



Formal Specification of Non-functional Properties of Component-Based Software

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Two trends underlying this work:

- **CBSE**
 - Current software systems have a high complexity
 - Modularity and component-based techniques can reduce complexity
- **Non-functional properties**
 - have been studied in the small, (Performance Engineering, ...)
 - How to scale up?

→ **Component-based technologies can be a key factor for scaling up non-functional specifications.**

- **Component-based systems open new ways to achieve non-functional properties**
 - Component-level scheduling
 - Buffers, Migration, Replication...



- General Principles
- The Specification
- Outlook/Conclusions

An Example: Timely Response From a System

- Depends on:

- The way the code is written (algorithmic issues)
- The time used other components take to do their part of the work
- Buffering of requests between components
- CPU scheduling and timely availability of other resources

internal properties of the component

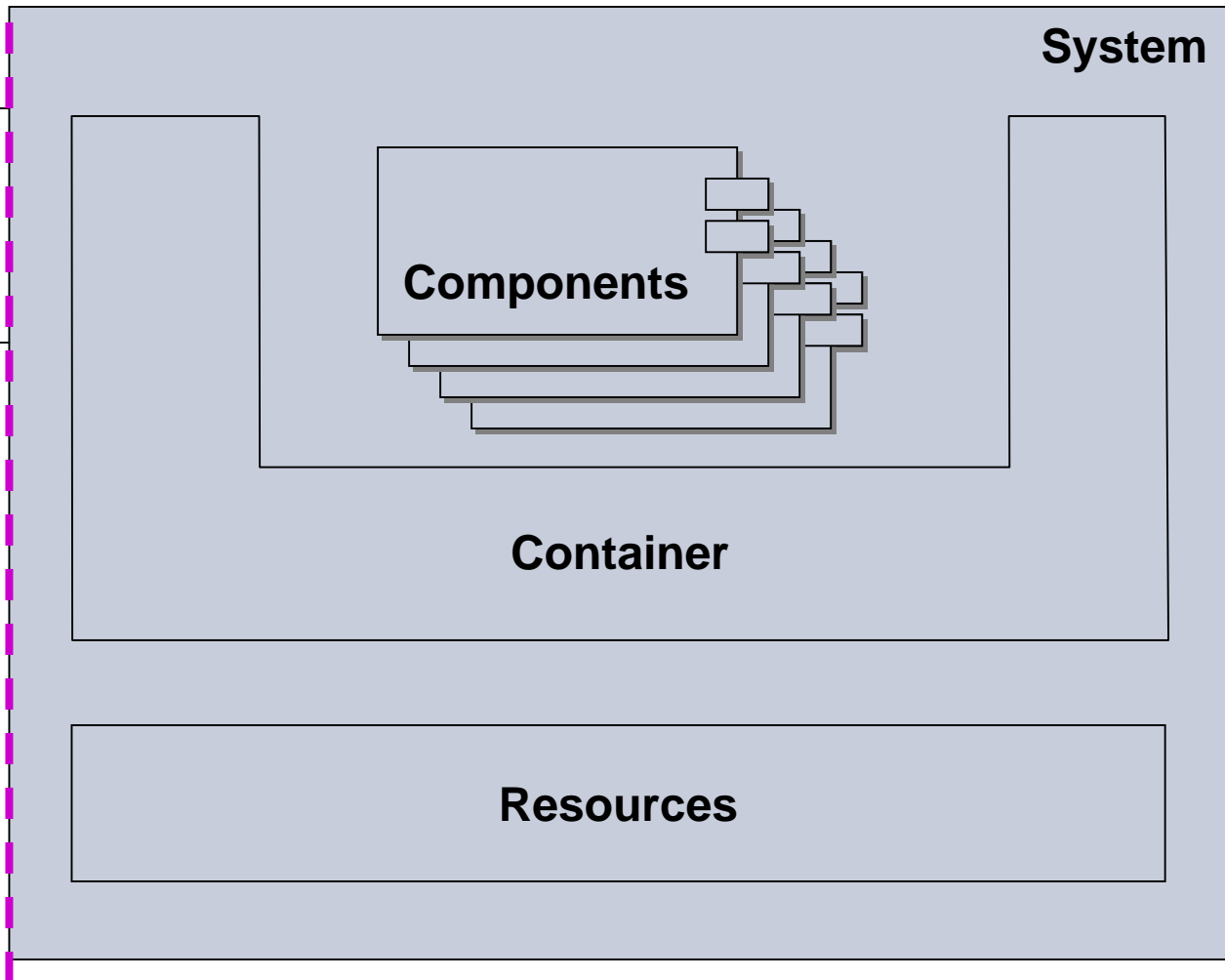
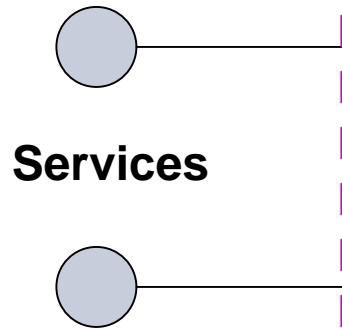
external properties depending on component usage

→ **We can talk about:**

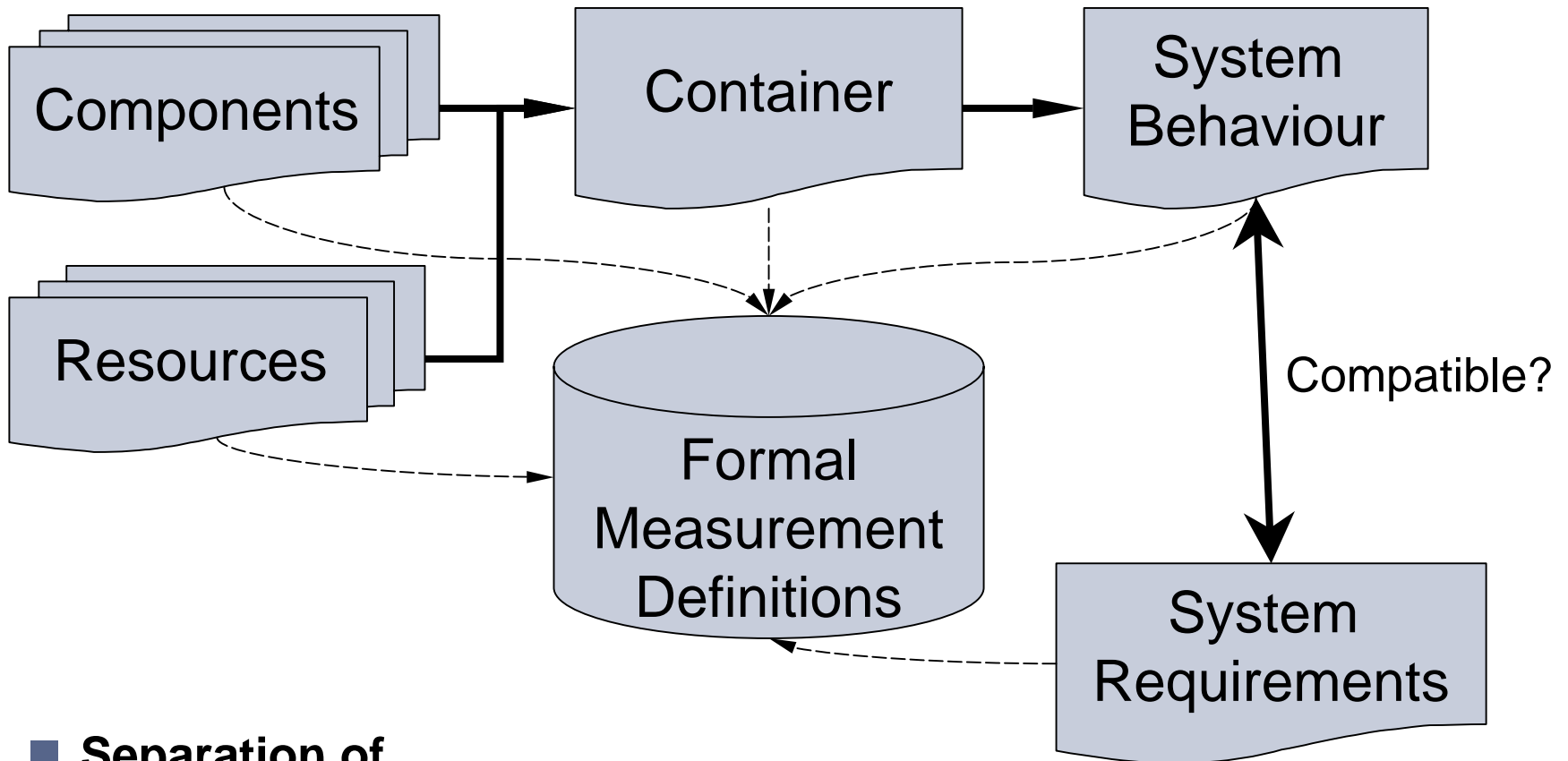
- Execution time = an *intrinsic property* of a component
- Response time = an *extrinsic property* of a system

User View (Extrinsic View)

System View (Intrinsic View)

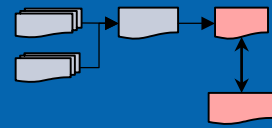


A Specification Framework



■ Separation of

- *Context models*: Models of the “system mechanics”
- *Measurement specification*: Definition of actual non-functional aspects



```

MODULE Service
VARIABLE inState
VARIABLE unhandledRequest

InitEnv ≜ unhandledRequest = FALSE

RequestArrival ≜ unhandledRequest' = TRUE
                ∧ UNCHANGED inState

NextEnv ≜ RequestArrival

EnvSpec ≜ InitEnv
          ∧ □[NextEnv]unhandledRequest

InitServ ≜ inState = Idle

StartRequest ≜ inState = Idle
              ∧ unhandledRequest = TRUE
              ∧ inState' = HandlingRequest
              ∧ unhandledRequest' = FALSE

FinishRequest ≜ inState = HandlingRequest
               ∧ inState' = Idle
               ∧ UNCHANGED unhandledRequest

NextServ ≜ StartRequest ∨ FinishRequest

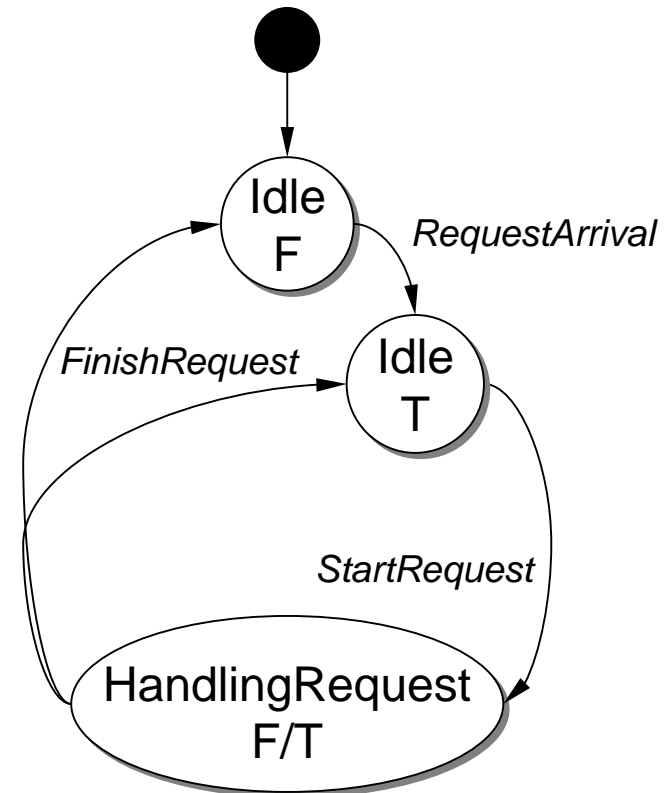
vars ≜ ⟨inState, unhandledRequest⟩

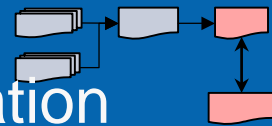
ServiceSpec ≜ InitServ
             ∧ □[NextServ]vars

Service ≜ EnvSpec ⇝ ServiceSpec
    
```

- Currently very simple model:

Service = Single Operation





MODULE *ResponseTimeConstrainedService*

EXTENDS *RealTime*

CONSTANT *ResponseTimeBound*

ASSUME ($ResponseTimeBound \in Real$) \wedge ($ResponseTimeBound > 0$)

VARIABLES *ResponseTime*, *inState*, *unhandledRequest*, *Start*

Serv \triangleq INSTANCE *Service*

Init $\triangleq Start = 0 \wedge ResponseTime = 0$

StartNext $\triangleq Serv!StartRequest \Rightarrow Start' = now$

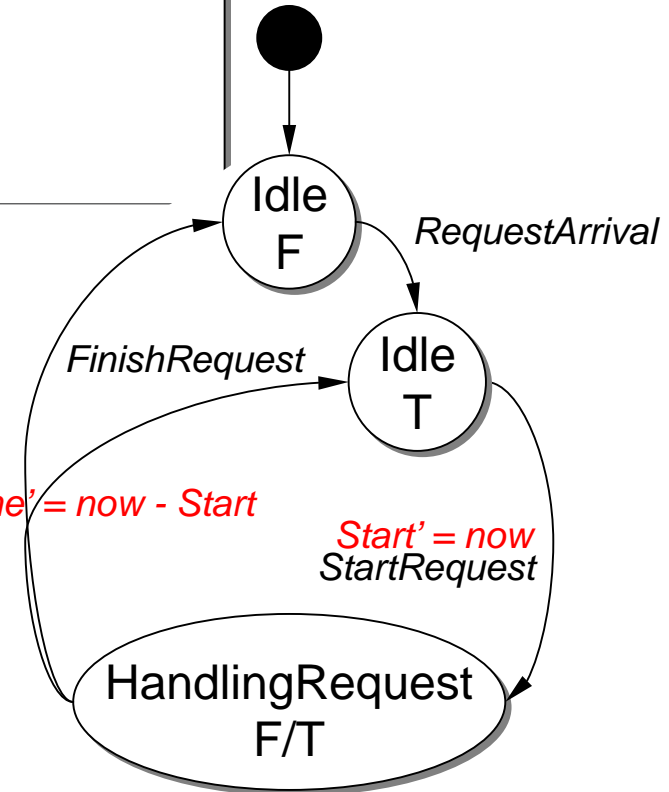
RespNext $\triangleq Serv!FinishRequest \Rightarrow ResponseTime' = now - Start$

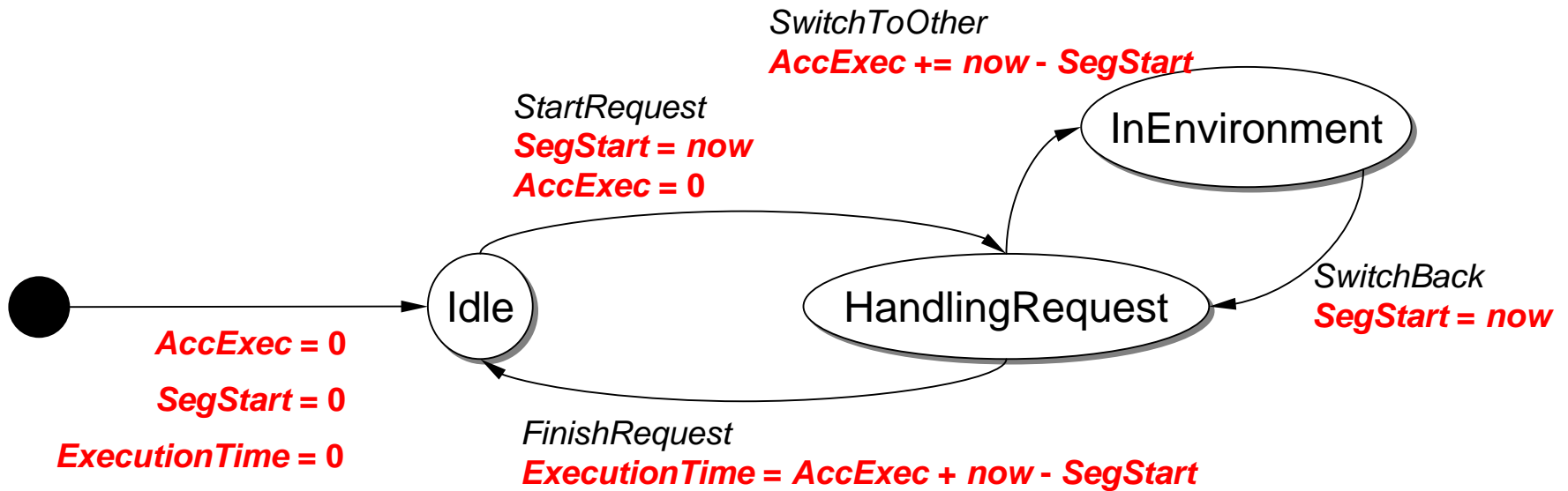
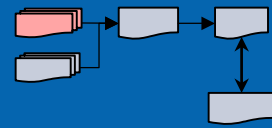
Next $\triangleq StartNext \wedge RespNext$

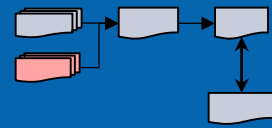
vars $\triangleq \langle inState, unhandledRequest, Start, ResponseTime \rangle$

RespSpec $\triangleq Init \wedge \square[Next]_{vars}$

Service \triangleq
 $Serv!Service$
 $\wedge RTnow(vars)$
 $\wedge RespSpec$
 $\wedge \square(ResponseTime \leq ResponseTimeBound)$







- Example: CPU scheduled by RMS (Rate-Monotonic Scheduling)

MODULE *RMSScheduler*

EXTENDS *Reals*

CONSTANT *TaskCount*
 ASSUME (*TaskCount* ∈ *Nat*) ∧ (*TaskCount* > 0)

CONSTANT *Periods*
 ASSUME *Periods* ∈ [{1 .. *TaskCount*} → *Real*]

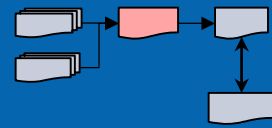
CONSTANT *Wcets*
 ASSUME *Wcets* ∈ [{1 .. *TaskCount*} → *Real*]

VARIABLES *MinExecTime*, *AssignedTo*, *now*

TimedCPUSched \triangleq INSTANCE *TimedCPUScheduler*

Schedulable $\triangleq \sum_{k=1}^{TaskCount} \frac{Wcets[k]}{Periods[k]} \leq TaskCount \left(TaskCount \sqrt{2} - 1 \right)$

RMSScheduler \triangleq *TimedCPUSched!**TimedCPUScheduler*
 ∧ *Schedulable* ⇒ □ *TimedCPUSched!**ExecutionTimesOk*

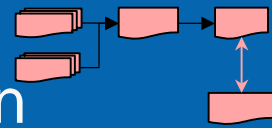


$$\begin{aligned} \text{ContainerPreCond} &\triangleq \text{ExecutionTime} \leq \text{ResponseTime} \\ &\wedge (\text{TaskCount} = 1 \wedge \\ &\quad \text{Periods} = [1 \rightarrow \text{ResponseTime}] \wedge \\ &\quad \text{Wcets} = [1 \rightarrow \text{ExecutionTime}] \wedge \\ &\quad \text{CPUCanSchedule}(\text{TaskCount}, \text{Periods}, \text{Wcets})) \\ &\wedge \text{ComponentMaxExecTime}(\text{ExecutionTime}) \\ &\wedge \text{MinInterrequestTime}(\text{ResponseTime}) \end{aligned}$$

$$\text{ContainerPostCond} \triangleq \text{ServiceResponseTime}(\text{ResponseTime})$$

$$\text{Container} \triangleq \text{ContainerPreCond} \Rightarrow \text{ContainerPostCond}$$

- So far no selection of
 - Concrete component(s)
 - Concrete resource realizations
 - We selected that we need CPU, but didn't say anything about RMS



VARIABLES $TaskCount, Periods, Wcets$

$System \triangleq MyComponent(20)$

$\wedge MyCPU(TaskCount, Periods, Wcets)$

$\wedge MyContainer(20, ResponseTime, TaskCount, Periods, Wcets)$

$ExternalService \triangleq Environment(RequestPeriod) \dashv\triangleright Service(ResponseTime)$

$IsFeasible \triangleq System \Rightarrow ExternalService$

- Future work:
 - Mapping context model \leftrightarrow application model

 - Extend to services delivered by networks of components
 - Extend to multiple properties per specification

 - Apply to other examples
 - other service models (stream based services)
 - stochastic extrinsic properties

- Distinction of **intrinsic/extrinsic** properties of **components/services**
- **System specification** = Composition of component, service, resource and container specifications
 - **Scalability** of the specifications through clear modularization
 - **Formal measurement definitions** as interface between specs.
- **Feasible System** = available components and resources allow the container to provide the required non-functional properties

- Non-functional specifications:
 - Specifically semantics:

Staehli, R., Eliassen, F., Agedal, J.Ø., Blair, G.: *Quality of service semantics for component-based systems*. In: Middleware 2003 Companion, 2nd Int'l Workshop on Reflective and Adaptive Middleware Systems.
 - Specification approaches:
 - Characteristic-specific
... lots
 - Measurement-based

Agedal: *Quality of service support in development of distributed systems* → CQML

Selic: *A generic framework for modelling resources with UML*

Skene et al: *Precise Service-Level Agreements*