of methods in all type  $c$  of a package  $p$ “

```
from Package p
where p.fromSource()
select p, avg(RefType c |
             c.getPackage() = p |
             c.getNumberOfMethods())
```





# Aggregation Functions for Computing SLOC Metrics

- ▶ Query: “Calculate a SLOC metrics on package “Billing” in the current compilation unit”
- ▶ Grammar rules:
- ▶ Aggr ::= aggregationFunction '('  
    localvars // FROM  
    '|' condition // WHERE  
    '|' aggregatedValue ')' // SELECT
- ▶ AggregationFunction ::= 'sum' | 'count' | 'avg' | 'max' | 'min'

```
from Package pkg
where pkg.hasName("Billing")
select sum(CompilationUnit comp | //FROM
         comp.getPackage()=pkg | // WHERE
         comp.getNumberOfLines()) // SELECT
```



# Statistics (Metrics) Uses Aggregation Functions

The screenshot shows the Eclipse IDE interface. On the left is the Package Explorer showing a project structure with various Java classes. The central editor displays a SQL query using aggregation functions. On the right, a window titled 'Average Method/Constructor Fan-In (jhotdraw)' displays a bar chart. The chart shows the average fan-in for various packages, with the highest values for 'org.jhotdraw.util' and 'org.jhotdraw.samples.net'.

```
from Package p, float average
where p.fromSource()
and average = avg(Class c, Callable member
| c.getPackage() = p and
member.getDeclaringType() = c
| member.getFanIn())
select p, average order by average desc
```

**Average Method/Constructor Fan-In (jhotdraw)**

Package	Average Fan-In
org.jhotdraw.applet	1.0
org.jhotdraw.application	1.5
org.jhotdraw.contrib	1.0
org.jhotdraw.contrib.dnd	1.0
org.jhotdraw.contrib.html	0.8
org.jhotdraw.contrib.zoom	0.8
org.jhotdraw.figures	1.2
org.jhotdraw.framework	0.5
org.jhotdraw.samples.javadraw	0.2
org.jhotdraw.samples.minimap	0.2
org.jhotdraw.samples.net	2.5
org.jhotdraw.samples.nothing	0.2
org.jhotdraw.samples.pert	0.8
org.jhotdraw.standard	0.2
org.jhotdraw.util	2.5
org.jhotdraw.util.collections.jdk11	0.2
org.jhotdraw.util.collections.jdk12	0.2

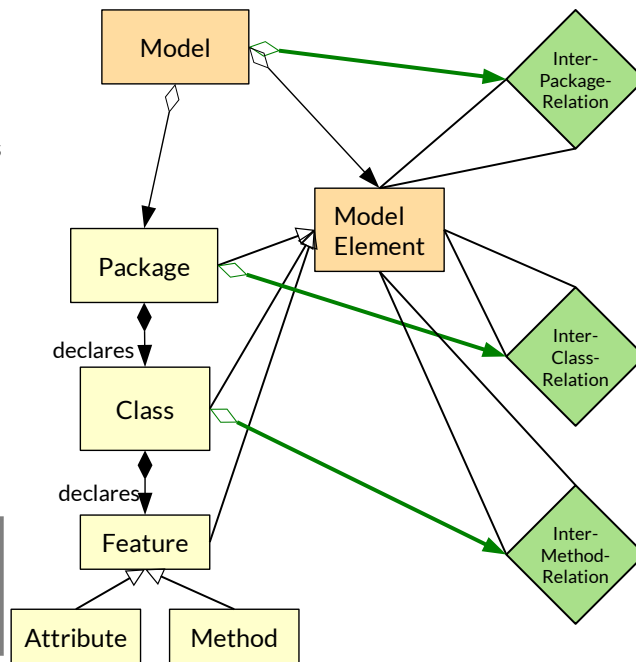


## 40.3. Lifting Information Up the Containment Hierarchy

# Block Containment Structure (Scope Structure in the ASG) of a Model

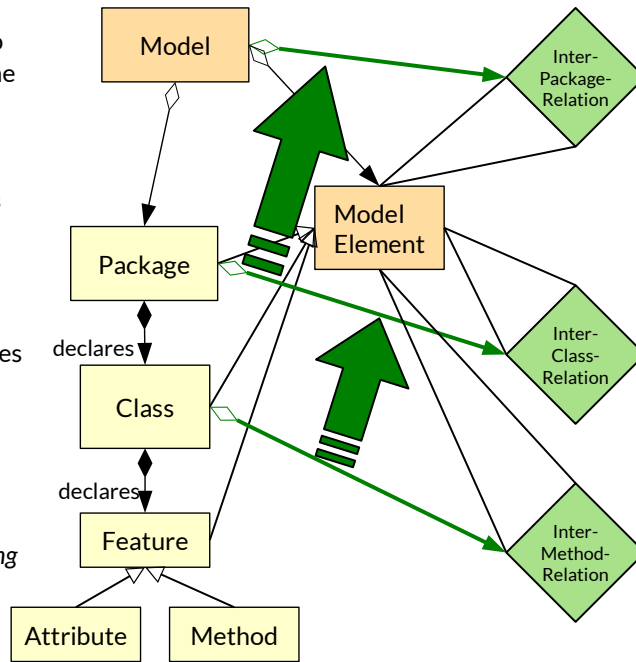
- ▶ Languages are block-structured, i.e., live in a **containment hierarchy**.
- ▶ A model has **model elements**
- ▶ A class has **inter-method relationship** (e.g., the call graph)
- ▶ A package has **inter-class relationships** between these model elements
- ▶ A model has **inter-package relationships** between these model elements

**A macromodel builds on graphs, at least on link trees, no longer on trees**



# Lifting Information Along the Block Containment Structure (Scope Structure of the ASG) by Synthesized Attribution

- ▶ **Dependency lifting** means to lift information **up** in along the containment hierarchy by a *synthesized attribution*
  - from an inter-method relationship to a inter-class relationship
  - from an inter-class relationship to a inter-package relationship
- ▶ Dependency lifting propagates information **up** the abstract syntax tree and the containment tree
- ▶ Dependency lifting is an important process to *summarize dependencies among siblings in containment hierarchies in models*



# Dependency Lifting Information Along the Block Containment Structure

- ▶ **Dependency lifting** lifts dependency information up the containment structure in a model, thereby summarizing the dependencies at the level of the model
- ▶ **Dependency lifting queries** are defined on an enclosed type
- ▶ result is an implicitly defined default return parameter of a query

```
// Lifting a pair of method dependencies
// on a pair of classes
// getDependentClass() is a synthesized
// attribution of Class.Method
class Method {
  Class getDependentClass() {
    exists (Method m |
      depends(this.getClass(),m)
      and result = m.getClass()
    )
  }
}
```



```
// Lifting a pair of class dependencies to
// a pair of packages
// getDependentPackage() is a synthesized
// attribution of Package.Class
class Class {
  Package getDependentPackage() {
    exists (Class cl |
      depends(this.getPackage(),cl)
      and result = cl.getPackage()
    )
  }
}
```

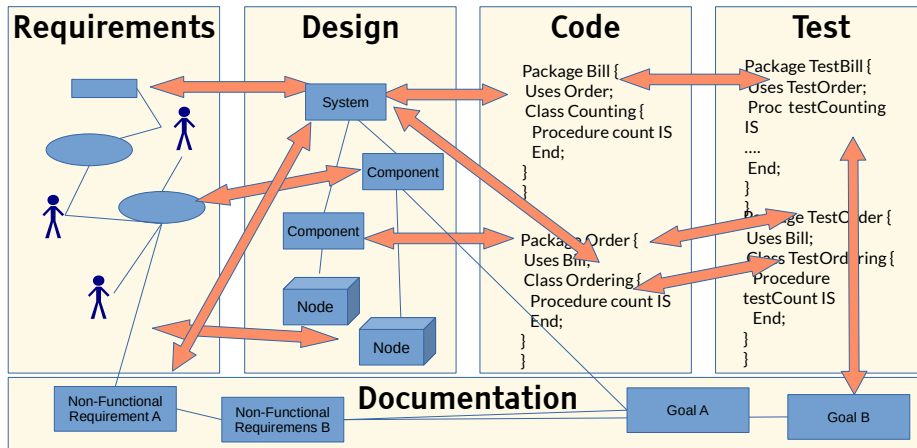


## 40.4 Macromodel Dependency Analysis

- Remember: A **macromodel** is a multimodel with consistent dependencies

## Q12: The ReDoDeCT Problem and its Macromodel

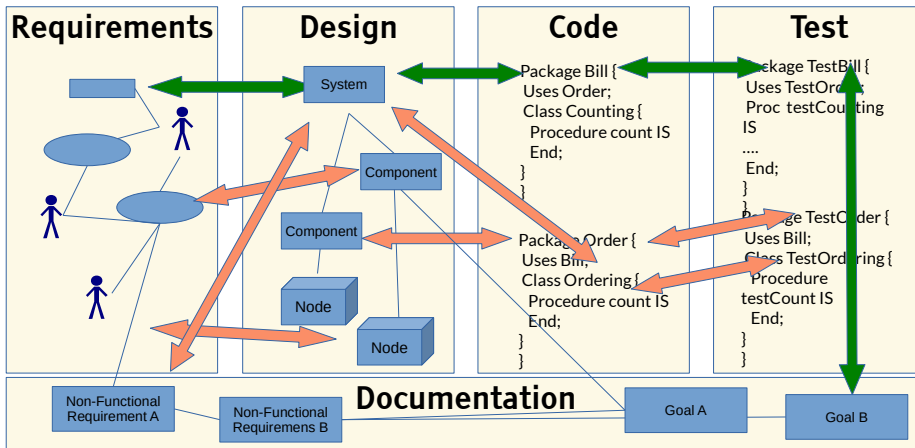
- ▶ The **ReDoDeCT problem** is the problem how requirements, documentation, design, code, and tests are related ( $\rightarrow$  V model)
- ▶ Mappings between the Requirements model, Documentation files, Design model, Code, Test cases
- ▶ A **ReDoDeCT macromodel** has maintained mappings between all 5 models





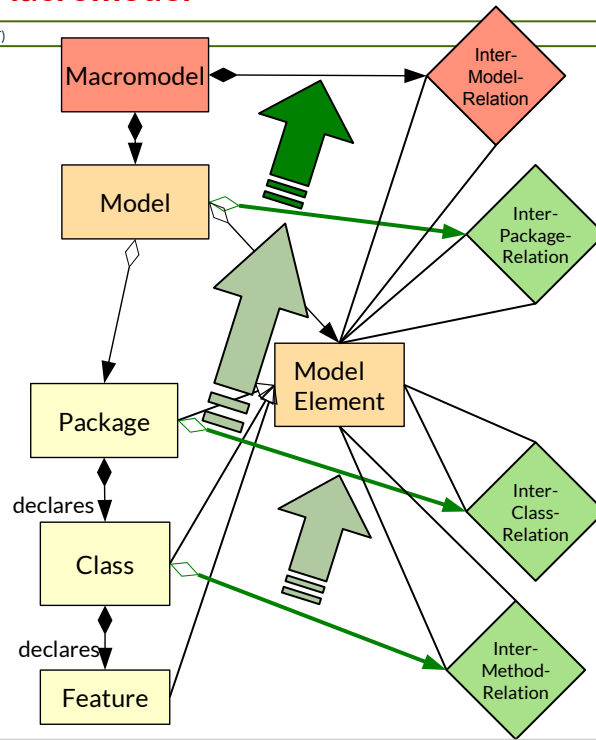
# Inter-Model Relationships in The ReDoDeCT Macromodel

- ▶ An **inter-model relationship** is a relationship between model elements of different models (usually link or graph relationship)
  - Here: expresses mapping between the Requirements model, Design model, Code, Test cases
- ▶ The ReDoDeCT macromodel relies on **inter-model relationships** between all 5 models



# Lifting Information Along the Block Containment Structure Between Models in the Macromodel

- ▶ **Macromodel-Dependency Lifting** means to lift information **up** in along the containment hierarchy **from between the packages of a model to between the models of the macromodel**
  - from an intra-model relationships to a inter-model relationship
- ▶ Megamodel-Dependency- Lifting propagates information **up** into the megamodel
- ▶ Megamodel-Dependency- Lifting is an important process to *summarize dependencies among models*
- ▶ Result: a **macromodel**



## Cultimodel Dependency Lifting in Semmlé QL

- ▶ The lifting procedure also works for lifting package dependencies within a model to model dependencies.
  - Consider models as “normal” objects in the repository
  - Formulate queries about model-element relationships and lift them to model relationships

```
// Lifting a pair of class dependencies to
// a pair of packages
class Class {
  Package getDependentPackage() {
    exists (Class cl |
      depends(this.getPackage(),cl)
      and result = cl.getPackage()
    )
    and result != this
  }
}
```



```
// Lifting a pair of package
dependencies to
// a pair of models
class Package {
  Model getDependentModel() {
    exists (Model mod |
      depends(this.getModel(),mod)
      and result = mod.getModel()
    )
    and result != this
  }
}
```

# How to Discover Dependencies Between Models in a Multimodel

- ▶ After analysis of all models, **lift the information up the containment hierarchy into the multimodel**
  - Construct inter-model relationships by lifting from inter-package relationships
- ▶ This turns the multimodel into a **macromodel**, a multimodel with model-element constraints
- ▶ The lifted dependencies allow for discovering dependencies between models in a multimodel
  - The precise detailed dependencies give tracing to update models in a multimodel, if something changes

**Macromodel dependency analysis** consists of lifting model-level dependency analysis to inter-model relationships by synthesized attribution

**Macromodel consistency** consists of updating all inter-model relationships and all induced model-level dependencies

# The End

- ▶ Why does ERD and MOF help to define link-consistent link trees?
- ▶ Explain why TgreQL and Xcerpt have similar query styles
- ▶ Why does a megamodel usually build on graphs, not on trees?
- ▶ Why do we need graph query and transformation languages?



## 40.5. Other Graph Query Languages

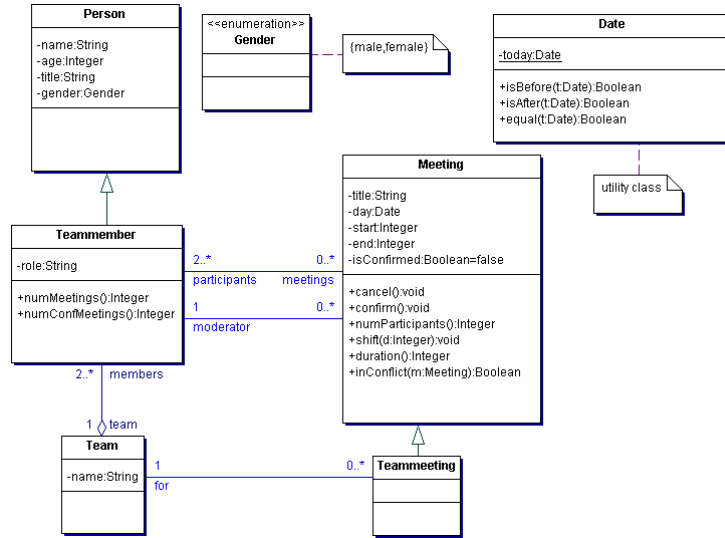


## 40.5.1. Writing Model Constraints by Graph Querying with OCL

- The DDL of OCL is MOF
- .QL is for Java and other GPL
- OCL is for UML-CD

# OCL for Invariants in UML-Class Diagrams

► → course Softwaretechnologie-II





## Examples OCL Invariants

79 Model-Driven Software Development in Technical Spaces (MOST)

- ▶ OCL queries usually start at a specific class; their results define *invariants* on the objects of the class
  - All attributes of a class are visible by default in OCL.
  - Relations between classes define functions
- ▶ Query language uses expressions over these functions

**Example of Invariant:**

```
context Meeting inv: self.end > self.start
```

**Equivalent:**

```
context Meeting inv: end > start
-- self is the context of the query, from which processing starts
```

**Equivalent named constraint:**

```
context Meeting inv startEndConstraint:
self.end > self.start
-- Constraints can constrain attribute values
```

- ▶ FROM and SELECT clauses are modeled via functions:

**Selection constraint:**

```
context Person inv searchForPerson:
allInstances() ->select(p:Person|p.name.StartsWith(„Uwe”))
-- FROM clause is modeled via allInstances() function
-- SELECT clause is modeled via select() function
```

## Examples OCL Invariants

- ▶ **Selection constraint:**  
`context Person inv searchNames:  
allInstances () ->collect(name)  
context Person inv countNames:  
allInstances () ->collect(name) ->size()`
- ▶ **Multiplicity constraint:**  
`context Person inv countNames:  
allInstances () ->collect(name) ->size() < 15`
- ▶ More on OCL: → Course Softwaretechnologie-II, Ch. “Konsistenzprüfung mit OCL”, Dr. Birgit Demuth
- ▶ [www.dresden-ocl.de](http://www.dresden-ocl.de)





## 40.5.2. Graph Querying with GReQL

- ▶ Open source, from University of Koblenz-Landau, Prof. Ebert
- ▶ Applicable to a subset of UML (GrUML)



## TGreQL is similar to .QL

- ▶ But uses a relational notation, from-with-report clauses

```
from RefType tsc, RefType s, RefType t
```

```
where
```

```
  tsc.hasQualifiedName("org.jfree.data.gantt", "TaskSeriesCollection")
```

```
  and s.hasSubtype*(tsc)
```

```
  and t.hasSubtype(s)
```

```
    and not(t.hasName("Object"))
```

```
select s,t
```

.QL

```
from RefType tsc, RefType s, RefType t
```

```
with
```

```
  s hasSubtype*->tsc,
```

```
  tsc.hasQualifiedName("org.jfree.data.gantt", "TaskSeriesCollection"),
```

```
  t hasSubtype->s,
```

```
    not t.hasName("Object")
```

```
report s,t
```

TGreQL



- ▶ TgreQL style is very similar to Xcerpt
- ▶ Implements F-Datalog incl. Transitive closure operator
- ▶ Prof. J. Ebert U Koblenz

Operators:

- \* Transitive closure operator
- + positive transitive closure
- → ← navigation direction
- [] optional path
- () sequence of paths or edges
- | alternative path

```
// construct a call graph
From caller, callee: V{Method}
With caller (
  {isStatementIn}
  [ {isReturnValueOf} ]
  {isActualParameterOf} *
  {isCalleeOf}
) + callee
Report
  caller.name as „Caller“
  callee.name as „Callee“
```

Result (example):

Caller	Callee
main	System.out.println
main	compute
main	twice
main	add
compute	twice
compute	add





## 40.5.3 Model Mappings with Query-View-Transformations (QVT)

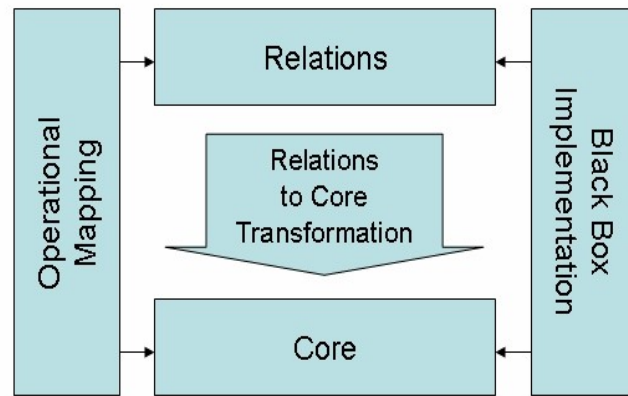
The language of the OMG for model transformations within MDA

OMG: MOF 2.0 Query / Views / Transformations RFP. ad/2002-04-10. Needham, MA: Object Management Group, April 2002.

<http://www.omg.org/cgi-bin/doc?ad/2002-4-10>



DRESDEN  
concept  
Erhalten aus  
Wissenschaft  
und Kultur



From: [[https://de.wikipedia.org/wiki/Datei:QVT-Language-Architecture\\_591x387.jpg](https://de.wikipedia.org/wiki/Datei:QVT-Language-Architecture_591x387.jpg)]

## Transitive Closure with QVT Relations

- ▶ **QVT relations** uses logic expressions on base and derived relations (graph-logic isomorphism)

```
// Transitive Closure in QVT relations,  
// Modeled with recursive relation  
"transitiverelation"  
relation transitiverelation {  
  domain node:Node {  
    // matching attributes  
    name = sameName;  
  }  
  domain node2:Node {  
    // node2 must have the  
    // same name as node  
    name = sameName;  
  }  
  domain node3:Node {  
    // node3 must also  
    // have the same name  
    name = sameName;  
  }  
}
```

```
when {  
  // conditions: base relation must exist  
  baserelation(node,node2) or  
  // or a transitive relation to a base relation  
  (transitiverelation(node,neighbor)  
  and baserelation(neighbor,node2));  
}  
where { // Aufruf einer Transformation  
  makeNodeSound(node);  
}  
}
```





# QVT Tools

Tool			
Eclipse M2M Project	Operational	<a href="http://www.eclipse.org/m2m/">http://www.eclipse.org/m2m/</a>	
Magic Draw	Operational		
MediniQVT	Relational	<a href="http://projects.ikv.de/qvt/wiki">http://projects.ikv.de/qvt/wiki</a>	



## QVT-R uses OCL for Model Search, Query, and Mapping

- ▶ OCL can be called within QVT scripts
  - Two different DQL are combined within a single language

```
// this is QVT
rule checkNoDoubleFeatureInSuperClasses(name:String) {
  from node: Class (
    - OCL query
    node->TransitiveClosure()->collect().exists(s | s.name() = name);
  )
  to
  System.out.println("Error: super class has doubly de ned feature:
  "+s.name());
}
```





## 40.5.4. Graph Invariant Specification with Spider Diagrams



# Spider Diagrams

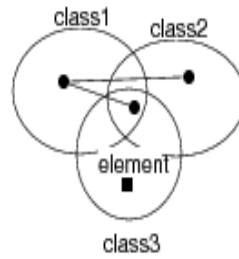
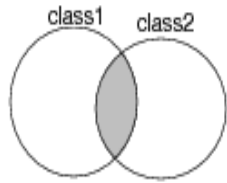
- ▶ [http://en.wikipedia.org/wiki/Spider\\_diagram](http://en.wikipedia.org/wiki/Spider_diagram)
- ▶ S. Kent. Constraint Diagrams: Visualizing Invariants in OO Modelling. Proceedings of OOPSA 97, ACM Press, Oct. 97, pp. 327-341.
- ▶ S. Kent and J. Howse. Mixing Visual and Textual Constraint Languages, UML 99, IEEE press, Oct 1999.
  
- ▶ Spider-Diagramme are equivalent to monadic second-order logic 2. Stufe (MSOL).
  - They include OCL (first-order logic)
- ▶ Source of diagrams: J. Lövdahl, Towards a Visual Editing Environment for the Semantic Web. Linköpings universitet, 2002.

## Simple Spider Diagrams are Extended Venn Diagrams

- ▶ Classes are visualized as venn ellipsoids
- ▶ Set algebra is expressed by intersection of ellipsoids
- ▶ Existential Logic (propositional logic with existential quantifiers) is expressed by **spiders** (hyperedges)

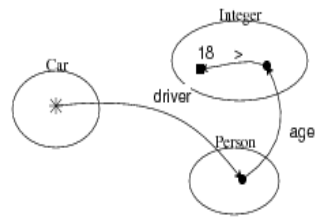
Result =  
 $class1 \wedge class2$

An object of class1 has an object of class2  
and an object in  $class1 \wedge class2 \wedge class3$   
and  $class3 \setminus class1 \setminus class2$  is not empty

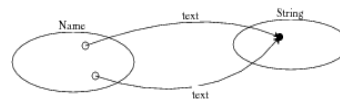


- ▶ All quantifiers are possible (star symbol)

All cars must be driven  
by a person older than 18

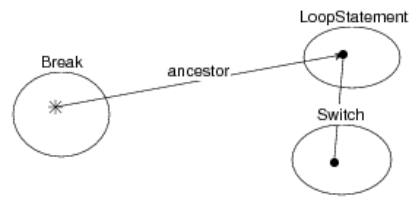


There are no two names that have the same string

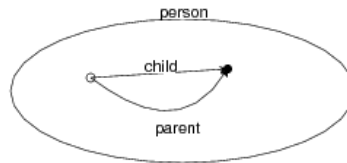


## Other Constraints

All Break statements must have a LoopStatement as ancestor,  
which is related to a Switch state

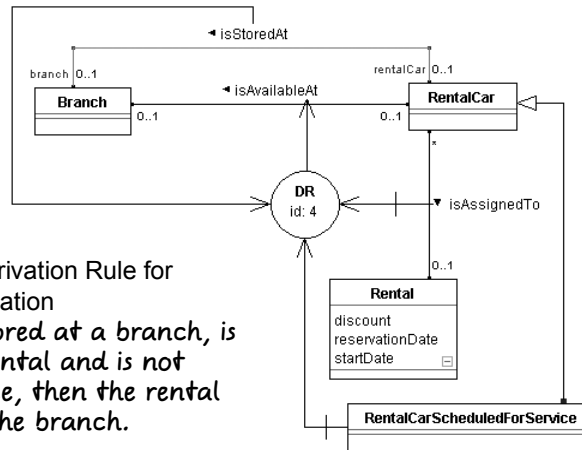


For every person, there is no child that has no parent



## 40.5.5. URML – A UML-like Spider Notation

- ▶ URML <http://oxygen.informatik.tu-cottbus.de/reverse-i1/?q=URML>
- ▶ Emilian Pascalau and Adrian Giurca. Can URML model successfully Drools rules? Proceedings of the 2nd East European Workshop on Rule-Based Applications (RuleApps 2008) at the 18th European Conference on Artificial Intelligence. Patras, Greece, July 23, 2008.
  - <http://ceur-ws.org/Vol-428/paper5.pdf>



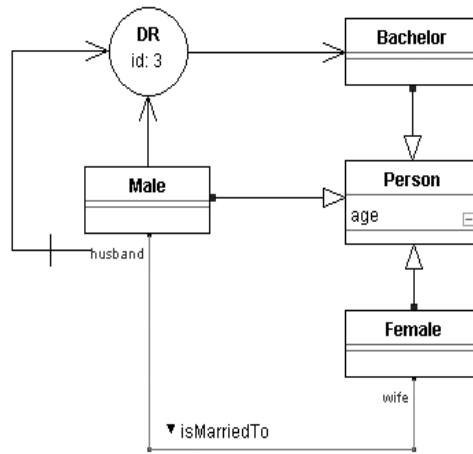
- ▶ Ex: Modeling a Derivation Rule for Defining an Association  
*If a rental car is stored at a branch, is not assigned to a rental and is not scheduled for service, then the rental car is available at the branch.*





# Modeling a Derivation Rule with a Role Condition

*A bachelor is a male that is not a husband.*



# The End

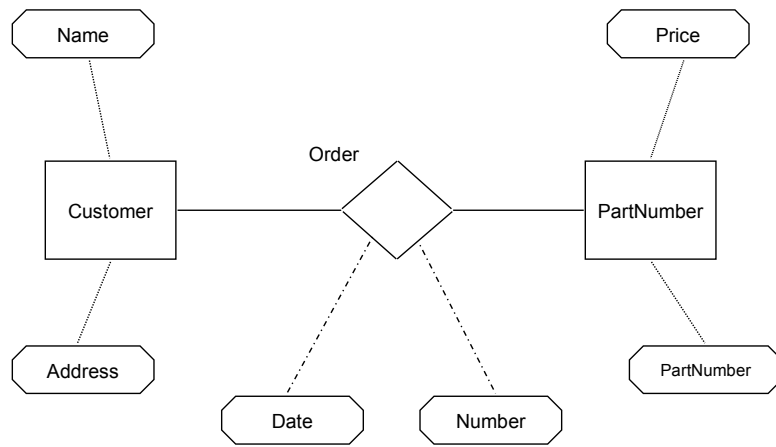
- ▶ Why does ERD and MOF help to define link-consistent link trees?
- ▶ Explain why TgreQL and Xcerpt have similar query styles
- ▶ Why does a megamodel usually build on graphs, not on trees?
- ▶ Why do we need graph query and transformation languages?




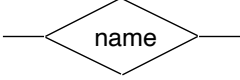
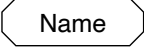
# Appendix

# A Simple ER-Model

- ▶ All “entities” (classes) are represented as “entity-”tables



# ERD Model Elements [Chen]

Notation	Meaning
	<b>Entity type:</b> Set of objects
	<b>Relationship type:</b> Set of relations between entity types
	<b>Attribute:</b> Describes a function or a predicate over an entity
$1, n$ $0 < n$	<b>Cardinality</b> of a relationship type: minimum and maximum amount of neighbors in a relation